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About This Book

Replication Server® maintains replicated data at multiple sites on a network. Organizations with geographically distant sites can use Replication Server to create distributed database applications with better performance and data availability than a centralized database system can provide.

This book introduces distributed database systems built upon replication technology and helps you design a replication system.

Audience

The Replication Server Design Guide is for everyone who uses Replication Server. If you are new to Replication Server, begin with this book for an introduction to Replication Server and the applications that use replicated data.

If you are designing a new application for Replication Server, you should read this book before you install the Replication Server software. Use the information in this book to plan your replication system so that you will know where to install the software components that make up your replication system.

How to use this book

The information in this book is organized as follows:

- Chapter 1, “Introduction” introduces Replication Server and its features.
- Chapter 3, “Implementation Strategies” describes models for implementing your replication system design.
- Chapter 4, “Planning for Backup and Recovery” describes the preventive and corrective measures you can use to recover from replication system failures.
- Chapter 5, “Introduction to Replication Agents” describes Sybase Replication Agent products that you can use to replicate data from a database that is not a Sybase® Adaptive Server or a Sybase SQL Server database.
• Chapter 6, “Replicating Data into Foreign Data Servers” describes the replication system components that you need to replicate data into a data server other than Adaptive Server or Sybase SQL Server.

• Chapter 7, “International Replication Design Considerations” describes how to configure languages, character sets, and sort orders for an international environment.

• Appendix A, “Capacity Planning” explains how to estimate the amount of disk space needed for Replication Server partitions.

• Appendix B, “Log Transfer Language” describes the Log Transfer Language (LTL) used by Replication Agents to send transaction operation and stored procedure invocation data to a Replication Server.

Related documents

• The Replication Server installation and configuration guides for your platform describe Replication Server installation and explain how to use the rs_init program.

• What’s New in Replication Server 12.0? contains information about the new features in this release.

• The Replication Server Administration Guide contains an introduction to replication systems and in-depth information about Replication Server administration issues. It also describes how to use the Sybase Central™ plug-in—Replication Server Manager™ (hereafter referred to as RSM)— and the Replication Server online help.

• The Replication Server Reference Manual contains the syntax and detailed descriptions of the commands in the Replication Command Language (RCL).

• The Replication Server Troubleshooting Guide contains information that helps you diagnose and correct problems in a replication system.

• The Replication Server online help contains information about using Replication Server with Sybase Central.

Other sources of information

Use the Sybase Technical Library CD and the Technical Library Product Manuals web site to learn more about your product:

• Technical Library CD contains product manuals and technical documents and is included with your software. The DynaText browser (included on the Technical Library CD) allows you to access technical information about your product in an easy-to-use format.

Refer to the Technical Library Installation Guide in your documentation package for instructions on installing and starting Technical Library.
• Technical Library Product Manuals web site is an HTML version of the Technical Library CD that you can access using a standard web browser. In addition to product manuals, you’ll find links to the Technical Documents web site (formerly known as Tech Info Library), the Solved Cases page, and Sybase/Powersoft newsgroups.

To access the Technical Library Product Manuals web site, go to Product Manuals at http://sybooks.sybase.com.

Sybase certifications on the web

Technical documentation at the Sybase web site is updated frequently.

❖ For the latest information on product certifications and/or the EBF Rollups:

2 In the Browse section, click on What’s Hot.
3 Select links to Certification Reports and EBF Rollups, as well as links to Technical Newsletters, online manuals, and so on.

❖ If you are a registered SupportPlus user:

2 In the Browse section, click on What’s Hot.
3 Click on EBF Rollups.
   You can research EBFs using Technical Documents, and you can download EBFs using Electronic Software Distribution (ESD).
4 Follow the instructions associated with the SupportPlus Online Services entries.

❖ If you are not a registered SupportPlus user, and you want to become one:

You can register by following the instructions on the Web.

To use SupportPlus, you need:

1 A Web browser that supports the Secure Sockets Layer (SSL), such as Netscape Navigator 1.2 or later
2 An active support license
3 A named technical support contact
4 Your user ID and password
Whether or not you are a registered SupportPlus user:

You may use Sybase’s Technical Documents. Certification Reports are among the features documented at this site.

1. Point your web browser to Technical Documents at http://techinfo.sybase.com
2. In the Browse section, click on What’s Hot.
3. Click on the topic that interests you.

Conventions

This section describes the style and syntax conventions, RCL command formatting conventions, and icons used in this book.

Style conventions

Syntax statements that display the syntax and options for a command are printed as follows:

```
alter user user
set password new_passwd
[verify password old_passwd]
```

See “Syntax conventions” on page xii for more information.

Examples that show the use of Replication Server commands are printed as follows:

```
alter user louise
set password somNIfic
verify password EnnuI
```

Command names, command option names, program names, program flags, keywords, functions, and stored procedures are printed as follows:

Use `alter user` to change the password for a login name.

Variables, parameters, and user-supplied words are in italics in syntax and in paragraph text, as follows:

The `set password new_passwd` clause specifies a new password.

Names of database objects such as databases, tables, columns, and datatypes, are in italics in paragraph text, as follows:

The `base_price` column in the `Items` table is a `money` datatype.

Names of replication objects, such as function-string classes, error classes, replication definitions, and subscriptions, are in italics.

Syntax conventions

Syntax formatting conventions are summarized in the following table.

Examples combining these elements follow.
Table 1: Syntax formatting conventions

<table>
<thead>
<tr>
<th>Key</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>{ }</td>
<td>Curly braces mean you must choose at least one of the enclosed options. Do not include braces in the command.</td>
</tr>
<tr>
<td>[ ]</td>
<td>Brackets mean you may choose or omit enclosed options. Do not include brackets in the command.</td>
</tr>
<tr>
<td></td>
<td>Vertical bars mean you may choose no more than one option (enclosed in braces or brackets).</td>
</tr>
<tr>
<td>,</td>
<td>Commas mean you may choose as many options as you need (enclosed in braces or brackets). Separate your choices with commas, to be typed as part of the command. Commas may also be required in other syntax contexts.</td>
</tr>
<tr>
<td>( )</td>
<td>Parentheses are to be typed as part of the command.</td>
</tr>
<tr>
<td>...</td>
<td>An ellipsis (three dots) means you may repeat the last unit as many times as you need. Do not include ellipses in the command.</td>
</tr>
</tbody>
</table>

Obligatory choices

- Curly braces and vertical bars – choose only one option.
  
  \{red | yellow | blue\}

- Curly braces and commas – choose one or more options. If you choose more than one, separate your choices with commas.
  
  \{cash, check, credit\}

Optional choices

- One item in square brackets – choose it or omit it.
  
  [anchovies]

- Square brackets and vertical bars – choose none or only one.
  
  [beans | rice | sweet_potoatoes]

- Square brackets and commas – choose none, one, or more options. If you choose more than one, separate your choices with commas.
  
  [extra_cheese, avocados, sour_cream]

Repeating elements

An ellipsis (...) means that you may repeat the last unit as many times as necessary. For the alter replication definition command, for example, you can list one or more columns and their datatypes for the add clause or the add searchable columns clause:

```
alter replication definition replication_definition
{add column datatype [, column datatype]... | add searchable columns column [, column]... | replicate {minimal | all} columns}
```
<table>
<thead>
<tr>
<th>RCL command formatting</th>
<th>RCL commands are similar to Transact-SQL® commands. The following sections present the formatting rules.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Command format and command batches</td>
<td>• You can break a line anywhere except in the middle of a keyword or an identifier. You can continue a character string on the next line by typing a backslash () at the end of the line.</td>
</tr>
<tr>
<td></td>
<td>• Extra spaces are ignored, except after a backslash. Do not enter any spaces after a backslash.</td>
</tr>
<tr>
<td></td>
<td>• You can enter more than one command in a batch unless otherwise instructed.</td>
</tr>
<tr>
<td></td>
<td>• RCL commands are not transactional. Each command is executed independently and is not affected by the completion status of other commands in the batch. However, syntax errors in a command prevent Replication Server from executing subsequent commands in a batch.</td>
</tr>
<tr>
<td>Case sensitivity</td>
<td>• Keywords in RCL commands are not case sensitive. You can enter them in any combination of uppercase or lowercase letters.</td>
</tr>
<tr>
<td></td>
<td>• Case sensitivity in identifiers and character data depends on the sort order that is in effect.</td>
</tr>
<tr>
<td></td>
<td>• If you use a case-sensitive sort order such as “binary,” you must enter identifiers and character data in the correct combination of uppercase and lowercase letters.</td>
</tr>
<tr>
<td></td>
<td>• If you use a sort order that is not case sensitive, such as “nocase,” you can enter identifiers and character data in any combination of uppercase or lowercase letters.</td>
</tr>
<tr>
<td>Identifiers</td>
<td>Identifiers are names you give to database or replication objects. They include object names, column names, variable names, and parameter names.</td>
</tr>
<tr>
<td></td>
<td>• Identifiers can be 1–30 bytes long (equivalent to 1–30 single-byte characters) and must begin with a letter, the @ sign, or the _ character.</td>
</tr>
<tr>
<td></td>
<td>• Replication Server function parameters are the only identifiers that can begin with the @ character. Function parameter names can include 30 characters after the @ character.</td>
</tr>
<tr>
<td></td>
<td>• After the first character, identifiers can include letters, digits, and the #, $, or _ characters. Spaces are not allowed.</td>
</tr>
</tbody>
</table>
| Parameters in function strings | Parameters in function strings have the same rules as identifiers, except that:
• They are enclosed in question marks (?). This allows Replication Server to locate them in the function string. Use two consecutive question marks (??) to represent a literal question mark in a function string.

• The exclamation point (!) introduces a parameter modifier that indicates the source of the data to be substituted for a parameter at runtime. See the Replication Server Administration Guide for a complete list of modifiers.

**Datatype support**

Replication Server supports all Adaptive Server datatypes.

User-defined datatypes are not supported. The timestamp, double precision, nchar, and nvarchar datatypes are indirectly supported; they are mapped to other datatypes. Columns using the timestamp datatype are mapped to varbinary(8).

See the Replication Server Reference Manual for more information about supported datatypes, including how to format them.

**Icons**

Illustrations in this book use icons to represent the components of a replication system.

**Replication Server**

![Replication Server icon]

This icon represents Replication Server, the Sybase server program that maintains replicated data on a local-area network (LAN) and processes data transactions received from other Replication Servers on the wide-area network (WAN).

**Adaptive Server or data server**

![Adaptive Server or data server icon]

This icon represents Adaptive Server, the Sybase data server. Data servers manage databases containing primary or replicate data. Replication Server also works with SQL Server and with heterogeneous data servers, so, unless otherwise noted, this icon can represent any data server in a replication system.
Replication Agent

This icon represents a Replication Agent, a replication system process or module that transfers transaction log information for a primary database to a Replication Server. The Replication Agent for Adaptive Server is RepAgent; the Replication Agent for Sybase SQL Server is Log Transfer Manager (LTM). Sybase provides Replication Agent products for Adaptive Server Anywhere, DB2, Informix, Microsoft SQL Server, and Oracle data servers.

Except for RepAgent, which is an Adaptive Server thread, all Replication Agents are separate processes. In general, this icon only appears when representing a Replication Agent that is a separate process.

Client application

This icon represents a client application. A client application is a user process or application connected to a data server. It may be a front-end application program executed by a user or a program that executes as an extension of the system.

Replication Server Manager

This icon represents Sybase Central’s Replication Server Manager (RSM), a software tool that lets a Replication System Administrator monitor, analyze, troubleshoot, and administer a replication system.

RSM consists of:
• An **RSM Client** (Sybase Central + the Replication Server plug-in): an Open Client application that lets you monitor, diagnose, and manage your replication system. The RSM Client includes a graphical user interface (GUI) integrated with Sybase Central.

• An **RSM Server**: an Open Server application that translates and communicates requests from an RSM Client. The RSM Server also monitors for and responds to replication system events.

**If you need help**

Each Sybase installation that has purchased a support contract has one or more designated people who are authorized to contact Sybase Technical Support. If you cannot resolve a problem using the manuals or online help, please have the designated person contact Sybase Technical Support or the Sybase subsidiary in your area.
CHAPTER 1

Introduction

This chapter introduces Replication Server and describes how it is used to create and maintain distributed data applications.

### Centralized and distributed database systems

In the traditional enterprise computing model, an Information Systems department maintains control of a centralized corporate database system. Mainframe computers, usually located at corporate headquarters, provide the required performance levels. Remote sites access the corporate database through wide-area networks (WANs) using applications provided by the Information Systems department.

Changes in the corporate environment toward decentralized operations have prompted organizations to move toward distributed database systems that complement the new decentralized organization.

Today’s global enterprise may have many local-area networks (LANs) joined with a WAN, as well as additional data servers and applications on the LANs. Client applications at the sites need to access data locally through the LAN or remotely through the WAN. For example, a client in Tokyo might locally access a table stored on the Tokyo data server or remotely access a table stored on the New York data server.
Advantages of replicating data

In a distributed database environment, mainframe computers may be needed at corporate or regional headquarters to maintain sensitive corporate data, while clients at remote sites use minicomputers and server-class workstations for local processing.

Both centralized and distributed database systems must deal with the problems associated with remote access:

• Network response slows when WAN traffic is heavy. For example, a mission-critical transaction-processing application may be adversely affected when a decision-support application requests a large number of rows.

• A centralized data server can become a bottleneck as a large user community contends for data server access.

• Data is unavailable when a failure occurs on the network.

Advantages of replicating data

The performance and availability problems associated with remote database access can be solved by replicating the data from its source database to a local database. Replication Server provides a cost-effective, fault-tolerant system for replicating data.

Replication Server keeps data up to date in multiple databases so that clients can access local data instead of remote, centralized databases. Compared to a centralized data system, a replication system provides improved system performance and data availability and reduces communication overhead.

Because it transfers transactions, not rows, Replication Server maintains the integrity of replicated data across the system, while also increasing data availability. Replication Server also allows you to replicate stored procedure invocations, further enhancing performance.

Improved performance

In a distributed replication system, data requests are completed on the local data server without the client having to access the WAN. Performance for local clients is improved because:

• LAN data transfer rates are faster than WAN data transfer rates.
• Local clients share local data server resources instead of competing for central data server resources.
• Traffic and contention for locks are reduced considerably because local decision-support applications are separated from centralized OLTP applications.

Greater data availability

In a distributed replication system, data is replicated at local and remote sites, so clients can continue to work regardless of what happens at the primary data source or over the WAN.
• When a failure occurs at a remote site, clients can continue to use local copies of replicated data.
• When a WAN failure occurs, clients can continue to use local replicated data.
• When the local data server fails, clients can switch to replicated data at another site.

When WAN communications fail, Replication Servers at other sites store transactions in stable queues (disk storage) so that replicated tables at the unavailable site can be brought up to date when communications resume. When a replicated function is initiated in a source database, it is stored in stable queues until it can be delivered to the destination site.

Data distribution with Replication Server

Replication Server works to distribute data over a network by:
• Providing application developers and system administrators with a flexible publish-and-subscribe model for marking data and stored procedures to be replicated
• Managing replicated transactions while retaining transaction integrity across the network

Because Replication Server replicates transactions—incremental changes instead of data copies—and stored procedure invocations, not the stored procedures themselves, it provides a high-performance distributed data environment while maintaining data integrity.
**Publish-and-subscribe model**

In a functioning Replication Server system, transactions occurring in a source database are detected by a Replication Agent and transferred to the local Replication Server, which distributes the information across LANs and WANs to Replication Servers at destination sites. These Replication Servers in turn update the target database according to the requirements of the remote client. If a network or system component fails, data in the process of being delivered is temporarily stored in queues. When the failed component returns to operation, the replication system resynchronizes copies of the data and normal replication resumes.

The primary data is the source of the data that Replication Server replicates in other databases. You “publish” data at primary sites to which Replication Servers at other sites “subscribe.” You first create a **replication definition** to designate the location of the primary data. The replication definition describes the structure of the table and names the database that contains the primary copy of the table. For easier management, you may collect replication definitions into **publications**.

The creation of a replication definition or publication does not, by itself, cause Replication Server to replicate data. You must create a subscription against the replication definition (or the publication) to instruct Replication Server to replicate the data in another database. A subscription resembles a SQL **select** statement. It can include a **where** clause to specify which rows of a table you want to replicate in the local database, allowing you to replicate only the necessary data.

Beginning with the 11.5 version of Replication Server, you can have multiple replication definitions for a primary table. Replicate tables can subscribe to different replication definitions to obtain different views of the data.

Once you have created subscriptions to replication definitions or publications, Replication Server replicates transactions to databases with subscriptions for the data.

**Replicated functions**

Replication Server lets you replicate Adaptive Server stored procedure invocations asynchronously between databases. This method can improve performance over normal data replication by encapsulating many changes in a single **replicated function**. Because they are not associated with table replication definitions, replicated functions can execute stored procedures that may or may not modify data directly.
You can replicate stored procedure invocations from a primary database to a replicate database, or from a replicate database to a primary database. See “Using applied functions” on page 39 and “Request functions” on page 67 for details.

With replicated functions, you can execute a stored procedure in another database. A replicated function allows you to:

- Replicate the execution of an Adaptive Server stored procedure to subscribing sites
- Improve performance by replicating only the name and parameters of the stored procedure rather than the actual changes

Like tables, replicated stored procedures may have replication definitions (called function replication definitions) and subscriptions. When a replicated stored procedure executes, the Replication Server passes its name and execution parameters to subscribing sites, where the corresponding stored procedure executes.

Create function replication definitions at the primary data site. Replication Server supports applied functions and request functions:

- An applied function is replicated from a primary to a replicate database. You create subscriptions at replicate sites for the function replication definition and mark the stored procedure for replication in the primary database.
- A request function is replicated from a replicate to a primary database. There is no subscription for a request function. You mark the stored procedure for replication in the replicate database.

**Transaction management**

Replication Server depends on data servers to provide the transaction processing services needed to protect their stored data. To guarantee the integrity of distributed data, data servers must comply with such transaction-processing conventions as atomicity and consistency.

Data servers that store primary data provide most of the concurrency control needed for the distributed database system. If a transaction fails to update a table with primary data, Replication Server does not distribute the transaction to other sites. When a transaction does update primary data, Replication Server distributes the changes and, unless a failure occurs, the update succeeds at all sites that have subscribed to the data.
Replication Server uses optimistic concurrency control to maintain replicated data consistency. This method differs from a pessimistic distributed concurrency control method—such as the two-phase commit—because it processes failures after they occur.

Optimistic concurrency control has these advantages in a replication system:

- It promotes high availability of data because it does not lock the data for the duration of the distributed transaction.
- It requires fewer system resources to process a transaction.
- It does not require data servers to have special distributed transaction processing features in order to participate in a distributed transaction.

Failed replicated transactions

A modification to primary data may fail to update a replicate copy of the data at another site. The primary version is the “official” copy, and updates that succeed at the primary database are expected to succeed at sites with replicate copies.

Some reasons for the failure of an update to a replicated table are:

- The data server’s maintenance user login name does not have the permissions needed to update the replicate data.
- The replicate and primary versions of the data are inconsistent after a system recovery.
- A client updates replicate data directly rather than updating the primary version.
- The data server storing the replicate table has constraints that are not enforced by the data server storing the primary version.
- The data server storing the replicated copy of the table rejects the transaction due to a system failure, such as lack of space in the database.

When a transaction fails, Replication Server receives an error from the data server. Data server errors are mapped to Replication Server error actions. The default action for a failed transaction is to write a message in the Replication Server error log (including the message returned by the data server) and then suspend the database connection. After you correct the cause of the failure, you can resume the database connection and Replication Server will retry the failed transaction.
You also can have Replication Server record a failed transaction in the exceptions log (a set of three tables in the RSSD) and continue processing the next transaction. Refer to “Replication Server” on page 8 for a description of the RSSD.

If you use the exceptions log, you must manually resolve the transactions that are saved there to make the replicated data consistent with the primary data. In some cases, the process can be automated by encapsulating the logic for handling the rejected transactions in an intelligent application program.

Transactions that modify data in multiple data servers and databases

A transaction that modifies primary data in more than one data server may require additional concurrency control. According to the transaction processing requirements, either all of the operations in the transaction are performed or none of them are performed. If a transaction fails on one data server, it must be rolled back on all other data servers updated in the transaction.

Normally, there is exactly one Replication Agent for each primary database. If a single transaction updates multiple primary databases, that transaction is replicated as multiple independent transactions, one for each primary database. Or, you can select to encapsulate such transactions in a single stored procedure, which then flows as an atomic unit to subscribing sites.

Replication system components

Replication Server has an open architecture that allows you to build a replication system from existing systems and applications and add to it as your organization grows and changes.

Figure 1-1 illustrates one site in a WAN-based, distributed database system that uses Replication Server. The sections that follow describe each component.
Replication Server

A Replication Server at each site coordinates the data replication activities for
the local data servers and exchanges data with Replication Servers at other
sites.

A Replication Server:

• Receives primary data transactions from databases via replication agents
  and distributes them to sites with subscriptions for the data

• Receives transactions from other Replication Servers and applies them to
  local databases

Replication Server system tables store information needed to accomplish these
tasks. The system tables include descriptions of the replicated data and
replication objects such as replication definitions and subscriptions, security
records for Replication Server users, routing information for other sites, access
methods for the local databases, and other administrative information.

Replication Server system tables are stored in an Adaptive Server database
called the Replication Server System Database (RSSD). An RSSD is assigned
to each Replication Server. An Adaptive Server data server with an RSSD can
also store application databases.
You use Replication Command Language (RCL) or the Replication Server Manager component of Sybase Central to manage information in Replication Server. RCL commands, which resemble SQL commands, can be executed on Replication Server using isql, the Sybase interactive SQL utility. \textit{Replication Server Reference Manual} is the complete reference for RCL. Information about Replication Server Manager can be found in the \textit{Replication Server Administration Guide} and in Replication Server’s online help.

\section*{Partitions and stable queues}

Replication Server stores messages on disk to make sure that they can be delivered following a failure. When you install a Replication Server, you allocate an initial \textit{disk partition} that Replication Server uses for its disk storage. You can add additional partitions when you have finished installing the Replication Server.

The partition is either a raw disk device or operating system file. Because UNIX operating systems buffer file I/O, you may not be able to completely recover data following a failure. On such a system, use operating system files for partitions only in a test environment. Use a raw disk partitions for production environments. See the \textit{Replication Server Reference Manual} for more information about adding partitions.

Replication Server allocates \textit{stable queues} from its disk partitions for the routes and connections it serves. Messages are saved in the stable queues at least until the messages are confirmed as received at their destination.

The amount of disk space you should allocate for Replication Server partitions depends on the size of transactions and the transaction rate for your application. Stable queues act as buffers for data as it flows through your replication system. If a remote site’s Replication Server is unreachable during a network failure, the primary Replication Server stores transactions in a stable queue until communication is restored. The more space allocated for disk partitions, the longer the Replication Server can queue data without interrupting operations in the primary database.

Appendix A, “Capacity Planning” explains in detail how to calculate partition space for a Replication Server.
Data servers

Data servers manage databases containing primary or replicated data. Clients use them to store and retrieve data and to process queries and transactions. Replication Server maintains replicated data in data servers by logging in as a database user.

Replication Server supports heterogeneous data servers through an open interface. Any system for storing data can be used as a data server if it supports a set of required data operations and transaction processing directives.

See “Heterogeneous data server support” on page 15 for more information on data server requirements.

Replication Agent

A Replication Agent transfers transaction log information, which represents changes made to primary data, from a data server to a Replication Server for distribution to other databases. The Replication Agent for Adaptive Server is RepAgent, which is an Adaptive Server thread. The Replication Agent for Sybase SQL Server is a separate process called the LTM (Log Transfer Manager).

The Replication Agent reads the database transaction log and transfers log records for replicated tables and replicated stored procedures to the Replication Server that manages the database. The Replication Server reconstructs the transaction and forwards it to the sites that have subscriptions for the data.

A Replication Agent is needed for each database that contains primary data or for each database where replicated stored procedures are executed. A database that contains only copies of replicated data and has no replicated stored procedures does not require a Replication Agent.

