Virtualization Technologies

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Outline

1. Hypervisor based – virtualization
2. Containers
3. Uni-kernel
4. Virtualization and serverless computing
Hypervisor Based - Virtualization
On Virtualization

- Key concepts
- Type I (bare metal) vs. Type 2 (hosted)
- Solutions for non virtualizable CPUs
  - Binary Translation
  - Para-virtualization
Basic concepts

1. On operating systems
2. Virtual machine, hypervisor
4. Examples of benefits
Operating systems

Some of the motivations

- Only one single thread of CPU can run at a time on any single core consumer machine
- Machine language is tedious
Operating systems

Operating systems bring a level of abstraction on which multiple processes can run at a time – Deal among other things with:

- Multiplexing
- Hardware management issues

However only one operating system can run on a bare single core consumer machine
virtual machines and hypervisors

- Systems virtualization dates back to the 60s
- IBM experimentation with “time sharing systems”
virtual machines and hypervisors

- Why virtual machines?
  - How to develop software that run on different operating systems without the purchase of several servers
  - How to run legacy applications that run on legacy operating systems
  - Job migrations
virtual machines and hypervisors

Virtual machine (VM)
- Software that provides same inputs / outputs and behaviour expected from hardware (i.e. real machine) and that supports operations such as:
  - Create
  - Delete
  - Migrate
  - Increase resources

Hypervisor
- Software environment that enables operations on virtual machines (e.g. XEN, VMWare) and ensures isolation
virtual machines and hypervisors

Hypervisors (earlier known as Virtual Machine Monitor (VMM))

- Software environment that enables operations on virtual machines (e.g. XEN, VMWare) and meeting the following requirements:
  - Virtual machines identical to physical machines (same input / same output)
  - Efficiency
  - Isolation
virtual machines, hypervisors


From reference [1] – Note: There is a small error in the figure
Examples of Benefits

All benefits are due to the possibility to manipulate virtual machine (e.g. create, delete, increase resources, migrate), e.g.

- Co-existence of operating systems
- Optimization of hardware utilization
- Job migration
Advanced concepts

1. Bare metal vs. hosted hypervisor
2. Full virtualization vs. Para-virtualization
3. Binary translation
Type I vs Type II Hypervisor

Some concepts

- Hardware
- Host OS
  - Runs on the hardware (Type 2)
- Guest OS
  - Runs on top of the hypervisor

Note: Type II hypervisor is sometimes called “Hosted Hypervisor”
Type I vs Type II Hypervisor

AS Tanenbaum an H Bos, Modern Operating Systems, 5th edition, Published by Pearson (May 29, 2022) © 2023

(a) Type 1 hypervisor
- Hardware
  - CPU, disk, network, interrupts, etc.

(b) Type 2 hypervisor
- Guest OS
  - (e.g., Windows)
- Host OS
  - (e.g., Linux)
- Hardware
  - (CPU, disk, network, interrupts, etc.)
Type I vs Type II Hypervisor

Types of hypervisor

- Type I – bare metal
  - Installed on bare hardware
  - Examples
    - Citrix XEN server
    - VMWARE ESX/ESXI
Type I vs Type II Hypervisor

Types of hypervisor

- Type 2 – hosted
  - Runs on top of host operating system
  - Examples:
    - VMWare workstation
    - VirtualBox
Type I vs Type II Hypervisor

Type I - Bare metal

- Hypervisor installed on bare hardware

  - Advantages (compared to type II)
    - Performance (No additional software layer to go through)
    - Security (No possible attack through host operating system)

  - Drawbacks (compared to type II)
    - Host operating system needs to be “ported” on top of hypervisor
    - Complexity depends on the type of virtualization (Full virtualization vs. para-virtualization)
Type I vs Type II Hypervisor

Type II - Hosted

- Hypervisor installed on top of host operating system

  - Drawbacks (compared to type I)
    - Performance (need to go through host operating system)
    - Security (i.e. Possibility to attack through host operating system)

  - Advantages (compared to type I)
    - Host operating system is re-used as it is (No need to port it)
    - No change required to applications running on top of host operating system
Full virtualization vs. Para-virtualization

More on operating systems fundamentals

- User process vs. Kernel process
- User mode vs. Kernel mode

Note: In user mode some instructions called sensitive instructions should not be executed
Full virtualization vs. Para-virtualization

