

ELEC 6131 – Error Detecting and Correcting Codes
Midterm
March 5, 2019
Time: 90 minutes

- 1) Consider the polynomial $p(x) = x^4 + x + 1$ over GF(2).
 - a) Show that $p(x)$ is primitive (2 Marks).
 - b) List all elements of GF(16) generated by $p(x)$ in power and polynomial form (4 Marks).
 - c) Show that α^3 a root of $f(x) = x^4 + x^3 + x^2 + x + 1$ (2 Marks).
 - d) What are other roots? (2 Marks)

- 2) Consider a (7, 4) code with the following Generating matrix:

$$G = \begin{bmatrix} 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 \end{bmatrix}.$$

Transform it to systematic form (3 Marks).

- 3) Consider an (8, 4) code with following equations:

$$c_0 = u_0 + u_1 + u_2$$

$$c_1 = u_1 + u_2 + u_3$$

$$c_2 = u_0 + u_1 + u_3$$

$$c_3 = u_0 + u_2 + u_3$$

$$c_j = u_{7-j} \text{ for } j = 4, 5, 6, 7.$$

- a) Find the generating matrix of the code. (1 Mark).
 - b) Find the parity check matrix of the code (2 Marks).
 - c) What is the minimum distance of the code (2 Marks).
 - d) How many errors can this code correct (1 Mark)? How many errors it can detect (1 Mark)?
- 4) Show that for any (n, k) binary linear code with minimum distance $2t+1$, following inequality holds (5 Marks):

$$n - k \geq \log_2 \left[1 + \binom{n}{1} + \binom{n}{1} + \dots + \binom{n}{t} \right].$$

- 5) A code has the following parity check matrix:

$$H = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \end{bmatrix}$$

Find the transmitted codeword if the received bits are:

$r = [1, e_1, 0, 1, 1, e_2, 0, 1, e_3, 0]$ where, e_1, e_2, e_3 are erasures. (5 Marks).