Because RepAgent is an Adaptive Server thread, most system diagrams in this book do not include a Replication Agent icon. In Chapter 5, “Introduction to Replication Agents” however, the Replication Agent described is a separate process and, for clarity, system diagrams contain the Replication Agent icon.
Client applications

A client application is a user program that accesses a data server. When the data server is an Adaptive Server, applications can be programs created with Open Client Client Server, Embedded SQL, PowerBuilder®, or any other front-end development tool compatible with Sybase Client/Server Interfaces (C/SI).

Client applications should update only the primary data. Replication Server distributes the changes to the other sites. Client applications that do not modify data do not need to distinguish between primary and replicated data.

Connecting replication system components

Replication Server, Replication Agent, and Adaptive Server use Sybase Client/Server Interfaces to communicate over a network. In addition, Replication Server uses routes and connections to send messages to other Replication Servers and databases. The following sections describe the interfaces file, routes, and connections.

Interfaces file

Server programs such as data servers, Replication Servers, and Replication Agents are registered in an interfaces file so that client applications and other server programs can locate them. Generally, one interfaces file at each site contains entries for all of the local and remote Replication Servers and data servers. The entry for each server includes its unique name and the network information that other servers and client programs need to connect to it.

Use ds_edit to maintain your interfaces file.

Note If you are using network-based security, available with Replication Server version 12.0, use the directory services of your network security mechanism (rather than the interfaces file) to register Replication Servers, Adaptive Servers, and gateway software. Refer to the documentation that comes with your network-based security mechanism for details.
**Connecting replication system components**

**Routes and connections**

Routes and connections allow Replication Servers to send messages to each other and to send commands to databases. A *route* is a one-way message stream that sends requests from one Replication Server to another Replication Server. A *connection* is a message stream from a Replication Server to a database. Replication Server uses a *logical connection* to represent the active and standby databases in a warm standby application.

To replicate data from one database into another, you must first establish the routes and connections that allow Replication Server to move the data from its source to its destination.

When you add a database to your replication system, Sybase Central or `rs_init` creates the connection for you. You never have to create a connection directly unless you are replicating data into a database that is not an Adaptive Server or SQL Server database.

If you have more than one Replication Server in your replication system, you must create routes between them. If you have only one Replication Server, you do not need to create routes.

Figure 1-2 illustrates connections and routes between three Replication Servers, one database storing primary data, and four databases storing replicated data.

*Figure 1-2: Connections and routes*
When you create a route from a primary Replication Server to a replicate Replication Server, transactions flow from the primary server to the replicate server.

If you plan to execute replicated stored procedures in a replicate database to update a primary database, you must also create a route from the replicate Replication Server to the primary Replication Server.

**Direct and indirect routes**

In a replication system with one primary Replication Server and many replicate Replication Servers, you can use indirect routes to reduce the load on the primary Replication Server. Indirect routes allow Replication Server to send messages to multiple destinations through a single intermediate Replication Server.

Routes with intermediate sites have important advantages:

- **Reduced WAN volume**
  
  Replication Server distributes one copy of a message to each intermediate site. Replication Servers at the intermediate sites duplicate the messages for each of their outgoing queues.

- **Reduced Replication Server load**
  
  Additional Replication Servers running on separate computers share the processing load, reducing the amount of processing required of Replication Servers at primary sites.

- **Fault tolerance**
  
  Messages stored at intermediate sites can be used to recover from partition failures at remote sites. See the *Replication Server Administration Guide* for details.

Figure 1-3 shows how message distribution is handled using intermediate sites. The message follows a direct route to the intermediate sites. From the intermediate site, it follows a direct route to the local site. With this routing arrangement, the primary site sends two messages rather than eight.
Connecting replication system components

**Figure 1-3: Routes in a hierarchical configuration**

Intermediate sites reduce primary site message volume, but they increase the time between updates at the primary and replicate servers. Plan your routes carefully; use only the number of intermediate sites required.

If intermediate sites are not used, routes are set up in a star configuration, as Figure 1-4 illustrates.

**Figure 1-4: Direct routes in a star configuration**
When a row is updated at a primary site, the primary Replication Server sends messages through the WAN to each remote site that has a subscription for the row. For example, in Figure 1-4, New York can send identical data through eight different routes. If there are many sites, the network is quickly overloaded with redundant messages.

Creating routes in a hierarchical arrangement allows load balancing by reducing the number of connections and messages distributed from the primary site. Additional Replication Servers running on separate computers share the processing load, reducing the amount of processing required of Replication Servers at primary sites.

**Heterogeneous data server support**

Replication Server’s open architecture allows the use of heterogeneous data servers in replication systems. Support for Adaptive Server is built into Replication Server. You build the interface for other data servers by creating the instructions Replication Server will use with the data server. This section contains a brief overview of open architecture components. See Chapter 6, “Replicating Data into Foreign Data Servers” for more information.

The open architecture includes:

- Client/Server Interfaces
- Replication Agents
- Error classes and error processing actions
- Functions, function strings, and function-string classes

**Client/Server Interfaces (C/SI)**

Sybase Client/Server Interfaces (C/SI) includes routines and protocols for client/server communication. Replication Server uses C/SI to connect with data servers as a client. If a data server doesn’t support C/SI, you can use a Sybase gateway product (such as DirectConnect) to allow Replication Server to access the data server. A gateway accepts connections from C/SI clients and, using the data server’s native interface, submits client requests to the data server for processing. When the data server returns the results, the gateway uses C/SI routines to return them to the client.
**Replication Agents**

A Replication Agent is required for every database that stores primary data or initiates replicated functions. A Replication Agent reads the data server transaction log to detect changes to primary data and executions of replicated stored procedures. The transaction log provides a reliable source of information about primary data modifications because it contains records of committed, recoverable transactions.

Chapter 5, “Introduction to Replication Agents” describes the Replication Agents for other data servers than Adaptive Server or Sybase SQL Server. Replication Agents for some non-Sybase data servers—such as DB2, Informix, Microsoft SQL Server, and Oracle—are available from Sybase.

**Processing data server errors**

Replication Server processes the errors and status codes returned by data servers according to your instructions. Each vendor’s data server has a different set of error codes. Replication Command Language (RCL) commands allow you to:

- Create an error class to group together the error code mappings for a database.
- Assign error actions, such as warn, retry_log, and stop_replication, to data server error codes.
- Associate an error class with a database.

**Functions, function strings, and function-string classes**

In order to operate in a heterogeneous database environment, Replication Server differentiates database commands from the functions it uses to distribute data server requests to other sites. A *function* is a Replication Server object that represents a data server operation such as insert, delete, and begin transaction. Replication Server uses function strings to convert functions into data-server-specific commands. A *function string* is a template that Replication Server uses to generate a command the data server can interpret as a transaction-control directive or data-modification instruction.
A function-string class is the set of all function strings used with a database. Function-string classes are provided for Adaptive Server and DB2 data servers. Function strings for transaction control directives are defined just once for each function-string class. Function strings to insert a row, delete a row, or update a row are defined once for each replicated table in a database.

A function string can contain variables—identifiers enclosed in question marks (?)—that represent the values of columns, procedure parameters, system-defined information, and user-defined variables. Replication Server replaces the variables with actual values before sending the function strings to the data server.

Function strings can be used to generate either RPC or database commands such as SQL statements, depending on their format. An RPC-formatted function string contains a remote procedure call followed by a list of data parameters. Embedded variables can be used to assign runtime values to the parameters. Replication Server interprets RPC function strings, builds a remote procedure call, and replaces the variables with runtime data values. An RPC can execute a registered procedure in an Open Server gateway to a data server or a stored procedure in an Adaptive Server.

A language-formatted function string passes a database command to the data server. Replication Server does not attempt to interpret the string, except to replace embedded variables with runtime data values. For example, several relational database servers use the SQL database language. A SQL command should be represented as a language function string.

RPC function strings can be more efficient than language function strings because the network packets sent from Replication Server are more compact.

**Replication Server security**

Replication Server security includes password-protected login names and a permission system based on the grant and revoke commands. In addition, Replication Server 12.0 supports third-party security services that ensure secure message transmission over the network and that enable user authentication.
Replication Server security

Login names
Each Replication Server uses login names that are distinct from data server login names. Many clients do not need a Replication Server login name because they accomplish their work through data server applications.

Replication Server login names
When you install a Replication Server, rs_init creates Replication Server login names that other Replication Servers and Replication Agents use to log in to Replication Server.

A Replication System Administrator creates and manages the Replication Server login names and passwords used to manage replicated data or replication system functions such as the addition of new users or a route change. Passwords can be encrypted.

Data server login names
Data server login names are used with a client application to connect to a data server. The application uses data stored by the data server, including data replicated by Replication Server. A Database Administrator creates and manages data server login accounts.

Client access to replicated copies of tables is also managed by the Database Administrator. Since Sybase recommends that replicated tables be read-only, clients may be permitted to view replicated data but should be prevented from inserting, deleting, or updating rows.

To modify replicated tables, a client must modify the primary data so that the Replication Server can distribute the changes to replicate databases that have subscriptions for the data. To modify a table, a client must have a login name on the data server where the primary data is stored as well as the permissions necessary to update the primary data.

Data server maintenance user login name
Replication Server uses a maintenance user login name for each local data server database that contains replicated tables. It is called the maintenance user login name because Replication Server uses it to maintain replicated tables. The Database Administrator must make sure that the maintenance user login name has the permissions needed to update the replicated tables in the database.
Normally, transactions applied by the maintenance user are filtered by the Replication Agent so they are not replicated out of a database. In certain applications, however, these transactions must be replicated. See Chapter 3, “Implementation Strategies” for more information on these application types.

Permissions

The grant and revoke commands are used to grant and revoke Replication Server permissions for clients. Table 1-1 lists the permissions that can be granted to clients.

<table>
<thead>
<tr>
<th>Permission</th>
<th>Capabilities</th>
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</thead>
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<td>sa</td>
<td>Gives recipient System Administrator capabilities. Clients with sa permission can perform any Replication Server command.</td>
</tr>
<tr>
<td>create object</td>
<td>Allows recipient to create, alter, or drop Replication Server objects, including replication definitions and subscriptions.</td>
</tr>
<tr>
<td>primary subscribe</td>
<td>Gives the recipient permission to create a subscription in a primary database, but not to create other objects. To create a subscription at a remote site, a client needs create object permission in the replicate database and create object or primary subscribe permission in the primary database.</td>
</tr>
<tr>
<td>connect source</td>
<td>This permission is required for login names used by the Replication Agent. It allows the recipient to execute the RCL commands that are reserved for Replication Agents.</td>
</tr>
</tbody>
</table>

Network-based security

With a third-party network-based security mechanism, users are authenticated by the security system at login. Authentication is the process of verifying that users are who they say they are. Users receive a credential that can be presented to remote servers in lieu of a password. As a result, users have seamless access to the components of the replication system through a single login.
Network-based security mechanisms also provide a variety of data-protection services, such as message confidentiality and out-of-sequence checking. Replication Server requests the service, prepares the data, and sends it to the network server for encryption or validation. Once the service is performed, data is returned to the requesting Replication Server for distribution.

Once a secure pathway has been established, data can move in both directions. Both ends of the pathway must support the same security mechanism and be configured the same. The security mechanism must be installed on all machines that make use of network security, and Replication Server 12.0 must be installed on all participating machines.

See the Replication Server Reference Manual for information about network-based security in the replication system.

Summary

- Replication Server maintains copies of tables in different databases on a network. The replicated copies provide two advantages to users at the sites: faster response and greater availability.
- A replication system built on Replication Server uses the Sybase Client/Server Interfaces to connect components—data servers, Replication Servers, Replication Agents, and client applications.
- Replication Server is designed with an open interface that allows non-Sybase data servers to be included in a replication system.
- One copy of a table is the primary version. All others are replicated copies.
- Using subscriptions, a replicated copy of a table may contain a subset of the rows in a table.
- Replication Server security consists of login names, passwords, and permissions. Replication Server also supports third-party network-based security mechanisms.
- Replication Server uses optimistic concurrency control that processes failures when they occur. Compared to other methods, optimistic concurrency control provides greater data availability, uses fewer resources, and works well with heterogeneous data servers.
This chapter discusses several topics that are important for you, as an application designer and user, to explore before you implement Replication Server.

### Application types

Determining the type of Replication Server application you build will in part determine the type of replication strategy you use to implement the application. Chapter 3, “Implementation Strategies” covers various replication scenarios.

Replication Server supports the following basic application types:

- Decision support
- Distributed online transaction processing (OLTP)
- Remote OLTP using request functions
- Warm standby

Each of these application types differs in the way it updates primary data and in the way it distributes primary and replicated data within the replication system.
Decision-support applications

Decision-support clients and production online transaction processing (OLTP) clients use data in different ways. Decision-support clients execute lengthy queries that hold locks on tables to guarantee serial consistency. OLTP clients, on the other hand, execute transactions that must complete quickly and cannot accept the delays caused by decision-support clients’ data locks. The two types of clients do not interfere with each other if they maintain separate copies of replicated tables.

Replication Server off-loads processing associated with decision-support applications from a centralized online transaction processing application onto local servers. The primary database manages the transaction processing, and the replicate databases at local sites handle requests for information from decision-support clients. Providing a separate, reference-only copy of the data allows OLTP systems to continue unobstructed.

Multiple copies of tables containing primary data used in decision-support applications can be maintained at a single site or at multiple sites over the network.

Multiple copies at a single site

Sometimes it is useful to maintain multiple replicate copies of a table at a single site. A subscription specifies the database where Replication Server maintains the replicated data. You can create subscriptions for multiple copies of a table by creating the table in different databases at the same site and then creating subscriptions for each one.

If OLTP and decision-support clients are on the same LAN, one Replication Server can manage both the primary data and the replicate data. Figure 2-1 illustrates such an arrangement.
For best performance, the databases are usually maintained by different data servers. The subscriptions can request different subsets of the data to be maintained in each database, so the replicated copies do not have to be identical.

If you must have two copies of a table in the same database, you can use multiple replication definitions for a primary table. One replication definition could have publishers as the replicate table name, and the other publishers2. Multiple replication definitions are also useful if you want different replicates to receive different column subsets. See “Multiple replication definitions” on page 60.

Another way to update multiple tables in an Adaptive Server database is to use stored procedures. Code the multiple updates in the stored procedures and write Replication Server function strings to execute the stored procedures. You can also use replicated functions and stored procedures to update multiple tables.
Application types

Multiple copies distributed over a network

When copies of tables are distributed over a WAN in a decision-support application, all updates are performed by applications executing at the primary site and are distributed to the remote sites that have subscriptions for the data. Figure 2-2 illustrates this arrangement.

*Figure 2-2: Multiple LAN decision support replicate*

![Diagram of Multiple LAN decision support replicate](image-url)

This type of system uses the centralized primary maintenance method of updating primary data. Clients at remote sites subscribe to replication definitions or publications of primary data. They do not update primary data. See “Centralized primary maintenance” on page 30 for information on this method.

Distributed OLTP applications

Although some distributed transaction-processing applications maintain centralized primary data, others fragment primary data among replicate sites.
A primary fragment is a horizontal segment of a table that holds the primary version of a set of rows. Updates are applied to the primary version first and are then distributed to sites that have replicated copies of the data.

Sites that are responsible for, or own, portions of a table by definition have multiple primary fragments. For example, the salesdetail table in Figure 2-3 has primary fragments in Chicago, New York, and San Francisco:

*Figure 2-3: Table with multiple primary fragments*

A key constructed from one or more columns identifies the primary fragment where a row belongs. The key for the salesdetail table is the stor_id column.

- Rows with “7067” in the stor_id column belong to the primary fragment at the Chicago site.
- Rows with “5023” in the stor_id column belong to the primary fragment at the New York site.
- Rows with “8042” in the stor_id column belong to the primary fragment at the San Francisco site.

There are three application models based on multiple primary fragments:

- Distributed primary fragments – in this model, tables at each site contain both primary and replicated data. Updates to the primary version are distributed to other sites. Updates to non-primary data are received from the primary site.
- Corporate rollup – in this model, multiple primary fragments maintained at remote sites are consolidated into a single aggregate replicate table at a central site.
- Redistributed corporate rollup – this model is the same as the corporate rollup model, except that the consolidated table is redistributed.
Remote OLTP using request functions

Replicated functions can be used to execute transactions remotely. Client applications at remote sites can update primary data asynchronously with request functions. The client application does not require a network connection to the primary site, and the request can be accepted by the Replication Server even when the primary site is not available.

Once the request function executes the stored procedure in the primary database, Replication Server may replicate some or all of the data changes made in the primary database. These changes can be propagated to replicate databases as data rows or as applied functions.

Local update applications

A local update application allows clients at a remote site to see the updates they have entered before the replication system returns them from the primary site. For example, if a customer account is updated at a remote site, clients at the site can see the results of the transaction even if the primary site is not accessible.

Local updates can be performed by using a pending updates table. For each replicated table, a corresponding local table contains provisional updates—updates that have been submitted to the primary site, but that have not been returned through the replication system. Client applications update the pending transactions table and, at the same time, send a request function to the primary site. See “An example using a local pending table” on page 70 for information on implementing this type of application.

When the update succeeds against the primary copy, it is distributed to remote sites, including the site where the transaction originated. You can create a function string or replicated stored procedure to update the replicated table and delete local updates from the pending table. This makes it possible for client applications to know which transactions have been confirmed and which are pending.
Standby applications

A **warm standby application** is a Replication Server application that maintains a pair of Adaptive Server or SQL Server databases, one of which functions as a standby copy of the other.

Client applications generally update the **active database**, while Replication Server maintains the **standby database** as a copy of the active database. Replication Server keeps the standby database consistent with the active database by replicating transactions retrieved from the active database transaction log.

If the active database fails, or if you need to perform maintenance on the active database or data server, you can switch to the standby database so that client applications can resume work with little interruption.

Figure 2-4 illustrates a warm standby system.

*Figure 2-4: A warm standby system*

The two databases in a warm standby application appear as a single logical database in the replication system. Depending on your application, this logical database may not participate in replication, or it may be a primary database or a replicate database with respect to other databases in the replication system.

Several Replication Server and RepAgent features explicitly support warm standby applications. See the *Replication Server Administration Guide* for more detailed information about warm standby applications.
Effects of loose consistency on applications

Data in a replicate database is “loosely consistent” with the data in the primary database. Replicate data lags behind primary data by the amount of time it takes to distribute updates from the primary database. The latency can be measured in seconds (or less) when the system is working properly. If a component fails—if, for example, a network connection is temporarily lost—updates can be delayed for minutes, hours, or days.

Although replicate data may lag behind primary data, it is transactionally consistent with the primary data. Replication Server delivers transactions to replicate databases in the order they are committed in the primary database. This ensures that the replicate data goes through the same series of states as the primary data.

The importance of loose consistency varies by application and even within an application. Some applications tolerate average system lag time and occasional, longer delays with no special provisions. Some require special handling when the lag time becomes too great, and some require special handling for certain types of transactions.

Controlling risks in high-value transactions

The delay introduced by data replication adds risk to some business decisions. For example, a banking application that approves cash withdrawals uses the most current account information available to verify that a customer’s balance is sufficient to cover the withdrawal. If withdrawals processed in a primary database have not reached the replicate database, an application using the replicate database risks approving a withdrawal that exceeds the funds available in the customer’s account.

To limit risk, the banking application can distinguish between high-value transactions and low-value transactions. For example, it might approve a $100 withdrawal based on the account balance in the local replicate database, but it would log in to the primary database to check the account balance before approving a $1000 withdrawal.
Measuring lag time

A measure of the latency for a replicate site can be used to limit risks for some transactions. A small lag time indicates that the primary and replicate data are nearly consistent. An extensive lag time indicates a greater potential difference between the primary and replicate data.

An application can use a measure of lag time to:

- Limit risk by restricting the transactions clients can execute as the lag time increases. For example, the banking application described in the previous section could include the lag time in its approval formula. It might allow withdrawals of up to $1000 based on the balance in the local replicated table when the latency is less than a minute. If the lag time is more than a minute, however, the application would log in to the primary database to approve withdrawals of more than $500.

- Provide clients with a “performance meter” for data replication. Clients can use an estimate of lag time as an advisory. For example, a decision-support user, noting that the lag time is high, might wait for the local data to catch up with the primary data, and for the lag time to decrease, before running an analysis based on replicate data.

Note  Sybase Central includes a heartbeat feature that allows you to track lag time. Refer to Replication Server’s online help for more information.

Methods for updating primary data

In a replication system, the primary copy of a data row is the definitive copy. An update committed in the primary database is authoritative and is distributed to all databases with subscriptions for the data.

Replication Server distributes transactions after they are committed in the primary database. Because changes made to replicate data are not distributed, make the data in replicate databases read-only for clients and route all client transactions to the primary database.

There are four ways to update primary data in a replication system based on Replication Server:

- Primary data maintenance is centralized at the primary site. Clients cannot update primary data from remote sites.
Methods for updating primary data

- Clients at remote sites update primary data through network connections.
- Clients at remote sites update primary data using request functions.
- Primary data maintenance is distributed at multiple primaries. Any resulting conflicts must be avoided or resolved.

Centralized primary maintenance

This method is the simplest, and the most restrictive, for remote clients. Client applications at remote sites use replicate data for reference only. This architecture can be used to create a copy of a production OLTP system that allows decision support applications to run separately from the OLTP system. See “Decision-support applications” on page 22 for information on applications that use this method for updating primary data.

Primary maintenance via network connections

For some applications, clients at remote sites must update primary data. The easiest way to do this is for the client to connect directly to the primary data server through the network. Replication Server distributes updates from the primary to the remote sites in the usual way.

Figure 2-5 illustrates this design.
This architecture uses client connections through the WAN. It is useful for applications with large amounts of data and low update rates. Updates are executed directly against the primary data so that they can be distributed to all sites that have subscriptions for the data. At remote sites, local replicated copies of the data can be used for decision-support applications to reduce the network load.

**Primary maintenance via request functions**

A replicated stored procedure executed in a replicate database is called a **request function**. It is called a request function because it requests a change in the primary database. You can use request functions to modify primary data without connecting to the primary data server through the WAN.

When you execute a user stored procedure in a replicate database, Replication Server sends the procedure to the primary database where you created the function replication definition.
When you execute the stored procedure, Adaptive Server records the procedure execution in the transaction log so that RepAgent can pass the request to the replication system. A request function is not required to perform any work in the local database, although it may. The real work is done when the primary Replication Server receives the procedure and executes it in the primary database.

In Figure 2-6, the gray arrows show the flow of the request function delivery. The black arrows show the flow of replicated data from the primary database to the replicate database.

**Figure 2-6: Modifying primary data with request function delivery**

The Replication Server at the replicate site stores request functions in a queue and delivers them to the primary Replication Server as soon as possible. If the primary is not available, the remote Replication Server stores the procedure until the problem is corrected.
Managing update conflicts for multiple primaries

When there are multiple primaries, remote updates should be designed so that they do not introduce errors resulting from simultaneous requests to update the same information from multiple remote sites.

If updates from two different sites conflict when they are applied at a primary site, one of the updates will be rejected, and the replicated data at the site where the update was rejected will be inconsistent with primary data.

To handle inter-site concurrency conflicts when there are multiple primaries, follow these guidelines:

- **When each row has an owner** – design the application so that inter-site conflicts are impossible. For example, restrict the updates performed at one site to rows that cannot be updated by clients at other sites. This guarantees that updates will not conflict at the primary site.

- **When there is no segmentation of ownership** – add version control information to function strings to allow conflicts to be detected and handled.

Designing conflicts out of an application

One way to handle conflicting updates from different sites is to construct the application and the environment so that conflicts cannot happen. For example, an application with customer account information distributed to each of several branch offices could require that an account be updated only at the customer’s home branch. This prevents two clients from updating the same account at the same time in different databases.

Another technique is to include a location key, such as a branch ID, in the primary key for replicated tables. If each site uses a unique location key for all of its transactions, inter-site conflicts cannot occur.

Version-controlled updates

You can use version-control updates to:

- Detect and resolve conflicting updates
- Resolve multiple primary conflicts when there is no single primary source
- Make multiple requests to a single primary
Updates are accepted or rejected based on a version column that changes each time the row is updated. The version column can be a number that increases with each update, a timestamp, or some other value from a set of unique values.