More on operating systems fundamentals

- Sensitive vs. non sensitive instruction
  - Sensitive
    - Has the capacity to interfere with supervisor software functioning (e.g. OS) and should be executed only in kernel mode (i.e. privileged mode)
    - Write OS memory vs. read OS memory

Note: When a user process sends a sensitive instruction, the instruction is trapped by the CPU and is not executed.
Full virtualization vs. Para-virtualization

Back to hypervisors

- In addition to user mode and kernel mode
  - Virtual user mode
  - Virtual kernel mode
Full virtualization vs. Para-virtualization

Back to hypervisors

- Scenarios discussions
  - CPU able to send trap to hypervisors (virtualizable CPUs)
  - CPU unable to send traps to hypervisors (non virtualization CPUs)
Full virtualization vs. Para-virtualization

Back to hypervisors (CPUs able to send traps to hypervisors)

AS Tanenbaum an H Bos, Modern Operating Systems, 5th edition, Published by Pearson (May 29, 2022) © 2023
Full virtualization vs. Para-virtualization

Could all CPU architectures be fully virtualized?

- The case of Intel x86 CPU architectures
  - Cannot be fully virtualized because they cannot generate convenient traps to hypervisors
    - Need to extended
Full virtualization vs. Para-virtualization

Definitions

**Full virtualization**
- Hypervisor enables virtual machines identical to real machine
  - Problematic for architectures such as Intel x86
Full virtualization vs. Para-virtualization

Definitions

Para-virtualization
- Hypervisor enables virtual machine that are similar but not identical to real machine
  - A solution to the problem of CPU architectures that cannot be virtualized
    - Prevents user programs from executing sensitive instructions
- Note:
  - Para-virtualization is not the only solution to the problem
Full virtualization vs. Para-virtualization

Full virtualization

- Advantages
  - Possibility to host guest operating systems with no change since virtual machines are identical to real machines

- Disadvantages
  - Not always feasible (e.g. Intel x86)
    - There are work around (e.g. binary translation)
  - Some guest operating systems might need to see both virtual resources and real resources for real time applications
Full virtualization vs. Para-virtualization

Para - virtualization

- **Advantages**
  - Feasible for all CPU architectures
  - Performance – Compared to:
    - Full virtualization
    - Other approaches to architectures that could not be virtualized (e.g. binary translation)

- **Disadvantages**
  - Need to modify guest operating systems
Full virtualization vs. Para-virtualization

Para-virtualization

- Alternatives to para-virtualization
  - Binary translation (e.g. VMWare ESX server)
    - Leads to full virtualization
    - No need to re-write “statically” guest operating systems
      - i.e. guest OS can be installed without change
    - Interpretation of guest code (OS + application)
      - “Rewrites” dynamically guest code and insert traps when necessary
Full virtualization vs. Para-virtualization

Para - virtualization

- Alternatives to para-virtualization
  - Binary translation
    - Disadvantages / penalties
      - Performance
      - However, optimization is possible, e.g.
        » Adaptive translation (i.e. optimize the code being translated)
Alternatives to Hypervisor Based - Virtualization
Containers and Unikernels

- Issues with hypervisors
- Alternatives (Containers and unikernels)
Hypervisor

In a hypervisor based – approach, a VM includes the application + full blown operating system (e.g. Linux Debian, Linux Red Hat)

- OS on virtual machine needs to boot
  - Slow starting time for application
- Resources are not used in an efficient manner
  - Linux kernel replicated in each VM that runs linux.
Proposed Solutions

Back to operating systems basics

- The two components of an operating system
  - Kernel
    - Interacts with the hardware and manages it (e.g. write/read a disk partition)
  - Libraries
    - Set of higher level functions accessible to programs via system calls
      - Enable function like create / read / delete file while hiding the low level operations on the hard disk
Alternatives

VM vs container vs Unikernel

T. Goethals et al., Unikernels vs. Containers: An In-Depth Benchmarking Study in the Context of Microservice Application, IEEE SC2 Conference, November 2018
On containers

Operating system (Kernel) virtualization:

- Kernel offers isolated spaces to run containers

  Containers
  - Applications packaged with their run time environment that run on a same kernel
  - Run as processes, but with isolated file system, networking, CPU and memory resources
On containers

Operating system (Kernel) virtualization:

- Kernel offers isolated spaces to run containers

  - Containers
    - Hosted by container engine (e.g. Docker Engine)
    - Need to be deployed, managed and orchestrated (e.g. Kubernetes)
On containers