To update a row, an application must provide the current value of the version column at the primary site. Typically, the value is provided as a parameter to a replicated stored procedure. The stored procedure at the primary site checks the version parameter and takes appropriate action if it detects a conflict. If the application chooses to roll back the transaction, it is written to the exceptions log.

**Note** Managing update conflicts using version control requires careful planning and design. When there are multiple primaries, it is usually simpler and more effective to establish owners for each row of a table.
CHAPTER 3

Implementation Strategies

This chapter describes several models and strategies you can use to implement your Replication Server application design. It includes procedures and sample scripts that you can adapt to your own application.

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</tr>
</tbody>
</table>

Overview of models and strategies

The models discussed in this chapter are:

- Basic primary copy model – centralized primary data, distributed replicate data
- Distributed primary fragments model – both primary and replicate data distributed through the replication system
- Corporate rollup – distributed primary data, centralized replicate data
- Redistributed corporate rollup – same as corporate rollup, but updates redistributed to replicate databases
- Warm standby applications – two databases, one serving as backup for the other, which together as a logical unit may participate in replication

In addition, this chapter describes model variations and other strategies you can use:

- Multiple replication definitions
- Publications
Basic primary copy model

- Request functions
- Pending tables
- Master/detail relationships

These methods are described in detail in “Model variations and strategies” on page 59.

The type of application you are building, the way you update primary data, and the way you manage potential update conflicts determine the model you use to implement your replication application.

For instance, you can use the basic primary copy model to implement either a decision-support application or a low-volume distributed OLTP system. You might implement a decision-support application using either the basic primary copy model or the redistributed corporate rollup model, depending on whether your primary data is centralized or fragmented. A distributed OLTP application might be implemented using the distributed primary fragment model, with or without corporate rollup, depending on additional decision-support needs.

Basic primary copy model

The basic primary copy model allows you to replicate data from a primary database to destination databases. This model is well suited to decision-support applications, although low-volume transaction-processing applications can update primary data remotely, either directly over the WAN or through request functions (replicated stored procedures). Primary data that is updated from remote sites can then be replicated back to subscribing sites.

You can implement the basic primary copy model by using any or all of the following:

- Table replication definitions
- Applied functions
- Request functions

This section provides basic examples for using table replication definitions and applied functions. For examples of request functions and other, more advanced uses of the primary copy model, see “Model variations and strategies” on page 59.
CHAPTER 3  Implementation Strategies

Using table replication definitions

Using table replication definitions allows you to replicate data from a primary source as read-only copies.

You can create one or many replication definitions for a primary table although a particular replicate table can subscribe to only one of them. See “Multiple replication definitions” on page 60 for an example using multiple replication definitions.

You also can collect replication definitions in a publication and subscribe to all of them at one time with a publication subscription. See “Publications” on page 62 for an example using publications.

For each table you want to replicate according to the basic primary copy model, you need to:

- Set up routes and connections between Replication Servers.
- Create the table you want to replicate in the primary database.
- Create the table (or tables) to which you want to replicate in destination databases.
- Create indexes and grant appropriate permissions on the tables.

At the primary site:
- Mark the primary table for replication using the `sp_setreptable` system procedure.
- Create one (or more) replication definitions for the table at the primary Replication Server.

At the replicate sites:
- Create subscriptions for the table replication definitions at each replicate Replication Server.

See the `Replication Server Administration Guide` for details on setting up the basic primary copy model.

In Figure 3-1, a client application at the primary (Tokyo) site makes changes to the publishers table in the primary database. At the replicate (Sydney) site, the publishers table subscribes to the primary publishers table—for those rows where pub_id is equal to or greater than 1000.
Marking the table for replication

This script marks the publishers table for replication.

```sql
-- Execute this script at Tokyo data server
-- Marks publishers for replication
sp_setreptable publishers, 'true'
go
/* end of script */
```

Replication definition

This script creates a table replication definition for the publishers table at the primary Replication Server.

```sql
-- Execute this script at Tokyo replication server
-- Creates replication definition pubs_rep
create replication definition pubs_rep
with primary at TOKYO_DS.pubs2
```
with all tables named 'publishers'  
(pub_id char(4),  
pub_name varchar(40)  
city varchar(20)  
state varchar(2)  
primary key (pub_id)  
go  
/* end of script */

**Subscription**

This script creates a subscription for the replication definition defined at the primary Replication Server.

```sql
-- Execute this script at Sydney replication server  
-- Creates subscription pubs_sub  
Create subscription pubs_sub  
for pubs_rep  
with replicate at SYDNEY_DS.pubs2  
where pub_id >= 1000  
go  
/* end of script */
```

**Using applied functions**

You can also use applied functions to replicate stored procedure invocations to remote sites with replicate data. Using applied functions to replicate primary data can:

- Reduce network traffic over the WAN
- Increase throughput and decrease latency because applied functions execute more rapidly
- Enable a more modular system design

In the following example, a client application at the primary (Tokyo) site executes a user stored procedure, `upd_publishers_pubs2`, which makes changes to the `publishers` table in the primary database. Execution of `upd_publishers_pubs2` invokes function replication, which causes the corresponding stored procedure, also named `upd_publishers_pubs2`, to execute on the replicate data server.

To create an applied function for an application that implements the basic primary copy model, you need to:
Basic primary copy model

At the primary site:

- Create the user stored procedure in the primary database.
- Mark the user stored procedure for replicated function delivery using `sp_setrepproc`.
- Grant the appropriate procedure permissions to the appropriate user.
- At the primary Replication Server, create the function replication definition for the stored procedure with name, parameters, and datatypes that match those of the stored procedure. You can specify only the parameters you want to replicate.

At the replicate site:

- Create a stored procedure in the replicate database with the same parameters (or a subset of those parameters) and datatypes as those created in the primary database. Grant appropriate permissions to the procedure to the maintenance user.
- Create a subscription to the function replication definition in the replicate Replication Server.

Figure 3-2: Basic primary copy model using applied functions
Stored procedures

This script creates stored procedures for the publishers table at the primary and replicate sites.

```sql
-- Execute this script at Tokyo and Sydney data servers
-- Creates stored procedure upd_publishers_pubs2
create procedure upd_publishers_pubs2
(@pub_id char(4),
 @pub_name varchar(40),
 @city varchar(20),
 @state char(2))
as
update publishers
set
    pub_name = @pub_name,
    city = @city,
    state = @state
where
    pub_id = @pub_id
go
/* end of script */
```

Function replication definition

This script creates a function replication definition for the publishers table at the primary Replication Server. The replication definition uses the same name, parameters, and datatypes as the stored procedure in the primary database.

```sql
-- Execute this script at Tokyo replication server
-- Creates replication definition
upd_publishers_pubs2
create function replication definition
upd_publishers_pubs2
with primary at TOKYO_DS.pubs2
(@pub_id char(4),
 @pub_name varchar(40),
 @city varchar(20),
 @state char(2))
go
/* end of script */
```

Subscriptions

You can create a subscription for a function replication definition in one of two ways:
Basic primary copy model

- Use the `create subscription` command and the no-materialization method. Use this method if primary data is already loaded at the replicate, and updates are not in progress.

- Use the `define subscription`, `activate subscription`, and `validate subscription` commands and the bulk materialization method. Use this method if you are coordinating loading data with updates.

Examples of both methods follow.

**Using the no-materialization method**

This script creates a subscription at the replicate Replication Server using the no-materialization method for the replication definition defined at the primary Replication Server.

```sql
-- Execute this script at Sydney replication server
-- Creates subscription using no-materialization for upd_publishers_pubs2
create subscription upd_publishers_pubs2_sub for upd_publishers_pubs2
with replicate at SYDNEY_DS.pubs2
without materialization
go
/* end of script */
```

**Using bulk materialization**

This script defines, activates, and validates a subscription at the replicate Replication Server for the replication definition defined at the primary Replication Server.

```sql
-- Execute this script at Sydney replication server
-- Creates subscription using bulk materialization for upd_publishers_pubs2
define subscription upd_publishers_pubs2_sub for upd_publishers_pubs2
with replicate at SYDNEY_DS.pubs2
go

activate subscription upd_publishers_pubs2_sub for upd_publishers_pubs2
with replicate at SYDNEY_DS.pubs2
go
/* Load data. If updates are in progress, use activate subscription with the "with suspension" clause and
```
Distributed primary fragments model

In this model, tables at each site contain both primary and replicated data. However, each site functions as a primary site for a particular subset of rows called a *fragment*. Updates to the primary fragment are distributed to other sites. Updates to nonprimary data are received from the primary sites of other fragments.

Applications that use the distributed primary fragments model have distributed tables that contain primary and replicated data. The Replication Server at each site distributes modifications made to local primary data to other sites and applies modifications received from other sites to the data that is replicated locally.

The following tasks must be performed at each site to replicate a table in the distributed primary fragments model:

- Create the table in each database. The table should have the same structure in each database.
- Create indexes and grant appropriate permissions on the tables.
- Allow for replication on the tables using the `sp_setreptable` system procedure.
- Create a replication definition for the table at each site.
- At each site, create a subscription for the replication definition at the other sites. If $n$ is the number of sites, create $n-1$ subscriptions.

Figure 3-3 diagrams the flow of data for distributed primary fragments:
Figure 3-3 illustrates a salesdetail table set up with distributed primary fragments at three sites. Each site receives replicated data via two subscriptions.
Replication definitions

These scripts create replication definitions for the salesdetail table at each site:

```
-- Execute this script at Chicago RSCHI.
-- Creates replication definition chi_sales.
create replication definition chi_sales_rep
  with primary at DSCHI.DBCHI
  with all tables named 'salesdetail'
    (stor_id char(4),
     ord_num varchar(20),
     title_id varchar6),
  qty smallint,
  discount float)
primary key (stor_id, ord_num)
searchable columns
  (stor_id, ord_num, title_id)
```
Distributed primary fragments model

/* end of script */
-- Execute this script at New York RSNY.
-- Creates replication definition ny_sales.
create replication definition ny_sales_rep
with primary at DSNY.DBNY
with all tables named 'salesdetail'
(stor_id char(4),
 ord_num varchar(20),
 title_id varchar6),
 qty smallint,
 discount float)
primary key (stor_id, ord_num)
searchable columns
 (stor_id, ord_num, title_id)
go
/* end of script */
-- Execute this script at San Francisco RSSF.
-- Creates replication definition sf_sales.
create replication definition sf_sales_rep
with primary at DSSF.DBSF
with all tables named 'salesdetail'
(stor_id char(4),
 ord_num varchar(20),
 title_id varchar6),
 qty smallint,
 discount float)
primary key (stor_id, ord_num)
searchable columns
 (stor_id, ord_num, title_id)
go
/* end of script */

Subscriptions

Each site has a subscription to the replication definitions at the other two sites. These scripts create the subscriptions:

-- Execute this script at Chicago RSCHI.
-- Creates subscriptions to ny_sales and sf_sales.
create subscription ny_sales_sub
 for ny_sales_rep
 with replicate at DSCHI.D BCHI
 where stor_id = '5023'
go
create subscription sf_sales_sub
for sf_sales_rep
    with replicate at DSCHI.DBCHI
    where stor_id = ’8042’
go
/* end of script */

-- Execute this script at New York RSNY.
-- Create subscriptions to chi_sales and sf_sales.
create subscription chi_sales_sub
    for chi_sales_sub
    with replicate at DSNY.DBNY
    where stor_id = ’7067’
go
create subscription sf_sales_sub
    for sf_sales_rep
    with replicate at DSNY.DBNY
    where stor_id = ’8042’
go
/* end of script */

-- Execute this script at San Francisco RSSF.
-- Creates subscriptions to chi_sales and ny_sales.
create subscription chi_sales_sub
    for chi_sales_rep
    with replicate at DSSF.DBSF
    where stor_id = ’7067’
go
create subscription ny_sales_sub
    for ny_sales_rep
    with replicate at DSSF.DBSF
    where stor_id = ’5023’
go
/* end of script */

**Corporate rollup**

In this model, multiple primary fragments maintained at remote sites are consolidated into a single aggregate replicate table at a central site.

The corporate rollup model has distributed primary fragments and a single, centralized consolidated replicate table. The table at each primary site contains only the data that is primary at that site. No data is replicated to these sites. The corporate rollup table is a “rollup” of the data at the primary sites.
The corporate rollup model requires distinct replication definitions at each of the primary sites. The site where the data is consolidated has a subscription for the replication definition at each primary site.

Replication Agents are required at the primary sites but not at the central site, since data will not be replicated from that site.

These tasks must be performed to create a corporate rollup from distributed primary fragments:

- Create the table in each primary database and in the database at the central site. The tables should have the same structure and the same name.
- Create indexes and grant appropriate permissions on the tables.
- In each remote database, allow for replication on the table with the `sp_setreptable` system procedure.
- Create a replication definition for the table at each remote site.
- At the headquarters site, where the data is to be consolidated, create subscriptions for the replication definitions at the remote sites.

Figure 3-5 illustrates the flow of data for a corporate rollup application model.
Figure 3-5: Corporate rollup model with distributed primary fragments

Figure 3-6 illustrates a salesdetail table with a corporate rollup at a headquarters site. The headquarters site receives data from the remote sites via three subscriptions.
Replication definitions

These scripts create replication definitions for the salesdetail table at each primary site:

```
-- Execute this script at Chicago RSCHI.
-- Creates replication definition chi_sales.
cREATE DEFINITION chi_sales_rep
with primary at DSCHI.DBCHI
with all tables named 'salesdetail'
(stor_id char(4),
 ord_num varchar(20),
 title_id varchar6),
 qty smallint,
 discount float)
primary key (stor_id, ord_num)
searchable columns
(stor_id, ord_num, title_id)
go
/* end of script */
-- Execute this script at New York RSNY.
-- Creates replication definition ny_sales.
cREATE DEFINITION ny_sales_rep
with primary at NYCHI.DBCHI
with all tables named 'salesdetail'
(stor_id char(4),
 ord_num varchar(20),
 title_id varchar6),
 qty smallint,
 discount float)
primary key (stor_id, ord_num)
searchable columns
(stor_id, ord_num, title_id)
go
/* end of script */
```
with primary at DSNY.DBNY
with all tables named 'salesdetail'
    (stor_id char(4),
    ord_num varchar(20),
    title_id varchar6),
    qty smallint,
    discount float)
primary key (stor_id, ord_num)
searchable columns
    (stor_id, ord_num, title_id)
go
/* end of script */
-- Execute this script at San Francisco RSSF.
-- Creates replication definition sf_sales.
cREATE REPLICATION DEFINITION sf_sales_rep
with primary at DSSF.DBSF
with all tables named 'salesdetail'
    (stor_id char(4),
    ord_num varchar(20),
    title_id varchar6),
    qty smallint,
    discount float)
primary key (stor_id, ord_num)
searchable columns
    (stor_id, ord_num, title_id)
go
/* end of script */

Subscriptions

The headquarters site has subscriptions to the replication definitions at each of the three primary sites. The primary sites have no subscriptions. This script creates the subscriptions in the RSHQ Replication Server:

-- Execute this script at Headquarters RSHQ.
-- Creates subscriptions to chi_sales, ny_sales, and sf_sales.
CREATE SUBSCRIPTION chi_sales_sub
for chi_sales_rep
    with replicate at DSHQ.DBHQ
    where stor_id = '7067'
go
CREATE SUBSCRIPTION ny_sales_sub
for ny_sales_rep
    with replicate at DSHQ.DBHQ
Redistributed corporate rollup

The redistributed corporate rollup is the same as the corporate rollup model, except that the consolidated table is redistributed back to the remote sites.

Primary fragments distributed at remote sites are rolled up into a consolidated table at a central site. At the site where the fragments are consolidated, RepAgent processes the consolidated table as if it were primary data.

The consolidated table is described with a replication definition. Other sites can create subscriptions for this table.

Normally, RepAgent for Adaptive Server filters out updates made by the maintenance user. This ensures that replicated data is not redistributed as primary data.

The RepAgent send_maint_xacts_to_replicate option is provided for the redistributed corporate rollup model. If you start RepAgent with send_maint_xacts_to_replicate set to “true,” RepAgent submits all updates to the Replication Server as if they were made by a client application. (If you are using the LTM, start the program with the -A option.)

If you use the redistributed corporate rollup model:

- Do not allow primary sites to resubscribe to their primary data. If they do, transactions could loop endlessly through the system.
- Do not allow applications to update the corporate rollup table. All updates should originate from the primary sites.

Figure 3-7 illustrates the flow of data in an application based on the Redistributed Corporate Rollup model.
The design of the redistributed corporate rollup model is identical to the corporate rollup model, except that:

- RepAgent must be installed at the headquarters site for the DBHQ database. RepAgent must be started with the `send_maint_xacts_to_replicate` option set so that it will transfer log records from the maintenance user.

  If you are using LTM, it must be started with the `-A` option.

- RepAgent is required for the RSHQ RSSD, since data will be distributed from that site.
Warm standby applications

- A replication definition must be created for the salesdetail table at the headquarters site. Other sites can create subscriptions to this replication definition, but the primary sites must not subscribe to their own primary data.

- The RSHQ Replication Server must have routes to the other sites that create subscriptions for the consolidated replicate table. If the primary sites create subscriptions, routes must be created to them from RSHQ.

Warm standby applications

In a warm standby application, Replication Server maintains a pair of Adaptive Servers (or SQL Servers), one of which acts as the backup of the other.

Typically, client applications update the active database while Replication Server maintains the other database as a standby copy of the active database. If the active database fails, or if you need to perform maintenance on the active data server or database, you can switch to the standby database (and back) with little interruption of client applications.

In a warm standby application, you create three connections:

- A logical connection that Replication Server maps to the currently active database
- A physical connection for the active database
- A physical connection for the standby database

The logical database in a warm standby application may, with respect to other databases in the replication system, function as one of the following:

- A database that does not participate in replication
- A primary database
- A replicate database

The procedure in this section demonstrates how to set up a warm standby system for a database that acts as a primary database in a replication system.

Figure 3-8 illustrates a warm standby application operating on the BOSTON_DS data server for a pubs2 database on the NY_DS data server. The database is replicated to TOKYO_DS.
In this scenario the pubs2 database acts as a primary database in a replication environment. The primary pubs2 database for which a standby is created is called the active database.

Setting up a warm standby application

The following procedure is used to set up a warm standby application for an active database. In this procedure, an active database is already established. The procedure will be somewhat different if the active database has not yet been created. Make sure you review the warm standby information in the Replication Server Administration Guide before proceeding.

Note You must use Adaptive Server or SQL Server databases for this procedure.

In the active database:

1. Mark the entire active database for replication to the standby database with the sp_reptostandby stored procedure.
**Warm standby applications**

sp_reptostandby enables replication of data manipulation language (DML) and supported data definition language (DDL) commands and stored procedures. Refer to Chapter 13, “Managing Warm Standby Applications,” in the *Replication Server Administration Guide* for detailed information.

2. Reconfigure RepAgent using the `sp_config_rep_agent` stored procedure with the `send_warm_standby_xacts` option. Restart RepAgent.

3. Grant `replication_role` to the active database maintenance user.

4. On the active data server, add the maintenance user of the standby database to the active database, and grant `replication_role` to the new maintenance user. This step ensures that the maintenance user ID exists in the standby database after the database is loaded (step 8).

5. Log in to the Replication Server that is to manage the warm standby database, and create a logical connection for the active database, using the `create logical connection` command. The name of the logical connection must be the same as the name of the active database.

   **Note** If you create the logical connection *before* you create the active database connection, use different names for the logical connection and the active database.

6. On the standby data server, create the standby database with the same size as the active database.

7. Use Sybase Central or `rs_init` to create the standby database connection. For more information, see the Replication Server online help and the Replication Server installation and configuration guide for your platform.

   After the connection is created, log in to Replication Server and use the `admin logical_status` command to make sure that the new connection is “active.”

8. Initialize the standby database using `dump` and `load` without the `rs_init “dump marker”` option. (Or you can use `bcp`. Refer to the *Replication Server Administration Guide* for more information.)

   a. On the Replication Server, suspend the active database connection.

      **Note** If you cannot suspend the active database, use `dump` and `load` with the `rs_init “dump marker”` option.

   b. On the active Adaptive Server, dump the active database.
c Load the active database dump into the standby database.

d On the standby Adaptive Server, put the standby database online.

9 On the Replication Server, resume connections to the active and standby
databases, using the resume connection command.

Check the logical status, using the admin logical_status command. Do not
continue unless both active and standby databases are marked “active.”

10 Verify that modifications occur from active to standby database.

Using isql, update a record in the active database and then verify the update
in the standby database.

Switching to the standby database

If it becomes necessary to switch from the active database to the standby
database, you need to take steps to prevent client applications from executing
transactions against or updating the active database. After the switch is
complete, clients can connect to the new active database to continue their work.
See “Switching clients to the new database” on page 58 for details.

Before switching to a standby database, you should determine whether a switch
is necessary:

• Don’t switch if the active data server is experiencing a transient failure. A
  transient failure is a failure from which the Adaptive Server recovers when
  restarted, without additional recovery steps.

• Do switch if the active database will be unavailable for a long period of
time.

You must use the switch active command to switch the active and standby
databases. The following procedure illustrates how to switch the warm standby
system illustrated in Figure 3-8 on page 55 from the active database to the
standby database.

1 On the Replication Server, use switch active to switch processing to the
standby database.

2 Monitor progress of the switch. The switch is complete when the standby
connection is active and the previously active connection is suspended.

a On the Replication Server, check the logical status, using the
admin_logical_status command.
Warm standby applications

b. To follow the progress of the switch, check the last several entries in the Replication Server error log.

3. Start the RepAgent for the new active database.

4. Decide what you want to do with the old active database. You can:
   - Bring the database online as the new standby database, and resume connections so that Replication Server can apply new transactions, or
   - Drop the database connection using the `drop connection` command. You can add it again later as the new standby database.

5. Using `isql`, update a record in the new active database, and then check the update on the new standby database.

Figure 3-9: Warm standby system, after switching

Switching clients to the new database

Switching from the active to the standby database does not switch client applications to the new active data server and database. You must devise a method to handle client switching. For example, you could:

- Set up two interfaces files, one for client applications and one for Replication Server. At switch time, modify the client interfaces file to point to the new active server.
• Create an interfaces file entry with a symbolic data server name for use by client applications. At switch time, modify the address information associated with the symbolic name.

• Use a mechanism, such as an intermediate Open Server, to map the client application data server connections to the currently active data server automatically.

See the Replication Server Administration Guide for more information.

Model variations and strategies

This section describes some model variations and other strategies you can use to implement your replication system design. They are:

• Multiple replication definitions – a strategy using multiple replication definitions for a primary table to specify different table names, column sets, and column names, thereby presenting differing views of the primary table to subscribing replicate tables

• Publications – a strategy that allows users to subscribe to a set of replication definitions with a single subscription

• Request functions – a variation of the basic primary copy model that allows remote sites to make changes to primary data without accessing the primary database directly

• Pending tables – a strategy used with request functions that allows users to see the results of updates to primary data before the update has been returned to the replicate site

• Implementing master/detail relationships – uses request functions and stored procedures to ensure proper subscription migration
Multiple replication definitions

You can create multiple replication definitions for a single primary table. Each replication definition can specify different table names, column sets, and column names, thereby presenting different views of the primary table to each of the subscribing replicate tables.

**Note** You can create multiple replication definitions for a primary table, and a replicate table can subscribe to multiple table replication definitions. However, a replicate table can subscribe only to one replication definition per primary table.

To set up a system using multiple replication definitions, follow the directions in “Using table replication definitions” on page 37, creating multiple replication definitions as needed and a subscription for each one. When you use multiple replication definitions, you are using a variation of the primary copy model.

In Figure 3-10, a client application at the primary (Tokyo) site makes changes to the publishers table in the primary database. At one replicate (Sydney) site, the publishers table subscribes to the complete table, and the pubs_copy table subscribes only to the pub_id and pub_name columns. At another replicate (Seattle) site, the pubs_copy table subscribes to the pub_id and pub_name columns, where pub_id is equal to or greater than 1000.
Replication definitions

These scripts create table replication definitions for the publishers table at the primary Replication Server. Each replication definition describes a different view of the primary table.

```
-- Execute this script at Tokyo replication server
-- Creates replication definitions pubs_rep and pubs_copy_rep
create replication definition pubs_rep
with primary at TOKYO_DS.pubs2
with all tables named 'publishers'
(pub_id char(4),
pub_name varchar(40)
city varchar(20)
state varchar(2)
primary key (pub_id)
goto
```
Model variations and strategies

create replication definition pubs_copy_rep
with primary at TOKYO_DS.pubs2
with primary table named ‘publishers’
with replicate table named ‘pubs_copy’
(pub_id char(4),
pub_name varchar(40)
primary key (pub_id)
gogo
/* end of script */

Subscriptions

These scripts create subscriptions for the replication definitions defined at the primary Replication Server.

-- Execute this script at Sydney replication server
-- Creates subscription pubs_sub and pubs_copy_rep
Create subscription pubs_sub
for pubs_rep
with replicate at SYDNEY_DS.pubs2
go
create subscription pubs_copy_sub
for pubs_copy_rep
with replicate at SYDNEY_DS.pubs2
go
/* end of script */

-- Execute this script at Seattle replication server
-- Creates subscription pubs_copy_sub
create subscription pubs_copy_sub
for pubs_copy_rep
with replicate at SEATTLE_DS.pubs2
where pub_id >= 1000
go
/* end of script */

Publications

Use publications to collect replication definitions for tables and/or stored procedures and then subscribe to them as a group.

With publications, you can monitor the status of one publication subscription for a set of tables and procedures. Publication usage is not a separate model; it provides a grouping technique that can be used by any model.
When you use publications, you create and manage the following objects:

- **Articles** – replication definition extensions for tables or stored procedures that let you put table or function replication definitions in a publication.
- **Publications** – groups of articles from the same primary database.
- **Publication subscriptions** – subscriptions to a publication. When you create a publication subscription, Replication Server creates a subscription for each of the publication’s articles.

The following steps summarize the procedure for replicating data using publications.

**At the primary site:**

1. Create or select the replication definitions to include in the publication.
2. Create the publication, using the `create publication` command.
3. Create articles that reference the replication definitions you have chosen, using the `create article` command.
4. Validate the publication using the `validate publication` command.

**At the replicate site:**

1. Create a subscription for the publication using the `create subscription` command.

In Figure 3-11, a table replication definition `pubs_rep`, referenced by two articles, and a function replication definition, referenced by one article, are collected in the publication `pubs2_pub`. 
Figure 3-11: Publications

Stored procedure

This script creates a stored procedure update_authors_pubs2 that updates the authors table in the pubs2 database. Create the same procedure at the primary and replicate sites.

```sql
-- Execute this script at the Tokyo and Sydney data servers
-- Creates the stored procedure
create procedure upd_authors_pubs2
(@au_id          id,
 au_lname           varchar(40),
 au_fname          varchar(20),
 phone                char(12),
 address           varchar((12),
 city                   varchar(20),
 state             char(2),
country              varchar(12),
```
Function replication definition

This script creates a function replication definition at the primary site.

```
-- Execute this script at the Tokyo replication server
-- Creates the function replication definition
create function replication definition upd_authors_pubs2
with primary at TOKYO_DS.pubs2
(id, au_lname VARCHAR(40), au_fname VARCHAR(20), phone CHAR(12),
    address VARCHAR(12), city VARCHAR(20), state CHAR(2),
    country VARCHAR(12), postalcode CHAR(10))
go
/* end of script */
```

Table replication definition

This script creates a table replication definition for the publishers table at the primary Replication Server.

```
-- Execute this script at Tokyo replication server
-- Creates replication definitions pubs_rep
```

postalcode char(10)
as
update authors
set
  au_lname = @varchar(40),
  au_fname = @varchar(20),
  phone = @char(12),
  address = @varchar((12),
  city = @varchar(20),
  state = @char(2),
  country = @varchar(12),
  postalcode = @char(10)
where au_id = @au_id
go
/* end of script */
Model variations and strategies

create replication definition pubs_rep
with primary at TOKYO_DS.pubs2
with all tables named ‘publishers’
(pub_id char(4),
pub_name varchar(40)
city varchar(20)
state varchar(2)
primary key (pub_id)
gogo
/* end of script */

Publication

This script creates the publication pubs2_pub at the primary Replication Server.

-- Execute this script at Tokyo replication server
-- Creates publication pubs_pub
create publication pubs2_pub
with primary at TOKYO_DS.pubs2
gogo
/* end of script */

Articles

This script creates articles for the publication pubs2_pub at the primary Replication Server. It creates two articles for the replication definition pubs_rep.

-- Execute this script at Tokyo replication server
-- Creates articles upd_authors_art,
-- pubs_art, and pubs_copy_art
create article upd_authors_art
for pubs2_pub
with primary at TOKYO_DS.pubs2
with replication definition upd_authors_rep
gogo

create article pubs_art
for pubs2_pub
with primary at TOKYO_DS.pubs2
with replication definition pubs_rep
go

create article pubs_copy_art
for pubs2_pub
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with primary at TOKYO_DS.pubs2
with replication definition pubs_rep
where pub_id >= 1000
go
/* end of script */

Validation

This script changes the status of the publication pubs2_pub to “valid.”

-- Execute this script at Tokyo replication server
-- Validates the publication pubs2_pub
validate publication pubs2_pub
with primary at TOKYO_DS.pubs2
go
/* end of script */

Subscription

This script creates the subscription pubs2_pub_sub for the publication pubs2_pub. When this script is run, Replication Server creates article subscriptions for upd_authors_art, pubs_art, and pubs_copy_art.

-- Execute this script at Sydney replication server
-- Creates publication subscription pubs2_pub_sub
create subscription pubs2_pub_sub
for publication pubs2_pub
with primary at TOKYO_DS.pubs2
with replicate at SF.pubs2
without materialization
go
/* end of script */

Request functions

You can use request functions to invoke stored procedures on primary data. Request functions allow clients at remote sites to update primary data without accessing the primary database directly. Replication Server invokes a stored procedure on the primary database, which carries out the requested transactions. This method preserves the basic primary copy model while letting you distribute transactions. You do not have to fragment primary data among replicate sites.

The following two examples illustrate:
Model variations and strategies

- A system that uses request functions
- A system that uses request functions, applied functions, and a local pending table

A basic example using request functions

In this example, a client application at the replicate (Sydney) site executes the stored procedure `upd_publishers_pubs2_req`, which makes no changes to the replicate database but causes an associated stored procedure, `upd_publishers_pubs2`, to execute and change data on the primary (Tokyo) site. `upd_publishers_pubs2_req` is an empty-body stored procedure.

Perform these tasks to create a request function. No subscriptions are necessary at the primary database; function executions at a site other than the primary for the replication definition automatically flow to the primary site.

At the replicate site:
- Create an empty-body stored procedure at the replicate data server, and
  - Mark the procedure for replicated function delivery using `sp_setrepproc`.
  - Grant appropriate permissions to the procedure to the appropriate user.

The empty-body procedure either does nothing or displays a message indicating a pending update. It may have a different name than the associated stored procedure in the primary database.

At the primary site:
- Create a stored procedure that updates primary data. Grant permissions to the appropriate user.
- Create a function replication definition at the primary Replication Server for the empty-body stored procedure in the replicate database. The function replication definition and the empty-body stored procedure have the same name.

Once the primary data has been updated via the replication of the request stored procedure, data can be delivered to replicate sites (including the requestor) via either table replication or applied function replication.

Note  It is not necessary to create a subscription for a request function.
Figure 3-12: Request functions

Stored procedures
At the primary site: The script on creates the user stored procedure `upd_publisher_pubs2` at the primary site.

At the replicate site: This script creates the empty-body stored procedure `upd_publishers_pubs2_req`.

```sql
-- Execute this script at Sydney replication server
-- Creates stored procedure
create procedure upd_publishers_pubs2_req
(@pub_id char(4),
@pub_name  varchar(40),
@city  varchar(20),
@state  char(2))
as
begin
if (select 1) > 1
    print "Submitting request."
end
```
Model variations and strategies

Function replication definition

This script creates the function replication definition at the primary Replication Server. It has the same name as the empty-body stored procedure. The deliver as clause tells the primary data server to execute `upd_publisher_pubs2` at the primary site.

```sql
-- Execute this script at Tokyo replication server
-- Creates replication definition
create function replication definition upd_publishers_pubs2_req
with primary at TOKYO_DS.pubs2
deliver as 'upd_publisher_pubs2'
(@pub_id char(4),
@pub_name varchar(40),
@city varchar(20),
@state char(2))
go
/* end of script */
```

An example using a local pending table

The pending table is a design enhancement of applied and request functions that allows clients at a remote site to update primary data and see the updates at the replicate site before they are returned from the primary site. Use this model to implement local update applications.

In this strategy, a client application at a remote site executes a user stored procedure that updates data at the primary site using a request function. Changes to the primary data are replicated to the remote site via an applied function. A local pending table lets clients at the remote site see updates that are pending at the replicate site before the replication system returns the updates.

When a client application executes the user stored procedure at the replicate data server, it:

- Causes an associated stored procedure to execute and update data at the primary site
- Enters those updates in the local pending table
When the update succeeds at the primary database, it is distributed to the remote sites, including the site where the transaction originated. At the replicate site, a stored procedure updates the replicated table and deletes the corresponding updates from the pending table.

To use applied functions, request functions, and a local pending table, you must complete these tasks.

**At the replicate site:**
- Create a pending table in the replicate database. Grant appropriate permissions.
- Create a user stored procedure in the replicate database that initiates the request function and inserts data updates into the pending table.
- Mark the user stored procedure for replicated function delivery using `sp_setrepproc`.
- Grant procedure permissions to the appropriate user.
- Create a user stored procedure in the replicate database that updates the replicated table and deletes the corresponding update from the pending table. Grant appropriate permissions to the maintenance user.
- Create a subscription to the function replication definition.

**At the primary site:**
- Create the stored procedure that modifies the primary data.
- Create the function replication definition for the applied function.
- Create the function replication definition for the request function.

In this example, a client application at the replicate (Sydney) site executes a stored procedure `upd_publishers_pubs2_req`, which inserts values in the `publishers_pend` table and causes an associated stored procedure, `upd_publishers_pubs2`, to execute at the primary (Tokyo) site. Execution of `upd_publishers_pubs2` at the primary site causes the stored procedure `upd_publishers_pubs` to execute at the replicate site, which updates the `publishers` table and deletes the corresponding information from the `publishers_pend` table.

Figure 3-13 illustrates the data flow when you use applied functions, request functions, and a local pending table. The gray arrows show the flow of the request function delivery. The black arrows show the flow of the applied function delivery.
Pending table

This script creates a pending table in the replicate database.

```sql
-- Execute this script at Sydney data server
-- Creates local pending table
create table publishers_pend
(pub_id char(4) not null,
 pub_name varchar(40) null,
 city varchar(20) null,
 state char(2) null)
go
/* end of script */
```

Stored procedures

The script creates the stored procedure `upd_publisher_pubs2` at the primary site.
The following script modifies the `upd_publishers_pub2_req` stored procedure at the replicate site. The `insert into` clause tells the replicate Replication Server to insert values into the `publishers_pend` table.

```sql
-- Execute this script at Sydney data server
-- Creates stored procedure
create procedure upd_publishers_pubs2_req
    (@pub_id    char(4),
     @pub_name    varchar(40),
     @city       varchar(20),
     @state   char(2))
as
    insert into publishers_pend
    values
    (@pub_id, @pub_name, @city, @state)
go
/* end of script */
```

This script modifies the `upd_publishers_pubs2` procedure for the replicate site. It updates the `publishers` table and deletes the corresponding information from the `publishers_pend` table.

```sql
-- Execute this script at Sydney data server
-- Creates stored procedure upd_publishers_pubs2
create procedure upd_publishers_pubs2
    (@pub_id char(4),
     @pub_name varchar(40),
     @city varchar(20),
     @state char(2))
as
    update publishers
    set
      pub_name = @pub_name,
      city = @city,
      state = @state
    where
      pub_id = @pub_id
    delete from publishers_pend
    where
      pub_id = @pub_id
go
/* end of script */
```
Model variations and strategies

Function replication definitions

The script “Function replication definition” on page 70 creates the function replication definition `upd_publishers_pubs2_req`—for the request function—at the primary site.

The script “Function replication definition” on page 41 creates the function replication definition `upd_publishers_pubs2`—for the applied function—at the primary site.

Subscription

The script “Subscriptions” on page 41 creates the subscription `upd_publishers_pubs2_sub`—for the function replication definition `upd_publisher_pubs2`—at the replicate site.

Implementing master/detail relationships

You can use applied functions to replicate only selected data to remote sites. Using applied functions in this way reduces network traffic.

To implement master/detail relationships, use applied functions to support selective subscription to the master/detail tables. In this example,

- The publishers and titles tables exist in the pubs2 database at the primary and replicate sites.
- NY_DS is the primary site data server and SF_DS is the replicate site data server.

Figure 3-14 describes the publishers (master) and titles (detail) tables at the primary and replicate sites:
The primary site contains all records, but the replicate site is interested only in records related to the state of California (CA). Only a selection of publishers and titles records need to be replicated, based on the state column. However, only the publishers table contains a state column.

Adding a state column to the titles table adds redundancy to the system. A second, more efficient solution ties updates to master and detail tables through stored procedures and then replicates the stored procedures using applied functions. The logic to maintain selective subscription is contained in the stored procedures.

For example, if, at the primary site, a publisher’s state is changed from NY to CA, a record for that publisher must be inserted at the replicate site. Having replicate rows inserted or deleted as a result of updates that cause rows in a subscription to change is called subscription migration.

To ensure proper subscription migration, subscriptions are needed for a set of “upper-level” stored procedures that control the stored procedures that actually perform the updates. It is the invocation for the upper-level stored procedure that is replicated from the primary site to the replicate site.

To handle changes in the state column, the replicate site must subscribe to updates when either the new state or the old state is CA.

The sections below list the activities you must perform to enable selective substitution at the replicate Replication Server.
At the primary and replicate sites:

- Create stored procedures that insert records into the publishers and titles tables and an upper-level stored procedure that controls the execution of both insert procedures.
- Create stored procedures that delete records from the publishers and titles tables and an upper-level stored procedure that controls the execution of both delete procedures.
- Create stored procedures that update records in the publishers and titles tables and an upper-level stored procedure that controls the execution of both update procedures.
- Grant appropriate permissions on all upper-level stored procedures.

At the primary site:

- Mark each upper-level stored procedure for replication, using `sp_setrepproc`.
- Create a function replication definition for each upper-level stored procedure.

At the replicate site:

- Create subscriptions to the function replication definitions.

Stored procedures with insert clauses

The insert procedures are identical at the primary and replicate sites. The upper-level stored procedure that controls the insert procedures and the insert procedures observe the following logic:

- A publisher record is inserted only when there is no title ID.
- A title record is inserted only when the publisher exists.

These scripts create the `ins_publishers` and `ins_titles` insert stored procedures and the upper-level stored procedure `ins_pub_title`.

```sql
-- Execute this script at NY and SF data servers
-- Creates stored procedure
create procedure ins_publishers
    (@pub_id char(4), @pub_name varchar(40)=null,
    city varchar(20)=null, @state char(2)=null)
as
    insert publishers values (@pub_id, 
    @pub_name, @city, @state)
/* end of script */
-- Execute this script at NY and SF data servers
-- Creates stored procedure
create procedure ins_titles
    (@title_id tid, @title varchar(80), @type char(12),
Stored procedures with delete clauses

The delete procedures are identical at the primary and replicate sites. The upper-level stored procedure that controls the delete procedures and the delete procedures observe the following logic:

- When a record is deleted, all dependent child records are also deleted.
- A publisher record is not deleted when a title record exists.

These scripts create the `del_publishers` and `del_titles` stored procedures and the upper-level stored procedure `del_pub_title`.

```sql
-- Execute this script at NY and SF data servers
-- Creates stored procedure
create procedure del_publishers
(@pub_id char(4))
```
as
begin
if exists (select * from titles where pub_id=@pub_id)
   raiserror 20005 "**FATAL ERROR: Existing titles**"
else
   delete from publishers where pub_id=@pub_id
end

/* end of script */

-- Execute this script at NY and SF data servers
-- Creates stored procedure
create procedure del_titles
(@title_id tid, @pub_id char(4)=null)
as
if @pub_id=null
   delete from titles where title_id=@title_id
else
   delete from titles where pub_id=@pub_id
end

/* end of script */

-- Execute this script at NY and SF data servers
-- Creates stored procedure
create procedure del_pub_title
(@pub_id char(4), @state char(2), @title_id tid=null)
as
begin
if @title_id != null
   begin
      exec delTitles @title_id
      return
   end
if @pub_id != null
   begin
      exec delTitles @title_id, @pub_id
      exec del_publishers @pub_id
   end
end

/* end of script */

---

**Stored procedures with update clauses**

The update procedures differ at the primary and replicate sites.

At the primary site:  
Update procedures observe the following logic:

- Raise an error on an unknown `pub_id`.  

---

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• If a title does not exist, insert one.

These scripts create the `upd_publishers` and `upd_titles` stored procedures and the upper-level stored procedure `upd_pub_title` that controls the execution of `upd_publishers` and `upd_titles`.

Note that the `upd_pub_title` stored procedure has an additional column, `old_state`, that enables replicates to subscribe to rows that migrate.

```sql
-- Execute this script at NY data servers
-- Creates stored procedure
create procedure upd_publishers
    (@pub_id char(4), @pub_name varchar(40),
     @city varchar(20), @state char(2))
as
if not exists
    (select * from publishers where pub_id=@pub_id)
raiserror 20005 "**FATAL ERROR: Unknown publishers id**"
else
    update publishers set
    pub_name=@pub_name,
    city=@city,
    state=@state
    where pub_id = @pub_id
end
/* end of script */
-- Execute this script at NY data servers
-- Creates stored procedure
create procedure upd_titles
    (@title_id tid, @title varchar(80), @type char(12),
     @pub_id char(4)=null, @price money=null, @advance money=null,
     @total_sales int=null, @notes varchar(200)=null,
     @pubdate datetime, @contract bit)
as
if not exists
    (select * from titles where title_id=@title_id)
raiserror 20005 "**FATAL ERROR: Unknown title id**"
else
    update titles set
    title=@title,
    type=@type,
    pub_id=@pub_id,
    price=@price,
    advance=@advance,
    total_sales=@total_sales,
    notes=@notes,
    pubdate=@pubdate,
```
contract=@contract
where title_id = @title_id
end
/* end of script */
-- Execute this script at NY data server
-- Creates stored procedure
create procedure upd_pub_title
(@pub_id char(4), @pub_name varchar(40)=null,
@city varchar(20)=null, @state char(2)=null,
@title_id tid=null, @title varchar(80)=null, @type char(12)=null,
@price money=null, @advance money=null,
@total_sales int=null, @notes varchar(200)=null,
@pubdate datetime=null, @contract bit, @old_state char(2))
as
begin
if not exists (select * from publishers where pub_id=@pub_id)
raiserror 20005 "**FATAL ERROR: Unknown publishers id**"
else
exec upd_publishers @pub_id, @pub_name, @city, @state
if @title_id != null
begin
if not exists
(select * from titles where title_id=@title_id)
exec ins_titles @title_id, @title, @type, @pub_id,
@price, @advance, @total_sales, @notes, @pubdate,
@contract
else
exec upd_titles @title_id, @title, @type, @pub_id, @price,
@advance, @total_sales, @notes, @pubdate, @contract
end
end
/* end of script */

At the replicate site; Update procedures observe the following logic:

- Raise an error on an unknown pub_id.
- If title does not exist, insert one.
- Implement correct update migration as shown in Table 3-1.

**Table 3-1: Migration strategy for replicate site (CA)**

<table>
<thead>
<tr>
<th>Old State</th>
<th>New State</th>
<th>Update Procedure Needs To</th>
</tr>
</thead>
<tbody>
<tr>
<td>CA</td>
<td>CA</td>
<td>Update publishers and titles tables normally.</td>
</tr>
<tr>
<td>CA</td>
<td>NY</td>
<td>Delete publisher and cascade delete of all titles associated with publisher.</td>
</tr>
</tbody>
</table>
These scripts create the `upd_publishers` and `upd_titles` stored procedures and the managing stored procedure `upd_pub_title` that controls the execution of `upd_publishers` and `upd_titles`.

```
-- Execute this script at SF data servers
-- Creates stored procedure
create procedure upd_publishers
(@pub_id char(4), @pub_name varchar(40),
 @city varchar(20), @state char(2))
    as
    if not exists
    (select * from publishers where pub_id=@pub_id)
    raiserror 20005 "**FATAL ERROR: Unknown publishers id**"
    else
    update publishers set
    pub_name=@pub_name,
    city=@city,
    state=@state
    where pub_id = @pub_id
end
/* end of script */
-- Execute this script at SF data servers
-- Creates stored procedure
create procedure upd_titles
(@title_id tid, @title varchar(80), @type char(12),
 @pub_id char(4)=null, @price money=null, @advance money=null,
 @total_sales int=null, @notes varchar(200)=null,
 @pubdate datetime, @contract bit)
    as
    if not exists
    (select * from titles where title_id=@title_id)
exec ins_titles @title_id, @title, @type, @pub_id,
     @price, @advance, @total_sales, @notes, @pubdate,
     @contract
    else
    update titles set
    title=@title,
    type=@type,
    pub_id=@pub_id,
    price=@price,
    advance=@advance,
    total_sales=@total_sales,
    notes=@notes,
```

<table>
<thead>
<tr>
<th>Old State</th>
<th>New State</th>
<th>Update Procedure Needs To</th>
</tr>
</thead>
<tbody>
<tr>
<td>NY</td>
<td>CA</td>
<td>Insert new publisher and title (if any).</td>
</tr>
</tbody>
</table>
pubdate=@pubdate,
contract=@contract
where title_id = @title_id
end

/* end of script */

-- Execute this script at SF data servers
-- Creates stored procedure
create procedure upd_pub_title
(@pub_id char(4), @pub_name varchar(40)=null,
@city varchar(20)=null, @state char(2),
@title_id tid=null, @title varchar(80)=null, @type char(12)=null,
@price money=null, @advance money=null,
@total_sales int=null, @notes varchar(200)=null,
@pubdate datetime=null,@contract bit, @old_state char(2))
as
declare @rep_state char (2)
begin
select @rep_state=state from publishers
where pub_id=@pub_id
if @old_state = @state
begin
exec upd_publishers @pub_id, @pub_name, @city, @state
if @title_id != null
exec upd_titles @title_id, @title, @type,
@pub_id, @price, @advance, @total_sales,
@notes, @pubdate, @contract
end
else if @rep_state = @old_state
begin
exec del_titles @title_id, @pub_id
exec del_publishers @pub_id
end
else if @rep_state = null
begin
exec ins_publishers @pub_id, @pub_name, @city,
@state
if @title_id != null
exec ins_titles @title_id, @title, @type,
@pub_id, @price, @advance, @total_sales,
@notes, @pubdate, @contract
end
end

/* end of script */
Function replication definitions

Create function replication definitions on the primary Replication Server for ins_pub_title, del_pub_title, and upd_pub_title. Note that for inserts and deletes, only state is a searchable column; for updates, old_state is also searchable.