Operating system (Kernel) virtualization:

- Kernel offers isolated spaces to run containers
  - Some pros / cons
    - Less memory footprint
      » Do not include kernel
    - Faster start up time
      » Kernel does not need to boot
On containers

Operating system (Kernel) virtualization:

- Kernel offers isolated spaces to run containers
- Some pros / cons
  - Works only in environments in which you have given operating system kernel + its libraries (e.g. Linux kernel + Linux distributions)
  - Less secure than VM
    - Challenge:
      - Trade-off between isolation and performance / efficiency
On Unikernels

Application + Tiny run time:

- Tiny run time
  - Not the whole OS like VM
  - Not the whole libraries like containers
    » Only the function required by the applications
    » Static binding
  - Can run as a tiny VM or a tiny container
On Unikernels

Pros and cons:

- Smaller footprint
- Boot up faster
- Less flexible
  - Addition / removal of functionality requires re-compilation
Virtualization and Serverless Computing
Server-less computing
Server-less Computing
(Function as a Service)

- Introduction
- Architecture
- Pros / Cons
Introduction

Server-less does not mean there is no server !!!

- There are indeed servers !!!
  - However the servers are completely transparent to the cloud users, unlike (Virtual Machine (VM), Containers, Uni-kernel)
    - Server-less computing might actual rely on VMs or containers or uni-kernels
  - Cloud users deal with functions (No need to deal with the infrastructure)
    - thus Functions as a Service (FaaS)
Architecture

Examples of platforms

- Amazon Lambda
- Microsoft Azure function
- Kuberless
Architecture

Y. Li et al., Serverless Computing: State of the Art, Challenges and Opportunities, IEEE Transactions on Services Computing, March/April 2023

<table>
<thead>
<tr>
<th>Serverless</th>
<th>Virtual Machine</th>
<th>Container</th>
</tr>
</thead>
<tbody>
<tr>
<td>Users' application</td>
<td>Users' application</td>
<td>Users' application</td>
</tr>
<tr>
<td>Serverless platform (OpenWhisk, OpenFaas ...)</td>
<td>Configured Environment</td>
<td>Docker</td>
</tr>
<tr>
<td>Orchestrator</td>
<td>Guest Kernel</td>
<td></td>
</tr>
<tr>
<td>VM / Container</td>
<td>Host Kernel</td>
<td>Host Kernel</td>
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</tbody>
</table>

Hardware Resource
Architecture

Principles

1) Applications built as a set of functions

2) When there is a request for a given function, a run time environment (e.g. VM, container, uni-kernel) is launched with the function code + libraries

3) The run time is terminated after the execution of the function
Architecture

Serverless front-end
- Function programming
- Function serving

Platform: Modules such as:
- Run time
- Repository
- Scheduler
Architecture

Y. Li et al., Serverless Computing: State of the Art, Challenges and Opportunities, IEEE Transactions on Services Computing, March/April 2023
(Flow view)
Architecture

PAditya et al, Will Serverless Computing Revolutionize NFV, Proceedings of the IEEE, April 2019

Fig. 1. Serverless platform architecture.
Architecture

Load balancer:
- Self explanatory

Front end:
- End user interface

Message bus and scheduler:
- Mediation between front ends and execution engines
Architecture

**Load balancer:**
- Self explanatory

**Front end:**
- End user interface

**Message bus and scheduler:**
- Mediation between front ends and execution engines
  - Relies on a publication / subscription principles
Architecture

**Execution engine:**
- Self explanatory
  - Might rely on VM, containers and uni-kernels

**Storage sub-system:**
- States
- Persistent data
Pros (Examples)

- No real / virtual server management by cloud users

- Resource Efficiency and low cost

- Built-in scalability
Cons (Examples)

- **Most cited:**
  - Start up latency

- **Others:**
  - Learning curve of the new programming model (e.g. stateless functions + events)
Pros vs Cons

P Aditya et al, Will Serverless Computing Revolutionize NFV, Proceedings of the IEEE, April 2019

- Decision to be made on case by case basis

![Diagram showing latency requirements for various applications](image-url)
Pros vs Cons

- Decision to be made on case by case basis

Fig. 4. Cost comparison between Amazon Lambda (serverless) and Amazon EC2 (VMs) for spiky workload. In the gray region, serverless is 100x cheaper.
The End