```sql
-- Execute this script at NY data servers
-- Creates replication definition del_pub_title
create function replication definition del_pub_title
with primary at MIAMI_DS.pubs2
  (@pub_id char(4),
   @state char(2),
   @title_id varchar(6))
searchable parameters (@state)
go
/* end of script */

-- Execute this script at NY data servers
-- Creates replication definition ins_pub_title
create function replication definition ins_pub_title
with primary at MIAMI_DS.pubs2
  (@pub_id char(4),
   @pub_name varchar(40),
   @city varchar(20),
   @state char(2),
   @title_id varchar(6),
   @title varchar(80),
   @type char(12),
   @price money,
   @advance money,
   @total_sales int,
   @notes varchar(200),
   @pubdate datetime,@contract bit)
searchable parameters (@state)
go
/* end of script */

-- Execute this script at NY data servers
-- Creates replication definition upd_pub_title
create function replication definition upd_pub_title
with primary at MIAMI_DS.pubs2
  (@pub_id char(4),
   @pub_name varchar(40),
   @city varchar(20),
   @state char(2),
   @title_id varchar(6),
   @title varchar(80),
   @type char(12),
   @price money,
   @advance money,
   @total_sales int,
   @notes varchar(200),
   @pubdate datetime,@contract bit)
searchable parameters (@state)
go
/* end of script */
```
Model variations and strategies

@advance money,
@total_sales int,
@notes varchar(200),
@pubdate datetime,
@contract bit,
@old_state char(2))
searchable parameters (@state, @old_state)
go
/* end of script */

Subscriptions

This script creates the subscriptions at the replicate Replication Server using the no-materialization method. Use this method when you don’t need to load data at the replicate site.

To ensure proper subscription migration, you must create two subscriptions for upd_pub_title.

-- Execute this script at SF data servers
-- Creates subscription for del_pub_title,
ins_pub_title,
    and upd_pub_title
create subscription del_pub_title_sub
for del_pub_title
with replicate at SF_DS.pubs2
where @state='CA'
without materialization
go

create subscription ins_pub_title_sub
for ins_pub_title
with replicate at SF_DS.pubs2
where @state='CA'
without materialization
go

create subscription upd_pub_title_sub1
for upd_pub_title
with replicate at SF_DS.pubs2
where @state='CA'
without materialization
go

create subscription upd_pub_title_sub2
for upd_pub_title
with replicate at SF_DS.pubs2
where @old_state = 'CA'
without materialization

/* end of script */
Model variations and strategies
Planning for Backup and Recovery

This chapter describes the tools and methods you can use to return primary and replicate sites to a consistent state after a system component failure.

<table>
<thead>
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<th>Page</th>
</tr>
</thead>
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<tr>
<td>Preventive measures</td>
<td>88</td>
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<tr>
<td>Recovery measures</td>
<td>91</td>
</tr>
</tbody>
</table>

Protecting against data loss

Replication Server runs in distributed database systems with many other hardware and software components, including Adaptive Servers and other data servers, Replication Agents, LANS, WANs, and client application programs.

Any of these components, including Replication Server, may occasionally fail. Replication Server is a fault-tolerant system, designed with this possibility in mind. During most failures, it waits for the failure to be corrected and then continues its work. When failures require restarting Replication Server or a Replication Agent, the start-up process guarantees that replication is resumed without loss or duplication of data.

Protecting against data loss is the same with a replication system as with a centralized database system. In both cases, there is one definitive version of the data, and you should invest considerable planning and resources to protect it.

In a replication system, if the primary data is protected, all replicate data can ultimately be recovered. A site can replace lost replicate data simply by re-creating its subscriptions.

If replicated transactions are lost at the primary site (as happens, for example, when the primary database is rolled back to a previous dump), consistency may be lost between primary and replicated data.
Preventive measures

The backup and recovery methods available in a replication system include preventive and recovery measures. This chapter describes both of them.

Preventive measures include:

- Warm standby
- Hardware data mirroring (hot standby)
- Longer save intervals
- Coordinated dumps

Recovery measures include:

- Subscription initialization
- Subscription reconciliation utility (rs_subcmp)
- Database recovery
- Restoring coordinated dumps

Preventive measures

The following sections describe the preventive measures you can take to protect data in your replication system.

Standby applications

You can protect the data in your replication system by maintaining separate (standby) copies of primary data. Two possible standby methods are:

- **Warm standby application** – a pair of Adaptive Server or SQL Server databases, one of which is a backup copy of the other. Client applications update the active database; Replication Server maintains the standby database by copying supported operations to the active database.

- **Hardware data mirroring** – a form of hot standby application. With the use of additional hardware, data mirroring maintains an exact copy of the primary data by reproducing all operations on the primary data.
Comparing methods

In a hot standby application, a standby database can be placed into service without interrupting client applications and without losing any transactions. A hot standby database guarantees that transactions committed on the active database are also committed on the standby. When both databases are up, the active database and the standby database are in sync, and the hot standby database is ready for immediate use.

Alternately, a warm standby application maintained by Replication Server:

- Can be used in environments where data mirroring applications cannot, especially when necessary hardware is not available.
- Tolerates temporary network failures better than some hot standby applications because committed transactions can be stored on the active database, even when the standby database is down.
- Minimizes overhead on the active database because the active database does not need to verify that the databases are in sync.

However, a warm standby application maintained by Replication Server also:

- Requires some interruption of client applications when switching to the standby database.
- May not have executed in the standby database the most recent transactions committed in the active database.

Warm standby

Replication Server’s warm standby application is described in detail in Chapter 13, “Managing Warm Standby Applications,” in the Replication Server Administration Guide.

**Note** Replication Server version 12.0 supports Sybase Failover available in Adaptive Server Enterprise version 12.0. Failover support is not a substitute for warm standby. While warm standby applications keep a copy of a database, Failover support accesses the same database from a different machine. Connections from Replication Server to warm standby databases work the same way.

For detailed information about how Failover support works in Replication Server, see “Configuring the Replication System to Support Sybase Failover” in Chapter 16, “Replication System Recovery,” of the Replication Server Administration Guide.
Hardware data mirroring

To ensure the highest data availability, you can mirror the most critical data in the replication system. Mirroring duplicates I/O operations to maintain two identical copies of the data.

If the active media fails, the standby is brought online instantly. Mirroring all but eliminates the possibility of transaction loss.

The most beneficial places to use mirroring in a replication system are listed here in priority order:

1. **Primary database transaction logs**
   
   Transaction logs store transactions that have not been dumped to tape. If the primary transaction log is lost, transactions must be resubmitted.

2. **Primary database**
   
   A database can be recovered by reloading a previous database dump and subsequent transaction dumps. However, recovering a database that stores primary data also requires recovering or reinitializing the data that has been replicated throughout the enterprise. Extended downtime is often catastrophic for OLTP systems. Mirroring the primary data can prevent this type of catastrophe.

3. **Replication Server stable queues**
   
   Replication Server stores transactions in store-and-forward disk queues called stable queues. It allocates the queues from disk partitions assigned to the Replication Server using the `add partition` command.

   The data stored in stable queues is redundant; it originates in the primary database transaction log. However, if a stable queue is lost, Replication Server cannot deliver transactions to replicate sites. As a result, subscriptions at replicate sites must be reinitialized. Mirroring disk partitions protects stable queues and minimizes potential downtime for replicate databases.

4. **Replication Server System Database (RSSD)**
   
   Recovering from a failure of the RSSD can be a complex process if data such as replication definitions, subscriptions, routes, or function or error classes have been modified since the last backup. Refer to Chapter 16, “Replication System Recovery,” in the *Replication Server Administration Guide* for detailed recovery information.
Mirroring the RSSD can prevent system data loss and the necessity of a complex recovery process. If you don’t mirror the RSSD, be sure to back up the RSSD after any RCL operation that changes system data.

**Save interval**

You can configure a route from one Replication Server to another, or a connection from a Replication Server to a database, so that the Replication Server stores stable queue messages for a period of time after delivery to the destination. This period of time is called the save interval.

The save interval creates a backlog of messages at the source, which allows the replication system to tolerate a partition failure if it is corrected within the save interval. If the stable queues at the destination fail, you can rebuild them and have the source Replication Server resend the backlogged messages.

Refer to the *Replication Server Administration Guide* for details.

**Coordinated dumps**

When a database must be recovered by restoring a backup, replicated data in the affected databases at other sites must somehow be made consistent with the primary data. Replication Server provides a method for coordinating database dumps and transaction dumps at all sites in a distributed system. A database dump or transaction dump is initiated from the primary database. The Replication Agent retrieves the dump record from the log and submits it to the Replication Server so that the dump request can be distributed to the replicate sites. This method guarantees that the data can be restored to a known point of consistency.

A coordinated dump can be used only with databases that store primary data or replicated data, but not both. It is initiated from within a primary database.

Refer to the *Replication Server Administration Guide* for instructions on creating coordinated dumps.

**Recovery measures**

The following sections describe methods for recovering data that is lost after a component failure in a replication system.
Recovery measures

Re-creating subscriptions

The primary version of the data is definitive, so all inconsistencies should be resolved in its favor. One way to recover from a failure at a remote site is to re-create the subscriptions. This recovery method is most expedient for small replicated tables. Large subscriptions and primary data failures require other recovery methods.

Subscription reconciliation utility (\textit{rs\_subcmp})

\textit{rs\_subcmp} tests for rows that are missing, orphaned, or inconsistent in a replicate table and corrects the discrepancies. Using \textit{rs\_subcmp} may be more appropriate for recovering from minor inconsistencies than a more disruptive recovery procedure such as a coordinated load. See the \textit{Replication Server Reference Manual} for instructions on executing \textit{rs\_subcmp}.

Database recovery

When a primary database fails, all committed transactions can be recovered if the database and the transaction log are undamaged. If the database or the transaction log is damaged, you must load a database dump and transaction dumps to bring the database to a known state, and then resubmit the transactions that were executed after the last transaction dump.

When you run Replication Server and Replication Agents in recovery mode, you can replay transactions from reloaded dumps and make sure that all transactions in the primary database are replicated and that no transactions are duplicated. The \textit{Replication Server Administration Guide} describes recovering primary databases from dumps.

Restoring coordinated dumps

Restoring database or transaction dumps created by the coordinated dump process returns the primary and replicated data to a previous, consistent state.

For more information about the coordinated load procedure, see Chapter 16, “Replication System Recovery,” in the \textit{Replication Server Administration Guide}.
CHAPTER 5

Introduction to Replication Agents

This chapter provides an overview of the Replication Agent component of a Sybase replication system. It also provides details about how the RepAgent for Adaptive Server and other Replication Agent products work.

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Replication Agent overview

A Replication Agent is a Replication Server client that retrieves information from the transaction log for a primary database and formats it for the primary Replication Server. RepAgent is the Replication Agent component for Replication Server version 11.5 and later. Log Transfer Manager (LTM) is the Replication Agent component for Sybase SQL Server version 11.0.x and earlier.

Replication Agent detects changes to primary data and ignores changes to nonprimary data. Using Log Transfer Language (LTL), a subset of Replication Control Language (RCL), Replication Agent sends changes in primary data to the primary Replication Server, which distributes the information to replicate databases.

Sybase provides Replication Agent components for other Sybase and non-Sybase data servers, including Adaptive Server Anywhere, DB2, Informix, Microsoft SQL Server, and Oracle.

Replication Agent for DB2 is the Replication Agent component for the MVS-based DB2 database server. Sybase Replication Agent is the Replication Agent component for Informix, Microsoft SQL Server, and Oracle. For database management systems that Sybase does not supply a Replication Agent component for, you must create your own.
RepAgent is an Adaptive Server thread. LTM is a separate process for Sybase SQL Server databases. Replication Agent for DB2 is a separate process that resides on the MVS host. Sybase Replication Agent for Informix, Microsoft SQL Server, and Oracle is a separate application that resides on a LAN server host.

A Replication Agent performs these tasks:

1. Logs in to the Replication Server.
2. Sends a connect source command to identify the session as a log transfer source and to specify the database for which transaction information will be transferred.
3. Gets the name of the maintenance user for the database from the Replication Server. RepAgent filters out operations executed by the maintenance user, unless the send_maint_xacts_to_replicate or send_warm_standby_xacts configuration parameter is set to true.
4. Requests the secondary truncation point for the database from the Replication Server. This returns a value, called the origin queue ID, that RepAgent uses to find the location in the transaction log where it is to begin transferring transaction operations. The Replication Server has already received operations up to this location.
5. Retrieves records from the transaction log, beginning at the record following the secondary truncation point, and formats the information into LTL commands.

Transaction log

A Replication Agent can use one of two types of transaction logs:

- The native transaction log of the primary database
- A transaction log created and maintained by the Replication Agent

Native transaction logs are used by the RepAgent thread (for Adaptive Server), LTM (for earlier versions of Sybase SQL Server), and the Replication Agent for DB2.

Sybase Replication Agent (for Informix, Microsoft SQL Server, and Oracle) uses a transaction log that it creates and maintains in the primary database. The Sybase Replication Agent transaction log consists of database objects (triggers, tables, and procedures) in the primary database.
CHAPTER 5   Introduction to Replication Agents

Native transaction logs

A native database transaction log is the ideal source of information about changes to primary data because it is maintained by the database itself and it contains permanent, recoverable updates. If a Replication Agent cannot access the native transaction log of the primary database, the Replication Agent must use another source of reliable information.

Figure 5-1 shows the flow of data between a data server, transaction log, Replication Agent, and Replication Server.

Figure 5-1: The database transaction log

Replication Agent transaction logs

A Replication Agent cannot always access the native transaction log of a database to acquire the information necessary to replicate transaction operations.

Sybase Replication Agent uses its own proprietary transaction log to capture and record transactions in the primary database for replication. The Replication Agent transaction log consists of a set of shadow tables, stored procedures, and triggers that the Replication Agent creates in the primary database.

Sybase Replication Agent generates a SQL script that runs in the primary database. This script creates the Replication Agent transaction log components. See Sybase Replication Agent Administration Guide for more information on the Replication Agent transaction log.
Replication Agent products

Sybase Replication Agent products extend the capabilities of Replication Server by allowing non-Sybase database servers to be primary data servers in a Sybase replication system.

Sybase offers the following Replication Agent products for non-Sybase databases:

- Replication Agent for DB2
- Sybase Replication Agent (for Informix, Microsoft SQL Server, and Oracle)

The following sections describe these Replication Agent products.

Replication Agent for DB2

Replication Agent for DB2 is a component of a replication system that moves primary database transactions from DB2 through Replication Server to replicate databases.

Figure 5-2 illustrates the role Replication Agent for DB2 plays in the replication system.

*Figure 5-2: Replication Agent for DB2 data flow*

Replication Agent for DB2 fits into a replication system as follows:

- With Replication Agent for DB2, the primary database is DB2, which runs as a subsystem in MVS. The transaction logs are DB2 logs.
- Replication Agent for DB2 provides a log extract, called Replication Extract, that reads the DB2 logs and retrieves the relevant DB2 active and archive log entries for tables marked for replication.
CHAPTER 5  Introduction to Replication Agents

- LTM for MVS receives the data marked for replication from Replication Extract and transfers this data to Replication Server using the TCP/IP communications protocol.
- Replication Server then applies the changes to the replicate databases.

**DB2 transaction log**

The DB2 database server logs changes to rows in DB2 tables as they occur. The information written to the transaction log includes copies of the data before and after the changes. In DB2, these records are known as undo and redo records. Control records are written for commits and aborts. These records are translated to commits and rollbacks.

The DB2 log consists of a series of data sets. Replication Extract uses these log data sets to identify DB2 data changes. Since DB2 writes change records to the active log as they occur, Replication Extract can process the log records immediately after they are entered.

**Sybase Replication Agent**

Sybase Replication Agent is the Sybase replication system component that captures transactions in an Informix, Microsoft SQL Server, or Oracle primary database, and then transfers transactions to a primary Replication Server for distribution to replicate databases.

Sybase Replication Agent uses a transaction log that records copies of the transactions (or procedure invocations) to be replicated. The Log Reader component of Replication Agent reads the transaction log to retrieve the transaction from the primary database.

After the transaction is retrieved from the primary database, the Log Transfer Interface (LTI) component of Replication Agent processes the transaction and the resulting "change set" data and generates LTL output, which Replication Server uses to distribute the transaction to the subscribing replicate database(s).
Sybase Replication Agent uses information stored in the Replication Server System Database (RSSD) of the primary Replication Server to determine how to process the replicated transactions to generate the most efficient LTL.

After it receives LTL from Replication Agent, the primary Replication Server sends the replicated transaction to a replicate database, either directly or by way of a replicate Replication Server. The replicate Replication Server converts the replicated data to the native language of the replicate database, and then sends it to the replicate database server for processing. When the replicated transaction is processed successfully by the replicate database, the replicate database is synchronized with the primary database.

Sybase Replication Agent runs as a stand-alone application, independent of the primary database server, the primary Replication Server, and any other components of a replication system.

Sybase Replication Agent can reside on the same host machine as the primary database or any other component of the replication system, or it can reside on a separate machine from any other replication system components.

**Replication Agent communications**

Sybase Replication Agent uses the Java Database Connectivity (JDBC) protocol for all its communications.

Replication Agent uses a single instance of jConnect for JDBC to manage all of its connections to Open Client/Open Server applications, including the primary Replication Server and its RSSD and Replication Server Manager (RSM).
In the case of the primary database server, Sybase Replication Agent connects to the JDBC driver provided by the primary database vendor.

While replicating transactions, Replication Agent maintains connections with both the primary database and the primary Replication Server. In addition, Replication Agent occasionally connects to the RSSD of the primary Replication Server to retrieve replication definition data.

**Java implementation**

Sybase Replication Agent components are implemented in the Java programming language. Therefore, to run Sybase Replication Agent, you must have a Java Runtime Environment (JRE) installed on the computer that will act as the Replication Agent host machine.
Replication Agent products
CHAPTER 6

Replicating Data into Foreign Data Servers

This chapter describes the replication system components required to replicate data into foreign data servers—those that are not Adaptive Server or Sybase SQL Server.

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Interfacing with foreign data servers

Replication Server updates the replicate data stored in databases by submitting requests to data servers. Support for Adaptive Server Enterprise and SQL Server is provided with Replication Server. If your database is managed by a data server other than Adaptive Server or SQL Server, you must provide an interface for Replication Server to use.

This interface includes:

- An Open Server gateway application to receive instructions from Replication Server and apply them to the data server.
- A maintenance account that Replication Server can use to log in to the gateway.
- A function-string class to use with the database. The function strings in the class tell Replication Server how to format requests for the data server.
Sybase database gateway products

- An error class and error action assignments to handle errors the data server returns to Replication Server via the gateway.
- An `rs_lastcommit` table in each database that has replicated data. Replication Server uses this table to keep track of the transactions that have been successfully committed in the database.
- An `rs_get_lastcommit` function call to retrieve the last transaction from each source primary database.

Sybase offers several Open Server gateway application products that you can use to access non-Sybase database servers for a replicate database.

The heterogeneous datatype support (HDS) feature in Replication Server version 12 provides several components of this interface for supported database servers (DB2, Informix, Microsoft SQL Server, and Oracle). The HDS feature provides function string classes, error classes, and scripts to create the necessary tables and procedures in the replicate database.

See the Replication Server Administration Guide and the Replication Server Configuration Guide for your platform for more information about the HDS feature.

**Sybase database gateway products**

The Sybase DirectConnect data access servers are Sybase's middleware building blocks for connectivity between LAN-based clients and enterprise data sources. Using a DirectConnect data access server simplifies the integration of non-Sybase replicate databases into a Sybase replication system.

You can use a DirectConnect server on its own to connect directly to a database server. The DirectConnect server acts as an Open Server gateway by interpreting the Open Client/Server protocol used by Replication Server to the native communication protocol used by the non-Sybase replicate database.

Sybase provides DirectConnect data access libraries for the following data sources and database servers:

- AS/400
- Informix
- Microsoft SQL Server
- MVS (DB2)
- Oracle
See the *DirectConnect Access Service User's Guide* or the *DirectConnect Server Administration Guide* for your platform for more information about DirectConnect data access servers.

**Open Server gateway application**

The gateway application uses Open Server Server-Library/C routines so that it can accept a login from Replication Server and process Replication Server requests. It forwards the requests to the foreign data server with the replicated data and returns the status of the request to Replication Server through Server-Library routines.

---

**Note** This book describes the requirements of an Open Server gateway to a data server, but it does not discuss Open Server programming. Refer to the *Open Server Server-Library/C Reference Manual* for information about programming with Open Server.

---

When you design a gateway for a data server, you must decide how the function strings will be written and how the gateway will process them. For example, if you use language function strings, the gateway will need a parser, and you will have to define the syntax of the function strings so that the gateway can parse them.

If you create RPC function strings, you can use Open Server registered procedures to process them. This is easier than creating a parser.

The gateway application must:

- Use the data server vendor’s application programming interface (API) to access the database
- Perform any translation needed to reformat the requests received from Replication Server function calls to the format required by the database API
- Translate row results, messages, and errors into the Open Server result formats and return them to Replication Server

You must define the Open Server gateway in the interfaces file for the Replication Server.

Use the `create connection` command to create a connection to the gateway from the Replication Server managing the database.
Maintenance user

Replication Server logs in to the gateway as the maintenance user specified in create connection for the database. The gateway can log in to the data server with the same login name, or it can use another login name. The only requirements are that the login name must have the permissions needed to modify the replicate data.

Function-string class

The Replication Server managing a database requires a function-string class. Replication Server provides function-string classes for Adaptive Server, and with its heterogeneous datatype support (HDS) feature, Replication Server version 12 provides function-string classes for DB2, Informix, Microsoft SQL Server, and Oracle data servers.

If you are replicating to a non-Sybase data server that is not supported by the HDS feature, you must create a function-string class for that data server. You can either:

- Create a function-string class that inherits function strings from a system-provided class, or
- Create all the function strings yourself.

Replication Server sends the gateway a command that it constructs by mapping runtime values into the function string supplied for the function. Depending on how the function string is written and the requirements of the data server, your gateway can pass the command directly to the data server or process it in some way before it sends the request to the data server.

Refer to Chapter 12, “Customizing Database Operations,” in the Replication Server Administration Guide for a list of Replication Server system functions that your database gateway may need to process.
Creating function-string classes using inheritance

Replication Server lets you share function-string definitions between function-string classes by creating relationships between classes using a mechanism called function-string inheritance.

The system-provided classes `rs_default_function_class` and `rs_db2_function_class` can serve as parent classes for derived classes that inherit function strings from the parent class. You can create a derived class in order to customize certain function strings for your data server while retaining all other function strings from the parent class.

Use the `create function string class` command to create a derived class from the parent class `rs_default_class` or `rs_db2_function_class` that inherits from the parent class. Create customized function strings only as needed.

**Note** `rs_db2_function_class` does not support replication of text or image data. To enable replication of text or image data for DB2 databases, you must customize the `rs_writetext` function string using the RPC method through a gateway. Refer to “Creating distinct function-string classes” on page 105 for information about `rs_writetext`.

Refer to Chapter 12, “Customizing Database Operations,” in the Replication Server Administration Guide for a detailed discussion of function-string inheritance.

Creating distinct function-string classes

If you use a class that does not inherit from a system-provided class, you must create all function strings yourself, and add new function strings whenever you create a new table or function replication definition.

Use the `create function string class` command to create a new function-string class, and then create function strings for all of the functions with function-string-class scope.

You must create `rs_insert`, `rs_update`, and `rs_delete` function strings for each table you replicate in the database.

If you are replicating columns with text or image datatypes, you must create `rs_datarow_for_writetext`, `rs_get_textptr`, `rs_textptr_init`, and `rs_writetext` function strings for each text or image column. The function-string name must be the text or image column name for the replication definition.
The `rs_select` and `rs_select_with_lock` function strings are needed only if the database has the primary data for a replication definition.

### Error class

An error class determines how Replication Server handles the errors that are returned by your gateway. You must use the `create error class` command to create an error class for your gateway.

Use the `assign action` command to tell Replication Server how to respond to the errors returned by your gateway. Table 6-1 lists the possible actions.

<table>
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<tr>
<th>Action</th>
<th>Description</th>
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<tr>
<td>ignore</td>
<td>Assume that the command succeeded and that there is no error or warning condition to process. This action can be used for a return status that indicates successful execution.</td>
</tr>
<tr>
<td>warn</td>
<td>Log a warning message, but do not roll back the transaction or interrupt execution.</td>
</tr>
<tr>
<td>retry_log</td>
<td>Roll back the transaction and retry it. The number of retry attempts is set with <code>configure connection</code>. If the error recurs after retrying, write the transaction into the exceptions log and continue, executing the next transaction.</td>
</tr>
<tr>
<td>log</td>
<td>Roll back the current transaction and log it in the exceptions log. Then continue, executing the next transaction.</td>
</tr>
<tr>
<td>retry_stop</td>
<td>Roll back the transaction and retry it. The number of retry attempts is set with <code>configure connection</code>. If the error recurs after retrying, suspend replication for the database.</td>
</tr>
<tr>
<td>stop_replication</td>
<td>Roll back the current transaction and suspend replication for the database. This is equivalent to using <code>suspend connection</code>, and this is the default action. Since this action stops all replication activity for the database, it is important to identify the data server errors that can be handled without shutting down the database connection and assign them another action.</td>
</tr>
</tbody>
</table>

The default error action is `stop_replication`. If you do not assign another action to an error, Replication Server shuts down the connection to your gateway.
See the Replication Server Reference Manual for more information about create error class and assign action.

**rs_lastcommit table**

Each row in the rs_lastcommit table identifies the most recent committed transaction that was distributed to the database from a primary database. Replication Server uses this information to ensure that all transactions are distributed.

The rs_lastcommit table should be updated by the rs_commit function string before the transaction is committed. This guarantees that the table is updated with every transaction Replication Server commits in the database.

Replication Server maintains the rs_lastcommit table as the maintenance user for the database. You must make sure that the maintenance user has all of the permissions needed for the table.

Table 6-2 lists the columns in the rs_lastcommit table.

<table>
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<tr>
<th>Column name</th>
<th>Datatype</th>
<th>Description</th>
</tr>
</thead>
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<tr>
<td>origin</td>
<td>int</td>
<td>An integer assigned by Replication Server that uniquely identifies the database where the transaction originated</td>
</tr>
<tr>
<td>origin_qid</td>
<td>binary(36)</td>
<td>The origin queue ID for the commit record in the transaction</td>
</tr>
<tr>
<td>secondary_qid</td>
<td>binary(36)</td>
<td>A queue ID for a stable queue used during subscription materialization</td>
</tr>
<tr>
<td>origin_time</td>
<td>datetime</td>
<td>Time at origin for the transaction</td>
</tr>
<tr>
<td>dest_commit_time</td>
<td>datetime</td>
<td>Time the transaction was committed at the destination</td>
</tr>
</tbody>
</table>

The origin_time and dest_commit_time columns are not required.

The origin column is a unique key for the table. There is one row for each primary database whose data is replicated in this database.

If you use a coordinated dump with the database, you should update rs_lastcommit with the rs_dumpdb and rs_dumptran function strings.
For Adaptive Server and SQL Server databases, the `rs_commit`, `rs_dumpdb`, and `rs_dumptran` function strings execute a stored procedure named `rs_update_lastcommit` to update the `rs_lastcommit` table. This is the text of that stored procedure:

```sql
/* Create a procedure to update the */
** rs_lastcommit table. */
create procedure rs_update_lastcommit
    @origin         int,
    @origin_qid     binary(36),
    @secondary_qid  binary(36),
    @origin_time    datetime
as
    update rs_lastcommit
        set origin_qid = @origin_qid,
        secondary_qid = @secondary_qid,
        origin_time = @origin_time,
        commit_time = getdate()
    where origin = @origin
    if (@@rowcount = 0)
    begin
        insert rs_lastcommit (origin,
            origin_qid, secondary_qid,
            origin_time, commit_time,
            pad1, pad2, pad3, pad4,
            pad5, pad6, pad7, pad8)
            values (@origin, @origin_qid,
                @secondary_qid, @origin_time,
                getdate(), 0x00, 0x00, 0x00,
                0x00, 0x00, 0x00, 0x00)
    end
    go
```

See the reference pages for the `rs_commit`, `rs_dumpdb`, and `rs_dumptran` functions in the *Replication Server Reference Manual* for more information.

**rs_get_lastcommit function**

Replication Server sends an `rs_get_lastcommit` function call to the gateway to retrieve the last transaction committed in the database from each source primary database. The gateway is expected to return the first three columns described in Table 6-2.
The function string for `rs_get_lastcommit` can execute a simple `select`:

```
select origin, origin_qid, secondary_qid
from rs_lastcommit
```

### Using Sybase Central with foreign data servers

You can use Sybase Central’s RSM plug-in to manage a replication system that includes non-Sybase data servers. The following requirements and restrictions apply:

- You must configure an Open Server gateway application for each non-Sybase data server and complete the interface described in “Interfacing with foreign data servers” on page 101.
- The logical view of the Replication Server tree will not be visible in Sybase Central, but you can view data-specific information about replication definitions, subscriptions, and publications in the property sheets for Replication Server and the database.
- You cannot use Sybase Central to create replication definitions or publications, but you can use Sybase Central to create subscriptions for replication definitions or publications created with RCL commands and a utility program such as `isql`. 
Using Sybase Central with foreign data servers
International Replication Design Considerations

This chapter discusses issues that pertain to setting up a replication system in an international environment.

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Designing an international replication system

Replication Server and Replication Server Manager (RSM), the Sybase Central graphical plug-in that you can use to manage your replication system, support international environments. They provide these features:

- Localization of messages into English, French, German, and Japanese
- Support for all Sybase-supported character sets, with character-set conversion between Replication Server sites
- Support for nonbinary sort orders

When you design a replication system for an international environment, it is important to understand the impact that language, character set, and sort order settings have on your system. Replication Server and the RSM Server provide great flexibility in the configuration of these settings. This flexibility may lead to unexpected or unwanted results unless you follow the configuration guidelines presented in this chapter.
Message language

You can configure Replication Server and RSM Server to print messages to the error log and to clients in English, French, German, or Japanese. The language you choose must be compatible with the chosen character set. English is the default language; it is compatible with all Sybase character sets.

Each server program in your replication system, including Replication Server, Adaptive Server, other data servers, RSM Server, and the LTM, writes messages to its error log in the configured language. However, whether messages are sent to a client in the client’s language depends on the server.

For example, Adaptive Server checks for the language setting of its client (Replication Server) and returns messages in that language. RepAgent, an Adaptive Server thread, also returns messages in the client language.

However, Replication Server, RSM Server, and LTM do not check for a client’s language; instead, they return messages to a client in their own language. Thus, error logs can contain messages in different languages if the servers are configured in different languages.

**Note** To avoid the confusion that can result from a mixed-language error log, configure the same language setting for all servers and clients at a given site.

Character sets

Replication Server and RSM Server support all Sybase-supported character sets and perform character-set conversion of data and identifiers between primary and replicate sites. Character sets must be compatible for character-set conversion to be successful. For example, single-byte character set data cannot be converted to a multibyte character set. For details about character-set compatibility, see the *Open Client Client-Library Reference Manual*.

Your choice of a character set for a given server is influenced by the languages the server supports, the hardware and operating system it runs on, and the systems with which it interacts.

These things are true of Sybase-supported character sets:

- They are all supersets of 7-bit ASCII.
- Some are completely incompatible with each other, meaning that they have no characters in common beyond 7-bit ASCII.
• Among compatible character sets, some characters are not common to both sets—that is, no two character sets have all the same characters.

Character-set conversion

Character-set conversion takes place at the destination Replication Server. Every message packed for the Replication Server Interface (RSI) includes the name of the source Replication Server’s character set. The destination Replication Server uses this information to convert the character data and identifiers to its own character set.

When it attempts character-set conversion, Replication Server checks to see whether the character sets are compatible. If they are incompatible, no conversion occurs. If they are compatible and one or more characters that are not common to both sets are encountered, a ? (question mark) is substituted for the unrecognized characters.

In the Replication Server, context determines whether a ? is substituted or a character-set conversion exception is raised. For example, if Replication Server detects an incompatibility in a character being replicated, it substitutes a ? for the character; if it detects an incompatibility when converting the configuration file parameters, it prints an error message and shuts down.

You can use the dsi_charset_convert configuration parameter to specify whether or not Replication Server attempts character-set conversion. For a description of this parameter, see the configure connection command in the Replication Server Reference Manual.

For details about how Replication Server handles character-set conversion during subscription materialization, resolution, and reconciliation, see “Subscriptions” on page 116.

RSM Server issues

In general, RSM Server and Replication Server support internationalization in the same way. This section describes additional character set issues for RSM Server.
Character sets

Reading error log files
You can configure RSM Server to read Replication Server and Adaptive Server (or SQL Server) error log files. However, RSM Server can directly monitor only the log files of servers that have the same character set.

To monitor a server with a character set that is different from that of your RSM Server—for example, the server log file is printed in SJIS and your RSM Server is configured to use EUCJIS—you must monitor the server indirectly through another RSM Server. Set up a local RSM Server and configure it with the other server’s character set. Your RSM Server can then read the log file through the local RSM Server.

Refer to “Rep Server: Monitoring error log locally” in Replication Server’s online help for more information.

Changing the RSM Server character set
RSM Server assumes that all configuration files are written in the RSM Server character set. (With Replication Server you can write the configuration file in one character set and configure Replication Server to run with a different character set.) If you change the RSM Server character set, you must run rsmgen (UNIX) or RSMSetup (Windows NT) to rewrite configuration parameters in the new character set. Do not change values in the configuration file directly.

Refer to the Replication Server Configuration Guide for your platform for information about running rsmgen and RSMSetup.

RSM Client issues
Sybase Central and the Sybase Replication Server plug-in—the RSM Client—use the local $SYBASE installation’s default locale (language and charset) settings to connect to the RSM Server. You cannot configure Sybase Central and the RSM Client locale settings directly. Refer to the Open Client/Server Configuration Guide for detailed information about configuring local $SYBASE locale settings.

It is strongly recommended that the machine on which the RSM Client is installed and the RSM Server use the same language and charset. Otherwise, messages returned from RSM Server will be in a different language or, if the charsets are not compatible, unintelligible.
Guidelines for using character sets

In setting up a replication system, it is strongly recommended that all servers at a given Replication Server site use the same character set. It is also recommended that all of the Replication Servers in your replication system use compatible character sets.

Follow these guidelines to minimize character-set conversion problems:

- Use 7-bit ASCII (if possible) for all data servers, data, and object names.
  
  If your data and object names are all 7-bit ASCII or if all of your data servers and Replication Servers use the same character set, character set conversion will not present problems.

- If you need to replicate data between a single-byte and a multibyte server, restrict character data and object names to 7-bit ASCII to avoid corruption. Otherwise, you may experience problems. For example, Replication Server does not restrict server names to 7-bit ASCII, but Adaptive Server or the Connectivity Libraries may do so.

- When replicating between servers with different but compatible character sets (for example, ISO_1 and CP850), make sure that object names and character data do not include any 8-bit characters that are not common to both character sets.

It is recommended that RSM Server and all servers managed by RSM Server use the same or compatible character sets. If character sets used by servers in the RSM Server domain are different, configure the RSM Server with a character set that is a superset of the character sets of the servers it manages. For example, if servers in an RSM Server domain are configured with ISO_1 and SJIS, configure the RSM Server with SJIS, which is a superset of ISO_1.
Sort order

Replication Server and RSM Server use sort orders, or collating sequences, to determine how character data and identifiers are compared and ordered. Replication Server and RSM Server support all Sybase-supported sort orders, including non-binary sort orders. Non-binary sort orders are necessary for the correct ordering of character data and identifiers in European languages.

Note To make sure that data and identifiers are ordered consistently across your replication system, configure the RSM Server and all Replication Server components with the same sort order.

Subscriptions

Subscriptions involve comparisons of data at the:

- Primary data server during materialization
- Primary Replication Server during resolution
- Replicate Replication Server during initialization and dropping
- Replicate data server during dropping

Sort order and character-set conversion play important roles in processing subscriptions and must be consistent everywhere for subscriptions to be valid.

Subscription materialization

Figure 7-1 illustrates the typical message flow during subscription materialization.
During subscription materialization:

- The replicate Replication Server logs in to the primary data server and issues a `select` statement to retrieve the primary data.
- The primary data server converts all character data to the replicate Replication Server’s character set. The replicate Replication Server’s character set, if it is different, must be installed at the primary data server.
- The replicate Replication Server inserts the data at the replicate data server.

**Note** In bulk materialization, a subscription is initialized by a user-chosen mechanism outside the replication system. Therefore, you must make sure that the initial data selection at the primary data server uses the correct sort order and that the character data is converted to the replicate data server character set, if need be.

**Subscription resolution**

Figure 7-2 illustrates the typical message flow during the normal lifetime of a subscription.
During subscription resolution:

- Adaptive Server’s RepAgent thread scans the log for updates.
- The primary Replication Server uses its sort order to determine what rows qualify for the subscription. The primary Replication Server also adds the name of the primary Replication Server’s character set to the RSI message.
- The replicate Replication Server converts the data to its character set, if necessary, and applies updates to the replicate data server.

**Subscription reconciliation**

Figure 7-3 illustrates the rs_subcmp process during subscription reconciliation.

During subscription reconciliation:

- rs_subcmp connects to the primary data server and to the replicate data server using the replicate data server’s character set.
• The primary data server converts all character data to the replicate data server’s character set (all rs_subcmp operations are performed in the replicate data server’s character set). The replicate data server’s character set, if it is different from the primary data server’s character set, must be installed at the primary data server.

• rs_subcmp sends a select statement to both data servers. The sort order of each data server must be the same for this process to produce sensible results.

Dematerialization

Figure 7-4 illustrates the typical message flow during subscription dematerialization.

Figure 7-4: Subscription dematerialization

During subscription dematerialization:

• The replicate Replication Server selects data from the replicate data server to construct the dematerialization queue. The replicate data server uses its sort order to select the rows.

• The replicate Replication Server uses its sort order to throw out rows belonging to other subscriptions.

• The replicate Replication Server deletes the remaining rows from the replicate database.
Summary

- Replication Server and RSM Server can be configured to print messages to error logs and to clients in English, French, German, and Japanese. English is the default language.

- Sybase recommends that all servers at a replication site be configured with the same language.

- Replication Server and RSM Server support all Sybase-supported character sets and sort orders, including non-binary sort orders.

- Replication Server performs character-set conversion of data and identifiers between primary and replicate Replication Servers and databases.

- Sybase recommends that all servers at a replication site use the same character set and that all Replication Servers and RSM Servers in your system use compatible character sets.

- Sort order plays an important role in processing subscriptions, and it must be consistent everywhere for subscriptions to be valid.
APPENDIX A

Capacity Planning

This appendix contains information to help you plan the CPU, memory, disk, and network resources you need for your replication system.

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Overview of requirements

A replication system consists of Replication Servers, Replication Agents (RepAgent, LTM, or other Replication Agent), and data servers.

Replication Server requirements

The minimum requirements for each Replication Server are:

- One Replication Agent for the Replication Server System Database (RSSD) if there will be a route from this Replication Server.
- At least one 20MB raw partition or operating-system file for the stable queues.
- An Adaptive Server for the RSSD, which must have:
  - At least 10MB of free device space for the RSSD
  - Another 10MB of free device space for the log
Overview of requirements

- 20 user connections for the RSSD, in addition to the number of user connections needed by Adaptive Server users. When Replication Server starts up, several threads attempt to read the RSSD at the same time. To accommodate this demand, increase the number of user connections by 20.
- One RSSD user connection for every RSM in the replication system and one user connection per data server for each RSM process.
- Two user connections for each database containing primary data.
- One user connection for each replicate-only database.
- At least 8MB of RAM for the Replication Server executable program, plus data and stack memory.
- At least 5MB of RAM for each LTM or other Replication Agent. (RepAgent is an Adaptive Server thread; it doesn’t require any Replication Server memory.)

Replication Server requirements for primary databases

For each primary database it manages, a Replication Server needs:

- RepAgent thread for Adaptive Server databases, LTM for Sybase SQL Server databases, or other Replication Agent for non-Sybase databases
- One inbound stable queue
- One outbound stable queue
- One or two connections to the data server (one for the LTM and one for the DSI)

See the release bulletin for more information about Adaptive Server compatibility requirements.

Note  If you are using a Replication Agent with a non-Sybase data source, see the appropriate Replication Agent documentation for information about compatibility requirements.

Replication Server requirements for replicate databases

For each replicate database it manages, a Replication Server needs:
• One outbound stable queue
• One connection to the data server (DSI)

Replication Server requirement for routes

For each direct route to another Replication Server, a Replication Server needs:
• One outbound stable queue

Data volume (queue disk space requirements)

The most significant components in estimating the amount of resources required by the replication system are the volume of the data being replicated and the rate at which it is being updated.

To calculate data volume, you need to know the following things about your replication system:
• The number of sites
• The widths of the rows in replicated tables
• The widths of parameters in replicated functions
• The number of modifications per second
• The duration of a typical transaction
• The selectivity of replication for replicated tables (the fraction of that table replicated through the queue)
• The length of time you want queues to hold transactions when a destination is unavailable

The following sections provide formulas for calculating queue size requirements. You can also use the rs_fillcaptable and rs_capacity stored procedures in an RSSD to get a queue size estimate. See the Replication Server Reference Manual for information about these stored procedures.
Overview of disk queue size calculation

Figure A-1 illustrates the sequence and flow of calculating data volume and queue disk usage

Figure A-1: Calculating queue disk usage

Message sizes

A Replication Server distributes database modifications using messages in ASCII format. Most replication system messages correspond to the delete, insert, update, and function execution operations.

Each table has one message size for inserts or deletes and another for updates. Each function has its own message size. Message sizes are expressed in bytes.

Message sizes for the same type of modification (insert, delete, or update) may vary depending on whether the message is in the inbound or outbound queue.

You can calculate message sizes in bytes for table updates, inserts, and deletes, for functions, and for begin/commit pairs using the formulas presented below. See “Formula components” on page 126 for a description of the components.
Table updates

This formula will give you a rough estimate of the upper limit of message size:

\[
\begin{align*}
\text{InboundMsgSize update} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + (\text{RowWidth} \times 2) \\
\text{OutboundMsgSize update} &= \text{OutboundMsgOverhead} + \\
&\quad (\text{RowWidth} \times 2) + (\text{NumSites} \times 8)
\end{align*}
\]

To get a more precise estimate, use the RowWidthChanged figure in the calculation, as follows:

\[
\begin{align*}
\text{InboundMsgSize update} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + \text{RowWidth} + \text{RowWidthChanged} \\
\text{OutboundMsgSize update} &= \text{OutboundMsgOverhead} + \\
&\quad \text{RowWidth} + \text{RowWidthChanged} + (\text{NumSites} \times 8)
\end{align*}
\]

If you use the minimal columns feature, use this formula to calculate message size:

\[
\begin{align*}
\text{InboundMsgSize update} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + (\text{RowWidthChanged} \times 2) + \\
&\quad \text{PrimaryKeyWidth} \\
\text{OutboundMsgSize update} &= \text{OutboundMsgOverhead} + \\
&\quad (\text{RowWidthChanged} \times 2) + \text{PrimaryKeyWidth} + \\
&\quad (\text{NumSites} \times 8)
\end{align*}
\]

See “Table update calculations” on page 135 for examples of these calculations.

Table inserts

Use the following formula to calculate message size for table inserts. (This formula also applies if you use the minimal columns feature.)

\[
\begin{align*}
\text{InboundMsgSize insert} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + \text{RowWidth} \\
\text{OutboundMsgSize insert} &= \text{OutboundMsgOverhead} + \\
&\quad \text{RowWidth} + (\text{NumSites} \times 8)
\end{align*}
\]

Table deletes

If you do not use minimal columns, use this formula to calculate message size in table deletes:

\[
\begin{align*}
\text{InboundMsgSize insert} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + \text{RowWidth} \\
\text{OutboundMsgSize insert} &= \text{OutboundMsgOverhead} + \\
&\quad \text{RowWidth} + (\text{NumSites} \times 8)
\end{align*}
\]

If you use minimal columns, use this formula to calculate table deletes:

\[
\begin{align*}
\text{InboundMsgSize delete} &= \text{InboundMsgOverhead} + \\
&\quad \text{ColOverhead} + \text{PrimaryKeyWidth} \\
\text{OutboundMsgSize delete} &= \text{OutboundMsgOverhead} + \\
&\quad \text{PrimaryKeyWidth} + (\text{NumSites} \times 8)
\end{align*}
\]
Data volume (queue disk space requirements)

Functions

Use these formulas to calculate message size for functions:

\[
\text{InboundMsgSize}_{\text{function}} = \text{InboundMsgOverhead} + \text{ParameterWidth} + (\text{RowWidth} \times 2)
\]

\[
\text{OutboundMsgSize}_{\text{function}} = \text{OutboundMsgOverhead} + \text{ParameterWidth} + (\text{NumSites} \times 8)
\]

In the formula for inbound message size, \textit{RowWidth} does not apply to the replicated functions feature because before and after images of replicated functions are not sent to the inbound queue.

Begin and commit pairs

Use these formulas to calculate message sizes for begins and commits:

\[
\text{InboundMsgSize}_{\text{begin}} = \text{OutboundMsgSize}_{\text{begin}} = 250
\]

\[
\text{InboundMsgSize}_{\text{commit}} = \text{OutboundMsgSize}_{\text{commit}} = 200
\]

The total size of a begin/commit pair is 450 bytes. If typical transactions have many modifications, omit the begin or commit message sizes from your calculations. Their contribution to overall message size is negligible.

Formula components

In the preceding formulas:

- \textit{InboundMsgOverhead} equals 380 bytes. Each message includes a transaction ID, duplicate detection sequence numbers, and so on.
- \textit{OutboundMsgOverhead} equals 200 bytes. Each message includes a transaction ID, duplicate detection sequence numbers, and so on.
- \textit{ColOverhead}, which applies to the inbound queue only, equals 30 bytes of overhead per column, as follows:
  
  For an update operation without minimal columns:
  \[(\text{NumColumns} + \text{NumColumnsChanged}) \times 30\]
  
  For an update operation with minimal columns:
  
  \[((\text{NumColumnsChanged} \times 2) + \text{NumPrimaryKeyColumns}) \times 30\]
  
  For an insert operation (with or without minimal columns):
  
  \[
  \text{NumColumns} \times 30
  \]
  
  For a delete operation without minimal columns:
  
  \[
  \text{NumColumns} \times 30
  \]
  
  For a delete operation with minimal columns:
  
  \[
  \text{NumPrimaryKeyColumns} \times 30
  \]

- \textit{RowWidth} is the size of the ASCII representation of the table columns. For example: a char(10) column uses 10 bytes, a binary(10) column uses 22 bytes.
For table updates, the RowWidth is multiplied by 2 because both the before and after images of the row are distributed to the replicate sites.

- **RowWidthChanged** is the width of the changed columns in the row. For example, you have 10 columns with a total RowWidth of 200 bytes. Half of the columns in the row change, giving you a RowWidthChanged measurement of approximately 100 bytes.

- **NumSites** is the number of sites that a message will go to. This matters only when small messages are distributed to many sites and may be insignificant if you do not have many sites. If small, you might want to omit the number of sites factor from the formulas that use it.

- **ParameterWidth** is the size of the ASCII representation of the function parameters.

- **Begin/Commit Pair** is the combined size of the begin message header and the commit trailer, which equals 450 bytes.

**Change rate (number of messages)**

The change rate is the maximum number of modifications (inserts, deletes, and updates) made to the table per time period. Change rate varies from table to table.

When planning capacity, always use a change rate greater than the maximum change rate recorded for that table.

**Note** Figures for change rates are expressed as operations per unit of time (for example, 5 updates per second). Always use the same time period when calculating rates. If you measure updates per second, then you must use seconds as the basis for all other time period calculations.

**Change volume (number of bytes)**

The change volume is the amount of data that is changed in the table per time period. This amount varies from table to table. It is a function of the change rate and data size for the table.
Calculating table volume

The volume of data being replicated is the size of each message times the number of messages replicated per second. If you know the message size of each modification to a table and of each replicated function, you can calculate the total volume of messages generated in a database by summing the volumes of the individual tables and replicated functions.

Table volume upper bound method

For purposes of later calculations, figure table volumes and database volumes for inbound and outbound queues separately. Or, calculate the volume for outbound queues and use that figure as an approximation for both inbound and outbound volumes.

The upper bound for each table can be generated by multiplying the update rate on the table by the size of the longest message. This gives you the upper bound of the table volume.

\[
\begin{align*}
\text{InboundTableVolume}_{\text{upper}} &= (\text{Max}[\text{InboundMsgSize}] \times \text{ChangeRate}) \\
\text{OutboundTableVolume}_{\text{upper}} &= (\text{Max}[\text{OutboundMsgSize}] \times \text{ChangeRate})
\end{align*}
\]

where:

- \( \text{Max}[\text{InboundMsgSize}] \) and \( \text{Max}[\text{OutboundMsgSize}] \) are the sizes of the largest inbound and outbound message in bytes (typically, the message size with the largest parameters).
- \( \text{ChangeRate} \) is the maximum number of modifications to the table per time period.

Table volume sum of values method

A more precise calculation of data volume can be achieved by summing all of the volumes for the different types of messages.

To calculate \( \text{InboundTableVolume} \):

\[
\begin{align*}
\text{InboundTableVolume} &= \ (\text{InboundMsgSize}_{\text{update}} \times \text{ChangeRate}_{\text{update}}) + \\
&\quad (\text{InboundMsgSize}_{\text{insert}} \times \text{ChangeRate}_{\text{insert}}) + \\
&\quad (\text{InboundMsgSize}_{\text{delete}} \times \text{ChangeRate}_{\text{delete}}) + \\
&\quad (\text{InboundMsgSize}_{\text{function1}} \times \text{ChangeRate}_{\text{function1}}) + \\
&\quad (\text{InboundMsgSize}_{\text{function2}} \times \text{ChangeRate}_{\text{function2}}) + \\
&\quad \ldots
\end{align*}
\]
To calculate *OutboundTableVolume*:

\[
\text{OutboundTableVolume} = \\
(\text{OutboundMsgSizeupdate} \times \text{ChangeRateupdate}) + \\
(\text{OutboundMsgSizeinsert} \times \text{ChangeRateinsert}) + \\
(\text{OutboundMsgSizedelete} \times \text{ChangeRatedelete}) + \\
(\text{OutboundMsgSizefunction1} \times \text{ChangeRatefunction1}) + \\
(\text{OutboundMsgSizefunction2} \times \text{ChangeRatefunction2}) + \\
\ldots
\]

where:

- *ChangeRate* is the maximum number of modifications to the table per time period. *ChangeRate*_{update} refers to data updates, *ChangeRate*_{insert} refers to data insertions, and so forth.
- *InboundMsgSize* and *OutboundMsgSize* are the different types of maximum inbound and outbound message sizes. (See “Message sizes” on page 124.)

**Transaction volume**

*TransactionVolume* estimates data volume generated by the begin and commit records.

The formula for *InboundTransactionVolume* is:

\[
\text{InboundTransactionVolume} = (\text{MsgSizecommit} + \text{MsgSizebegin}) \times \text{InboundTransactionRate}
\]

where:

- \(\text{MsgSize}_{\text{begin}} + \text{MsgSize}_{\text{commit}}\) equals 450 bytes per transaction.
- *InboundTransactionRate* is the total number of transactions per time period at the primary database. This includes all transactions—those that update replicated tables and those that do not. One transaction may include one or more modifications to one or more tables.

The formula for *OutboundTransactionVolume* is:

\[
\text{OutboundTransactionVolume} = (\text{MsgSizecommit} + \text{MsgSizebegin}) \times \text{OutboundTransactionRate}
\]

where:

- \(\text{MsgSize}_{\text{begin}} + \text{MsgSize}_{\text{commit}}\) equals 450 bytes per transaction.
• **OutboundTransactionRate** is the total number of replicated transactions per time period replicated through a particular outbound queue. Each transaction may include one or more modifications to one or more tables.

  *OutboundTransactionRate* depends on:
  
  • How many transactions actually update the replicated data
  • How many of those transactions are replicated through a particular outbound queue

  For example, if the *InboundTransactionRate* is 50 transactions per second, and half of those transactions update replicated tables, and 20 percent of the updates to the replicated tables are replicated through queue 1, the *OutboundTransactionRate* for queue 1 is 5 transactions per second (50 * 0.5 * 0.2).

  Calculating the *OutboundTransactionRate* can be complicated if transactions contain updates to many different replicated tables with different replication selectivities. In such cases, the *InboundTransactionRate* provides a convenient upper bound for the *OutboundTransactionRate*, and you may use it in your formulas instead of the *OutboundTransactionRate*.

  If all transactions in your database are replicated through an outbound queue, then the *OutboundTransactionRate* is the same as the *InboundTransactionRate*.

### Database volume

A rough estimate of the upper limit of database volume can be calculated by summing the upper bound message rates for each table.

\[
\text{InboundDatabaseVolume}_{\text{upper}} = \text{sum(InboundTableVolume}_{\text{upper}}) + \text{InboundTransactionVolume}
\]

See “Inbound database volume” on page 137 for an example of these calculations.

A more accurate method of determining the *InboundDatabaseVolume* is to sum the Table Volumes that were calculated using the method described in “Table volume upper bound method” on page 128.

\[
\text{InboundDatabaseVolume} = \text{sum(InboundTableVolume)} + \text{InboundTransactionVolume}
\]
Inbound queue size

An inbound queue contains the updates to all of the replicated tables in a primary database. The inbound queue keeps these updates for the duration of the longest open transaction. Queue sizes are expressed in bytes.

To calculate the average size of an inbound queue:

\[
\text{InboundQueueSize}_{\text{typical}} = \text{InboundDatabaseVolume} \times \text{TransactionDuration}_{\text{typical}}
\]

To calculate the maximum size of an inbound queue:

\[
\text{InboundQueueSize}_{\text{longest}} = \text{InboundDatabaseVolume} \times \text{TransactionDuration}_{\text{longest}}
\]

where:

- \(\text{InboundDatabaseVolume}\) is the volume of messages calculated by the formulas described in “Database volume” on page 130.
  
  You can use the \(\text{OutboundDatabaseVolume}\) to calculate the \(\text{InboundQueueSize}\) if you do not need to be precise.

- \(\text{TransactionDuration}_{\text{typical}}\) represents the number of seconds for a typical transaction.

- \(\text{TransactionDuration}_{\text{longest}}\) represents the number of seconds for the absolute transaction.

See “Inbound queue size example calculation” on page 138 for an example of calculating the maximum size of an inbound queue.

In practice, the maximum size of the inbound queue is slightly longer than the calculation, because while a long transaction is being read, new transactions are being added to the log. Since the duration of the longest transaction is an approximation, when determining its value add an estimate for the short period required to read and process it from the stable queue.

Also note that the size of the queue may be larger if messages trickle slowly into the Replication Server. Replication Server writes messages to stable queues in fixed blocks of 16K. To reduce latency, it writes even partially full blocks every second. (You can use the init_sqm_write_delay configuration parameter to change the time delay to something other than 1 second.) If several messages arrive almost simultaneously, all of them will be written into one block. But if the same messages arrive at staggered intervals, each may occupy a block of its own.
Outbound queue volume

An outbound queue sends data to another Replication Server or to a data server. The destination is called a Direct Destination Site in these calculations. For each database replicated to (or through) the Direct Destination Site, some part of the Table Volumes passes through the queue. The fraction of a table replicated through a queue is called replication selectivity.

An upper bound of the OutboundQueueVolume can be calculated by assuming the highest possible replication selectivity of 1 (or 100 percent) for all tables. Then, the OutboundQueueVolume is the sum of all the OutboundTableVolumes for all tables replicated through the queue.

\[
\text{OutboundQueueVolume}_{\text{upper}} = \text{sum}(\text{OutboundTableVolumes}) + \text{OutboundTransactionVolume}
\]

For example, if two tables with OutboundTableVolumes of 20K per second and 10K per second and an assumed OutboundTransactionVolume of 1K per second are replicated through an outbound queue, the OutboundQueueVolume would be:

\[
20\text{K/Sec} + 10\text{K/Sec} + 1\text{K/Sec} = 31\text{K/Sec}
\]

A more accurate measure of queue volume can be obtained by factoring in a value for replication table selectivity. The formula for OutboundQueueVolume using ReplicationSelectivity is:

\[
\text{OutboundQueueVolume} = \text{sum}(\text{OutboundTableVolume} \times \text{ReplicationSelectivity}) + \text{OutboundTransactionVolume}
\]

For example, in the previous example, if only 50 percent of the first table was replicated and 80 percent of the second table was replicated, the OutboundQueueVolume would be:

\[
(20\text{K/Sec} \times 0.5) + (10\text{K/Sec} \times 0.8) + 1\text{K/Sec} = 19\text{K/Sec}
\]

See “Outbound queue volume example calculation” on page 138 for an example of this calculation.

Failure duration

If the direct destination site is unavailable, the queue buffers its messages. The failure duration figure you use to establish queue size determines how long a failure the queue can withstand.
Save interval

The save interval specifies how long data is kept in outbound queues before it is deleted. The save interval helps in recovering from a disk or node loss.

Use the `configure connection` command to specify the save interval for a database connection. Use the `configure route` command to specify the save interval for a route to another Replication Server. See the `Replication Server Reference Manual` for additional information.

Outbound queue size

The size of an outbound queue is calculated by the formula:

\[
\text{OutboundQueueSize} = \text{OutboundQueueVolume} \times (\text{FailureDuration} + \text{SaveInterval})
\]

where:

- \( \text{OutboundQueueVolume} \) is the amount of data in bytes entering the queue per time period.
- \( \text{FailureDuration} \) is the maximum time the queue is expected to buffer data for an unavailable site.
- \( \text{SaveInterval} \) is the time the messages are kept in the queue before being deleted.

If an outbound queue with an \( \text{OutboundQueueVolume} \) of 18K per second has a \( \text{SaveInterval} \) of 1 hour, and is intended to withstand 1 hour of unavailability at the destination site, the size of the queue should be:

\[
\text{OutboundQueueSize} = 18\text{K/sec} \times (60\text{min} + 60\text{min}) = 0.018\text{MB/sec} \times 120\text{min} \times 60\text{sec/min} = 130\text{MB}
\]

See “Outbound queue size” on page 139 for a more detailed example of these calculations.

In practice, if the actual change rates and table selectivities result in less than a full 16K block being filled each second, then the figures calculated above are low. Replication Server writes a block every second, whether or not the block is full. The worst usage of queue space occurs when small transactions are executed more than a second apart. Statistics indicate that most blocks are from 50 percent to 95 percent full.
Data volume (queue disk space requirements)

Overall queue disk usage

To calculate the overall total disk space required for worst-case queue usage, use the formula:

\[ \text{Sum}[\text{InboundQueueSize}] + \text{Sum}[\text{OutboundQueueSize}] \]

where:

- \( \text{Sum}[\text{InboundQueueSize}] \) is the sum of the inbound queue sizes for each of the databases.
- \( \text{Sum}[\text{OutboundQueueSize}] \) is the sum of the outbound queue sizes for each direct destination.

Additional considerations

When planning your replication system, take these considerations into account:

- When one of the remote sites is unavailable because of a network failure, it is likely that the other sites on the same network will be unavailable as well, so many queues will grow simultaneously. Thus, you need to allocate enough disk space for all queues to withstand a simultaneous failure of the planned-for duration.
- When the queues start filling up, you can always add more partitions if you have spare disk space.
- If all the queues fill up and the inbound queue cannot accept more messages, the primary database will not be able to truncate its log. If you manually override this restriction, replicate databases will not receive the truncated updates.

Example queue usage calculations

The following sections present a series of example calculations using the preceding formulas.

Examples, calculation parameters

These calculations use the following assumptions about the example replication system:

- Tables T1 and T2 are in database DB1.
• Table T3 is in database DB2.
• All three tables are replicated to sites S1, S2, and S3.
• The parameters are shown in the following tables (note that not all messages are replicated to all sites).

**Table A-1: Table parameters**

<table>
<thead>
<tr>
<th>Database</th>
<th>Table</th>
<th># Cols</th>
<th>Columns Changed</th>
<th>Change Rate (Num/sec)</th>
<th>Row Width (bytes)</th>
<th>Row Width Changed</th>
<th># Sites</th>
</tr>
</thead>
<tbody>
<tr>
<td>DB1</td>
<td>T1</td>
<td>10</td>
<td>5</td>
<td>20</td>
<td>200</td>
<td>100</td>
<td>3</td>
</tr>
<tr>
<td>DB1</td>
<td>T2</td>
<td>10</td>
<td>5</td>
<td>10</td>
<td>400</td>
<td>200</td>
<td>2</td>
</tr>
<tr>
<td>DB2</td>
<td>T3</td>
<td>120</td>
<td>100</td>
<td>2</td>
<td>1500</td>
<td>1000</td>
<td>2</td>
</tr>
</tbody>
</table>

**Table A-2: Site parameters—table replication selectivity**

<table>
<thead>
<tr>
<th>Table</th>
<th>Percent of Updates to S1</th>
<th>Percent of Updates to S2</th>
<th>Percent of Updates to S3</th>
</tr>
</thead>
<tbody>
<tr>
<td>T1</td>
<td>10%</td>
<td>40%</td>
<td>40%</td>
</tr>
<tr>
<td>T2</td>
<td>40%</td>
<td>80%</td>
<td>0%</td>
</tr>
<tr>
<td>T3</td>
<td>20%</td>
<td>0%</td>
<td>20%</td>
</tr>
</tbody>
</table>

**Table A-3: Transaction rates**

<table>
<thead>
<tr>
<th>Transaction Rate</th>
<th>DB1</th>
<th>DB2</th>
<th>S1</th>
<th>S2</th>
<th>S3</th>
</tr>
</thead>
<tbody>
<tr>
<td>Inbound</td>
<td>20/sec</td>
<td>20/sec</td>
<td>-</td>
<td>-</td>
<td>-</td>
</tr>
<tr>
<td>Outbound</td>
<td>-</td>
<td>-</td>
<td>5/sec</td>
<td>10/sec</td>
<td>8/sec</td>
</tr>
</tbody>
</table>

**Message size example calculations**

These examples use the RowWidthChanged formula (without minimal columns) for the calculation. See “Message sizes” on page 124 for other formulas that can be used to calculate message size.

**Table update calculations**

Using the following formulas:

\[
\text{InboundMsgSize}_{update} = \text{InboundMsgOverhead} + \text{ColOverhead} + \text{RowWidth} + \text{RowWidthChanged} \\
\text{OutboundMsgSize}_{update} = \text{OutboundMsgOverhead} + \text{RowWidth} + \text{RowWidthChanged} + (\text{NumSites} \times 8)
\]

The calculations for updates for each site are:
Data volume (queue disk space requirements)

InboundMsgSizeT1 = 380 + 450 + 200 + 100 = 1130 bytes
InboundMsgSizeT2 = 380 + 450 + 400 + 200 = 1430 bytes
InboundMsgSizeT3 = 380 + 6600 + 1500 + 1000 = 9480 bytes
OutboundMsgSizeT1 = 200 + 200 + 100 + 24 = 500 bytes
OutboundMsgSizeT2 = 200 + 400 + 200 + 16 = 800 bytes
OutboundMsgSizeT3 = 200 + 1500 + 1000 + 16 = 2700 bytes

Begin/commit pairs

MsgSizebegin = 250
MsgSizecommit = 200

Change rate

The change rate is the maximum number of data modifications performed on a table per time period. Naturally, this number is different for every table in your replication system.

These example calculations use the change rates shown in Table A-1 on page 135.

Table volume example calculations

Now calculate the table volumes for tables T1, T2, and T3. For simplicity, what is presented here is a worst-case analysis that assumes all changes are due to SQL update statements plus their associated begin/commit pairs.

Using the upper bound formula for InboundTableVolume:

\[
\text{InboundTableVolume} = (\text{Max[InboundMsgSize]} \times \text{ChangeRate})
\]

The inbound table volumes for tables T1, T2, and T3 are:

InboundTableVolumeT1 = 1130 \times 20 = 23K/sec
InboundTableVolumeT2 = 1430 \times 10 = 14K/sec
InboundTableVolumeT3 = 9480 \times 2 = 19K/sec

Using the upper bound formula for OutboundTableVolume:

\[
\text{OutboundTableVolume} = (\text{Max[OutboundMsgSize]} \times \text{ChangeRate})
\]
The outbound table volumes for tables T1, T2, and T3 are:

\[
\begin{align*}
\text{OutboundTableVolume}_{T1} &= 524 \times 20 = 10K/\text{sec} \\
\text{OutboundTableVolume}_{T2} &= 816 \times 10 = 8K/\text{sec} \\
\text{OutboundTableVolume}_{T3} &= 2716 \times 2 = 5K/\text{sec}
\end{align*}
\]

where:

- \text{InboundMsgSize} and \text{OutboundMsgSize} are taken from the calculations performed in “Table update calculations” on page 135.
- \text{ChangeRate} is taken from Table A-1 on page 135.

### RSSD and log size example calculations

Each Replication Server has an RSSD with its own inbound queue. The RSSD also contributes to the queue volumes of the outbound queues. However, since very little activity is expected in the RSSD, and all transactions are very small, you can assume that the disk space requirements for an RSSD are minimal:

- Inbound Queue = 2MB
- Outbound Queues = nothing

### Inbound database volume

Based on the InboundTableVolumes calculated above, you can calculate the InboundDatabaseVolume using the upper bound formula:

\[
\text{InboundDatabaseVolume} = \text{sum(InboundTableVolume)} + \text{InboundTransactionVolume}
\]

where:

- InboundTableVolumes are taken from the calculations on the previous page.
- TransactionVolume is the sum of the Message_{begin} and Message_{commit} pairs (250 bytes + 200 bytes) times the TransactionRate.
- In these examples, the transaction rates for each database are taken from Table A-3 on page 135.

Thus, the InboundDatabaseVolumes are:

\[
\begin{align*}
\text{InboundDatabaseVolume}_{DB1} &= 23K + 14K + \\
& (450 \text{ bytes/tran} \times 20 \text{ tran/sec}) = 46K/\text{sec} \\
\text{InboundDatabaseVolume}_{DB2} &= 19K +
\end{align*}
\]
Data volume (queue disk space requirements)

(450 bytes/tran * 20 tran/sec) = 28K/sec

Inbound queue size example calculation

Based on the InboundDatabaseVolumes, you can calculate the maximum sizes of inbound queues for DB1 and DB2 using the formula:

\[
\text{InboundQueueSize} = \text{InboundDatabaseVolume} \times \text{TransactionDuration}
\]

where:

- \( \text{Inbound DatabaseVolume} \) is the volume of messages calculated above.
- \( \text{TransactionDuration} \) is the number of seconds for the absolute longest transaction. For the purpose of this example, assume that the \( \text{TransactionDuration} \) for DB1 is 10 minutes, and for DB2 it is 5 minutes.

The InboundQueue sizes are:

\[
\begin{align*}
\text{InboundQueueDB1} &= 46\text{K/sec} \times 600\text{sec} = 28\text{MB} \\
\text{InboundQueueDB2} &= 28\text{K/sec} \times 300\text{sec} = 8\text{MB}
\end{align*}
\]

Outbound queue volume example calculation

Now calculate the OutboundQueueVolume for tables T1, T2, and T3. Because selectivity is measured on a table basis (rather than a database basis), you must calculate outbound queue size for each table replicated through it. The formula for calculating the maximum OutboundQueueVolume is:

\[
\text{OutboundQueueVolume} = \text{sum} (\text{OutboundTableVolume} \times \text{ReplicationSelectivity}) + \text{OutboundTransactionVolume}
\]

where:

- \( \text{OutboundTableVolume} \) is taken from the calculations performed in “Calculating table volume” on page 128.
- \( \text{ReplicationSelectivity} \) values are from “Site parameters—table replication selectivity” on page 135.
- \( \text{TransactionVolume} \) is the size of the Message\text{begin} and Message\text{commit} pairs (250 bytes + 200 bytes) times the \( \text{TransactionRate} \).

The formula for calculating the OutboundQueueVolume for Site 1 is:

\[
\text{OutboundQueueVolumeS1} = \text{OutboundTransactionVolume} + (\text{TableVolumeT1} \times \text{ReplicationSelectivityT1,S1}) +
\]
The OutboundQueueVolumes for the three sites are:

\[
\begin{align*}
\text{Site1} & = (450 \times 5) + (10 \text{K/sec} \times 0.1) + (8 \text{K/sec} \times 0.4) + (5 \text{K/sec} \times 0.2) = 7 \text{K/sec} \\
\text{Site2} & = (450 \times 10) + (10 \text{K/sec} \times 0.4) + (8 \text{K/sec} \times 0.8) + (5 \text{K/sec} \times 0) = 15 \text{K/sec} \\
\text{Site3} & = (450 \times 8) + (10 \text{K/sec} \times 0.4) + (8 \text{K/sec} \times 0) + (5 \text{K/sec} \times 0.2) = 9 \text{K/sec}
\end{align*}
\]

**Outbound queue size**

The size of an outbound queue is calculated by the formula:

\[
\text{OutboundQueueSize} = \text{OutboundQueueVolume} \times (\text{FailureDuration} + \text{SaveInterval})
\]

where:

- **OutboundQueueVolume** is the amount of data in bytes entering the queue per time period.
- **FailureDuration** is the maximum time the queue is expected to buffer data for an unavailable site. For the purpose of these example calculations, assume that failure duration is set at 4 hours (14,400 seconds).
- **SaveInterval** is the time configured for that particular queue. For the purpose of these calculations, assume that the save interval is 2 hours (7200 seconds).

The **OutboundQueueSizes** for the three sites are:

\[
\begin{align*}
\text{Site1} & = 7 \text{K/sec} \times 21,600 \text{sec} = 151 \text{MB} \\
\text{Site2} & = 15 \text{K/sec} \times 21,600 \text{sec} = 324 \text{MB} \\
\text{Site3} & = 9 \text{K/sec} \times 21,600 \text{sec} = 194 \text{MB}
\end{align*}
\]

**Total disk queue usage example calculation**

To calculate the total disk space necessary to handle all queues during a worst-case failure (4 hours), use the formula:

\[
\text{Sum(OutboundQueueSizes)} + \text{Sum(OutboundQueueSizes)}
\]

where:
Other disk space requirements

- **InboundQueueSizes** are the two queues calculated in “Inbound queue size example calculation” on page 138.

- **OutboundQueueSizes** are the three queues calculated above.

The total disk space needed in the worst case, including 2MB for the RSSD inbound queue, is:

\[
2\text{MB} + 28\text{MB} + 8\text{MB} + 151\text{MB} + 324\text{MB} + 194\text{MB} = 707\text{MB}
\]

Sybase recommends that you allocate enough space for your worst case scenario. If you use the save interval feature (outbound queues are not truncated even after messages are delivered to the next site), you must be sure to allocate enough queue space to sustain your peak transaction activity. If you do not use save interval, then under normal circumstances, your queue utilization will be very small, perhaps 1MB or 2MB per queue.

The example calculations assumed that all outbound queues would have to tolerate the same duration of failure. This assumption may not be true in your environment. Typically, connections across a WAN must tolerate longer duration failures than local connections.

Other disk space requirements

This section covers space requirements for Replication Servers, Replication Agents, and data servers in a replication system. LTMs do not require disk space, except for their error log files.

**Stable queues**

The calculations for stable queues are explained in the previous sections of this chapter.

**RSSD**

For the RSSD, allocate at least 12MB for the data and 12MB for the log. These replication system defaults are sufficient for hundreds of replication definitions and thousands of subscriptions.
Errors and rejected transactions are also placed in the RSSD. Tables holding them should be periodically truncated by the administrator. If you dump stable queues to the RSSD to help diagnose problems, be sure to truncate those tables when they are no longer needed.

Logs

Replication Servers and LTM write information, error messages, and trace output into their error log files. You should allocate 1MB to 2MB of disk storage for these log files. If you are asked by Sybase Technical Support to turn on trace flags, you may need to make substantially more log space available.

RepAgent writes messages to the Adaptive Server error log file.

Memory usage

Replication Servers and LTM are separate processes in the replication system. RepAgent is an Adaptive Server thread.

The following sections estimate the approximate memory requirements of each one.

Replication Server memory requirements

As a general guideline, on a Sun SPARCstation, estimate that:

- A newly installed Replication Server uses about 7MB of memory for data and stacks.
- Each DSI connection adds about 500K. This value increases if you configure larger values for MD memory (md_source_memory_pool configuration parameter).
- Every RepAgent (or LTM) connection adds 500K. This value increases if you configure larger values for SQT memory (sqt_max_cache_size configuration parameter).
- If you have thousands of subscriptions, or if you increase the function string cache size, you must account for that by adding more memory.
Memory usage

- Memory for subscription rules is a function of the columns referenced in the subscription and the number of rules. A typical subscription adds less than 80 bytes plus the combined size of all subscription predicate values.
- Function-string cache size can grow to 200K or the value of the fstr_cachesize configuration parameter.
- Remaining system table cache size is based on the number of objects defined. Each replication definition consumes memory equal to approximately 250 times the number of its columns. This might be a factor if you have numerous replication definitions with many columns.

RepAgent memory requirements

Most RepAgent memory comes from allocated Adaptive Server procedure cache (shared memory). It is used for:

- Overhead
- Schema cache
- Transaction cache
- Text and image cache

Refer to the *Sybase Adaptive Server Administration Guide* for information about increasing server memory.

Overhead

Adaptive Server allocates 5612 bytes for each database for log transfer—whether or not RepAgent (or LTM) is enabled.

When RepAgent is enabled, Adaptive Server allocates an additional 2332 bytes of memory at startup for each RepAgent.

Schema cache

The amount of memory used for the schema cache depends on number of objects (tables and stored procedures) and descriptors (columns and parameters) that are replicated. Each object and descriptor requires 128 bytes.

At startup, Adaptive Server allocates 8K of memory—which is sufficient for 64 object/descriptors. Thereafter, memory is allocated in 2K chunks. Adaptive Server also allocates 2048 bytes for the hash table at startup.
A “least recently used” (LRU) mechanism keeps the schema cache size manageable by removing from memory those objects descriptors not recently referenced. Thus, RepAgent does not need enough memory to describe all replicated objects. At minimum, RepAgent needs enough memory to describe one replicated object.

**Transaction cache**

RepAgent requires 256 bytes for each open transaction. Transaction cache memory is allocated in 2K chunks. As transactions are committed or aborted, free memory is placed in a pool which must be used before new memory can be allocated.

RepAgent requires memory for the maximum number of open transactions. If sufficient memory is unavailable, RepAgent shuts down.

At startup, Adaptive Server allocates 2048 bytes for a hash table.

**text and image cache**

The text and image cache is not affected by the size of the text and image data. Rather, the amount of memory used is dependent on the number of replicated tables containing text and image data and the number of columns in the tables that contain text and image data. Each replicated table containing text and image data requires 170 bytes; each replicated column requires 52 bytes.

text and image cache memory is allocated in 2K chunks. Memory is allocated only when replicated tables exist that contain text and image data.

The text and image cache uses a free memory pool and requires sufficient memory for all text and image data.

**Other memory**

- If RepAgent is enabled, Adaptive Server allocates one extra process descriptor for RepAgent.
- RepAgent uses ct-lib to connect to Replication Server. ct-lib allocates memory directly (dynamic memory) as needed.
LTM memory requirements

In addition to the executable program, there are two components to LTM memory requirements:

- The data LTM is maintaining in memory
- Some base minimum requirements for fixed data structures (4MB)

For example, if the largest amount of data that an LTM will need to retain in memory is 1MB, that LTM’s memory requirement is 5MB.

LTM data memory requirement

An LTM keeps in memory the context of every transaction since the oldest open transaction. The oldest open transaction does not have to update any replicated tables. A good rule of thumb to approximate LTM memory usage is:

\[
\text{TransactionRate} \times \text{DurationOfLongestTransaction} \times \text{LTMTransactionOverhead}
\]

where:

- \( \text{TransactionRate} \) is the number of transactions per time period for your replication system.
- \( \text{DurationOfLongestTransaction} \) is based on your replication system’s actual performance.
- \( \text{LTMTransactionOverhead} \) is about 250 bytes.

For example, if the database \( \text{TransactionRate} \) is 10 transactions per second and the \( \text{DurationOfLongestTransaction} \) is 5 minutes, the LTM maximum memory usage would be:

\[
10 \text{ tran/sec} \times 300 \text{ seconds} \times 250 \text{ bytes/tran} = 750K
\]

CPU usage

Replication Server runs on one processor. It can run on an SMP system, but does not make use of additional CPUs.

A Replication Server plays three roles:

- Getting and processing distributions from a primary database via a Replication Agent
• Applying distributions to a replicate database
• Receiving messages from another Replication Server and routing them to other Replication Servers

Each role uses different amounts of CPU time. One Replication Server can play all roles, or the roles can be split among multiple Replication Servers.

Overall performance depends on the number of columns in the tables, the sizes of the columns, the number of subscriptions, the number of sites subscribing, and so forth.

Network requirements

When you have calculated the insert rates into the outbound queues, you can get an idea of the network throughput required. The throughput should allow a Replication Server to send messages faster than messages are being added to the queues. Network messages are usually fractionally larger than what is written to the outbound queues.

Sometimes the queues are filled in bursts and are slowly drained through a potentially low bandwidth network. You should take this behavior into account when you calculate queue sizes and make sure that your queues can handle bursts of incoming data in addition to connection and destination site failures.
Network requirements
This appendix describes the Log Transfer Language (LTL) that a Replication Agent component sends to the primary Replication Server.

### Log Transfer Language overview

Table B-1 lists the LTL commands used by RepAgent (the Replication Agent component for the Adaptive Server Enterprise database).

<table>
<thead>
<tr>
<th>Topic</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>Log Transfer Language overview</td>
<td>147</td>
</tr>
<tr>
<td>connect source</td>
<td>148</td>
</tr>
<tr>
<td>get maintenance user</td>
<td>150</td>
</tr>
<tr>
<td>get truncation</td>
<td>151</td>
</tr>
<tr>
<td>distribute</td>
<td>153</td>
</tr>
<tr>
<td>Sample RepAgent session</td>
<td>165</td>
</tr>
</tbody>
</table>

#### Table B-1: Log Transfer Language commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>connect source</td>
<td>Identifies the Replication Server connection as a log transfer session.</td>
</tr>
<tr>
<td>get maintenance user</td>
<td>Retrieves the login name for the maintenance user.</td>
</tr>
<tr>
<td>get truncation</td>
<td>Retrieves the truncation point for the database from Replication Server.</td>
</tr>
<tr>
<td>distribute</td>
<td>Submits a log record to Replication Server. This command has several subcommands. See “distribute” on page 153.</td>
</tr>
</tbody>
</table>

The first three commands in Table B-1 coordinate the Replication Agent session with the Replication Server. The last command, distribute, submits log records to the Replication Server.

The following sections describe the syntax and use of the LTL commands.
After logging in to Replication Server, RepAgent sends a `connect source` command to identify itself. The command specifies the data server and database log RepAgent is forwarding, and the version of the LTL it proposes to use.

Here is the syntax for the `connect source` command:

```
connect source lti data_server.database version_no
[sendallxacts] [warmstdb] [in recovery]
```

where `data_server` and `database` identify the data server and the database whose log is being forwarded by the RepAgent. The `version_no` parameter identifies the LTL version to be used with the `connect source` command.

RepAgent (or LTM) and Replication Server negotiate the version of LTL to use for the session. This ensures that Replication Agents written for earlier releases of Replication Server can be used with later releases.

The Replication Agent proposes the LTL version to be used in the `version_no` parameter. Replication Server returns the highest supported LTL version that is less than or equal to the proposed version. The maximum LTL version number is 100 for Replication Server version 10.0.x, 101 for Replication Server version 10.1, 102 for Replication Server version 10.1.1, 103 for Replication Server version 11.0, and 200 for Replication Server version 11.5 and later.

- If the Replication Agent uses Replication Server version 11.5 features, it must use LTL version 200 and connect to Replication Server version 11.5 or later.
- If the Replication Agent uses Replication Server 11.0 features, it must use LTL version 103 and must connect to Replication Server version 11.0.
- If the Replication Agent does not use Replication Server 11.0 features, but uses Replication Server 10.1 features, it must use LTL version 101 or 102 and can connect to Replication Server version 10.1 or later.
- If the Replication Agent uses only Replication Server 10.0 features, it can use LTL 100 and can connect to any Replication Server.
Keywords

RepAgent uses different keywords with the connect source command, based on the parameters set for RepAgent. These keywords are:

- **sendallxacts** – when RepAgent is started with the send_maint_xacts_to_replicate flag, it submits all updates on replicated tables, including updates made by the maintenance user, to the Replication Server for distribution to subscribing replicate sites.

- **warmstdb** – when RepAgent is started with the send_warm_standby_xacts flag, it submits all updates on replicated tables, including updates made by the maintenance user, to the Replication Server for application to the standby database.

- **in recovery** – when RepAgent is started with the for_recovery flag it is in recovery mode. When Replication Server is in recovery mode, it permits connections only from RepAgents that are also in recovery mode. Recovery mode is used to replay restored transaction logs so that lost messages can be recovered. For more information about the role of the RepAgent in recovery mode, see the Replication Server Administration Guide.

Upgrade locator

In Replication Server version 11.0 and later, connect source returns an additional row that provides the Replication Server system version number and an upgrade locator.

The upgrade locator gives the origin queue ID of the last message written into the inbound queue before the system was upgraded. For details about the origin queue ID, see “Format of the origin queue ID” on page 152.

The upgrade locator is useful when you are creating a Replication Agent, upgrading from a pre-11.0 version of Replication Server to version 11.0 or later, and mixed-mode transactions (both applied and request functions) are in progress.
Example of connect source

The following `connect source` example identifies the session as a RepAgent session for the `Stocks_db` database in the NY_DS data server. RepAgent and Replication Server agree to use LTL version 200, for version 11.5 of the Replication Server:

```
connect source lti NY_DS.Stocks_db 200
VERSION
         200
Sysversion UpgradeLocator
         1150 0x00000000...
```

get maintenance user

After connecting to Replication Server, RepAgent sends a `get maintenance user` command to find the login name of the maintenance user for the database. Replication Server updates replicated copies of data as the maintenance user. RepAgent uses the login name to distinguish primary data updates from updates distributed through the replication system.

- Changes made by the maintenance user are the result of a distribution, and Replication Server does not redistribute them (except to the warm standby database, if RepAgent is configured to do so using `send_warm_standby_xacts`).
- Changes made to primary data by users other than the maintenance user are primary updates that Replication Server distributes to other databases.

The RepAgent that processes the log of a database containing replicated data must filter out all changes made by the maintenance user.

**Note**  RepAgent for Adaptive Server or LTM for SQL Server can be run in a mode that does not filter out transactions executed by the maintenance user. This feature is used when a consolidated replicated table is replicated to other sites.

Use this syntax for the `get maintenance user` command:

```
get maintenance user for data_server.database
```
Replication Server returns one row with a `char(30)` column that contains the maintenance user login name for the database.

The following example finds the maintenance user for the `pubs2` database in the NY_DS data server:

```sql
get maintenance user for NY_DS.pubs2_db
Maintenance_user
-------------------
pubs2_db_maint
```

**get truncation**

The `get truncation` command returns a 36-byte binary value called the origin queue ID. The format of the origin queue ID is described in the next section.

RepAgent uses the origin queue ID to:

- Locate the last log record saved in the Replication Server inbound queue
- Update the truncation point in the database log

The data server log must be truncated periodically to make room for more log records. RepAgent uses the origin queue ID to update the truncation point and to allow the data server to truncate the log records already received by Replication Server.

Use this syntax for the `get truncation` command:

```sql
get truncation data_server.database
```

where `data_server` and `database` identify the database whose log is being forwarded by the Replication Agent.

Replication Server returns a single row with one 36-byte column that contains the origin queue ID of the last command saved in its inbound queue. The first 32 bytes of the ID are generated by the Replication Agent. The last 4 bytes are appended by Replication Server for its own use and are ignored by the Replication Agent.
Format of the origin queue ID

The origin queue ID is a unique 32-byte binary string that increases in value as each new log record is transferred. The value must increase, because Replication Server ignores records with an ID lower than the highest ID stored in the inbound queue.

When the Replication Server is restarted, it must be able to map the origin queue ID to the original log record so that it can send the next log record with an increased ID value.

Table B-2 shows the format of the origin queue ID generated by the Adaptive Server RepAgent. A Replication Agent for a different data server can generate origin queue IDs in any format as long as the value increases and can be used to find the original log record.

<table>
<thead>
<tr>
<th>Bytes</th>
<th>Contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>1–2</td>
<td>Database generation number used for recovering after reloading coordinated dumps.</td>
</tr>
<tr>
<td>3–8</td>
<td>Log page timestamp for the current record.</td>
</tr>
<tr>
<td>9–14</td>
<td>Row ID of the current row. Row ID = page number (4 bytes) + row number (2 bytes).</td>
</tr>
<tr>
<td>15–20</td>
<td>Row ID of the begin record for the oldest open transaction.</td>
</tr>
<tr>
<td>21–28</td>
<td>Date and time of the begin record for the oldest open transaction.</td>
</tr>
<tr>
<td>29–30</td>
<td>An extension used by the RepAgent to roll back orphan transactions.</td>
</tr>
<tr>
<td>31–32</td>
<td>Unused.</td>
</tr>
</tbody>
</table>

Bytes 21–28 contain an 8-byte datetime datatype value that is the time of the oldest partially transferred transaction in the database log. Replication Server prints the date and time in this field in a message that helps the Database Administrator locate the offline dumps needed for recovery. If you do not store a valid date in this field, the date and time printed in the message are meaningless. However, the message also contains the entire origin queue ID printed in hexadecimal, so if you put the date and time in a location other than bytes 21–28, the Replication System Administrator can extract it from the origin queue ID.


**distribute**

The `distribute` command describes transaction control and data manipulation operations. It also conveys information that Replication Server uses to ensure that transaction log information is transferred without loss or duplication, in the event of failure.

In general, after establishing a session, RepAgent generates a `distribute` command for each operation it retrieves from the database log.

The format of the `distribute` command is:

```
distribute command_tags subcommand [values]
```

The command has three parts:

- The `command_tags` field is used by Replication Server to associate the command with its transaction and to ensure reliable transfer of the database log.
- A `subcommand` name specifies the operation.
- The data `values` associated with the operation are required with all subcommands except `commit transaction` and `rollback transaction`.

**Command tags**

Replication Server normally uses command tags for two purposes:

- To reassemble the commands in a transaction so that data servers at remote sites execute the transaction as a unit
- To ensure that each command is processed only once

The syntax for `command_tags` is:

```
[@origin_time=datetime_value]
@origin_qid=binary_value
@tran_id=binary_value
[@mode=0x08]
[@standby_only={1 | 0}]
```

**origin_time**

The `origin_time` parameter is a `datetime` value that specifies the time when the transaction or data manipulation operation occurred. It is used to report errors. `origin_time` is used only with the transaction control subcommands: `begin transaction`, `commit transaction`, and `rollback transaction`.
**origin_qid**

The `origin_qid` parameter is a 32-byte binary value that uniquely identifies the command in the log. It is a sequence number used by Replication Server to reject duplicate commands after a RepAgent connection has been reestablished. The value is generated as shown in Table B-2 on page 152.

**tran_id**

The `tran_id` parameter is a 120-byte binary value that identifies the transaction the command belongs to. The transaction ID must be globally unique. One way to guarantee this is to first construct a unique transaction ID for the database log, and then attach the data server name and database name to it.

For example, RepAgent constructs a `tran_id` in the following format:

```
tran_id.data_server.database
```

where `tran_id` is a value generated by the RepAgent from information in the log. It contains the generation number, log page timestamp, and the row number of the log record. It is guaranteed to be unique within the database. `data_server` and `database` identify the database whose log is being transferred by the RepAgent.

**mode**

The `mode` parameter is set if the owner name is to be used when Replication Server looks up replication definitions. This parameter is optional for applied commands. It should not be set if the owner name is unavailable.

`mode` is an LTL version 200 parameter; it is available with Replication Server version 11.5 or later.

**standby_only**

The `standby_only` parameter determines whether the command is sent to the standby and/or replicate databases. If `standby_only` is set to 1, the command is sent to the standby database and not to the replicate database. If `standby_only` is set to 0, the command is sent to the standby and replicate databases.

`standby_only` is an LTL version 200 parameter and is available with Replication Server version 11.5 or later. It is optional for applied commands.
Transaction-control subcommands

The transaction-control subcommands are begin transaction, commit transaction, and rollback transaction.

**begin transaction**

The begin transaction subcommand has the following syntax:

```
distribute command_tags begin [system] transaction ['tran_name'] [for 'user'[/'password' | no_password]]
```

- For the syntax and description of `command_tags`, see “Command tags” on page 153.
- The `system` keyword tells Replication Server not to apply this transaction inside begin transaction/commit transaction pairs. In Adaptive Server, it is used for transactions started internally or started in system stored procedures. For LTL version 200 or later, `system` is available with Replication Server version 11.5 or later.
- `tran_name` is an optional `varchar(30)` value that identifies the transaction. The transaction name does not have to be unique. Replication Server makes the transaction name available to function strings in a system-defined parameter.
- `user` and `password` are `varchar(30)` values that identify the login name and password of the user executing the transaction. `user` and `password` should be enclosed in quotes. They are both required for asynchronous procedure calls submitted from a non-primary site.

For LTL version 101 or later, `password` is optional and can be omitted when `user` is supplied. For LTL version 100, `password` must be supplied if `user` is supplied.

Use the `no_password` option when the primary database employs a “unified login” or when the user on the primary database has set a proxy. In both cases RepAgent does not recognize a user password. For LTL version 200 or later, `no_password` is available with Replication Server version 11.5 or later.

**commit transaction, rollback transaction, and rollback**

After RepAgent has submitted all of the commands in a transaction to the Replication Server, it sends either a commit transaction, rollback transaction, or rollback command.
Use this syntax for commit transaction, rollback transaction, and rollback:

```
distribute command_tags
   {commit transaction | rollback transaction | rollback [from ogid to] ogid}
```

Values are not used by commit transaction and rollback transaction.

The rollback subcommand, without the transaction keyword, requires specification of origin queue ID (ogid) values. The three possible forms of this subcommand are:

- **rollback ogid** – rolls back a single log record corresponding to the specified origin queue ID. This option supports the mini-rollback capability in DB2.
- **rollback to ogid** – rolls back all log records between the specified origin queue ID and the current log record.
- **rollback from ogid1 to ogid2** – rolls back a sequence of log records whose origin queue IDs fall in the specified range.

**Note** Replication Server ignores rolled back transactions received from a RepAgent.

**applied subcommand**

The applied subcommand describes operations recorded in the database, including:

- Row updates
- Row inserts
- Row deletes
- Executions of applied stored procedures (request stored procedures use the execute subcommand)
- Modifications to text or image data
The syntax for the applied subcommand is:

\[
\text{distribute command_tags applied [owner=owner_name]}
\]

\[
\{'table'.rs_update
\quad \text{yielding before param_list after param_list} |
\quad 'table'.rs_insert yielding after param_list |
\quad 'table'.rs_delete yielding before param_list |
\quad 'table'.function_name [param_list]
\quad \text{yielding after param_list before param_list} |
\quad 'table'.rs_datarow_for_writetext
\quad \text{yielding datarow column_list} |
\quad 'table'.rs_writetext
\quad \text{append [first] [last] [changed] [with log]}
\quad \text{[textlen=100] column_list}
\]

• For syntax and description of command_tags, see “Command tags” on page 153.

• table is the name of the database table to which the operation was applied. It must be enclosed in quotation marks.

• Replication Server uses table to associate the command with a replication definition. Beginning with Replication Server version 11.5 and version 200 LTL, if the tag \@mode=0x08 is set, Replication Server also associates the owner name with the replication definition. The create replication definition command’s with all tables named table_identifier clause determines how table is mapped to a replication definition:

  • If the replication definition has a with all tables named table_identifier or with primary table named table_identifier clause, table above is matched to the table_identifier or with the primary table named.

  • If the with all tables named table_identifier clause and the with primary table named table_identifier clauses were omitted, then table above is the name of the replication definition.

RepAgent does not need to be aware of replication definitions. It can use the table name on the data source.

yielding clause

For rs_update, rs_insert, and rs_delete, the yielding clause introduces before and after images of the row affected by the operation. Depending on the operation, the before image, the after image, or both, must be provided. Table B-3 shows which operations require before and after images:
The table function_name form of the applied subcommand is used to distribute replicated stored procedures when you use the method associated with table replication definitions. This method is described in Appendix A, “Asynchronous Stored Procedures,” in the Replication Server Administration Guide.

**Note** The preferred method for replicating stored procedures, which uses applied and request functions, is described in Chapter 9, “Managing Replicated Functions,” in the Replication Server Administration Guide. This method uses the execute subcommand to distribute replicated stored procedures (known as replicated functions).

If the stored procedure execution results in an insert or delete operation, RepAgent converts it to an rs_insert or rs_delete LTL command. If the execution results in an update operation, RepAgent uses the function_name form and supplies the before and after images of the updated row to the Replication Server.

A Replication Server function with the same name and parameters as the stored procedure is defined with the create function command, and the function_name in the applied command references this function. The param_list following the function name is the list of parameters of the stored procedure.

The yielding clause contains before and after images of the table row modified by the function. Subscriptions on that table determine where the function is distributed.

Before and after images are specified by a param_list, which is a list of column or parameter values. The syntax for param_list is:

```plaintext
[@param_name=]literal, [@param_name=]literal...
```

- **param_name** is a column name or, for replicated stored procedures, a parameter name.
- **literal** is the value of the column or parameter.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Before Image</th>
<th>After Image</th>
</tr>
</thead>
<tbody>
<tr>
<td>rs_update</td>
<td>Yes</td>
<td>Yes</td>
</tr>
<tr>
<td>rs_insert</td>
<td>—</td>
<td>Yes</td>
</tr>
<tr>
<td>rs_delete</td>
<td>Yes</td>
<td>—</td>
</tr>
</tbody>
</table>
All column names in the replication definition must appear in the list. Replication Server ignores any additional columns. Column or parameter names can be omitted if the values are supplied in the same sequence as they are defined in the replication definition. If the column names are included, you can list them in any order, although there is a performance advantage if the columns are supplied in replication definition order.

Replication Server version 10.1 and later supports an optimized yielding clause. An after image value can be omitted if it is the same as the before image value. For example, if a table has three columns a, b, and c, for an update where only column b changes, the yielding clause could be:

```
yielding before @a=5, @b=10, @c=15 after @b=12
```

**Note** If the minimal columns feature is used, a RepAgent using LTL version 10.1 or later must omit identical after images. See the `create replication definition` command in the *Replication Server Reference Manual* for more information about replicating minimal columns.

---

**Modifications to text or image data**

The `rs_datarow_for_writetext` and `rs_writetext` forms of the applied subcommand are used to distribute modifications to text or image data. These subcommands are built on the Replication Server version 10.1 performance feature of packing data as structured tokens. Each text or image column has a special character and length field, followed by the actual data value. Packing the data in this way eliminates the need for Replication Server to interpret every byte of data, which provides performance benefits.

`rs_datarow_for_writetext` carries an image of the data row associated with a text or non-image column that has been modified by the Transact-SQL `writetext` command, by the Client-Library function, `ct_send_data`, or by the DB-Library™ functions `dbwritetext` and `dbmoretext`. The image is used by Replication Server to construct the primary key for subsequent modification at the replicate database. The syntax for `rs_datarow_for_writetext` is:

```
distribute command_tags applied
'table'.rs_datarow_for_writetext
yielding datarow column_list
```
yielding `datarow column_list` carries the column names, the values of non-text or image columns, and the replication status of text or image columns. The replication status can be `always_rep`, `rep_if_changed`, or `never_rep`. The `column_list` field also carries additional information, called `text_status`, about the text or image columns. The `text_status` can be one of the following keywords:

**Table B-4: Text_status values for text and image data**

<table>
<thead>
<tr>
<th>Keyword</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>tnonnull</td>
<td>The column has a null text pointer. There are no modifications to text or image columns.</td>
</tr>
<tr>
<td>tpininit</td>
<td>Modifications were made at the primary database, which caused a text pointer allocation.</td>
</tr>
<tr>
<td>hastext</td>
<td>The current text or image data value follows.</td>
</tr>
<tr>
<td>notrep</td>
<td>The text or image column is not replicated. No commands are required in the replicate database because the data did not change value and the text or image column has a replicate_if_changed status.</td>
</tr>
<tr>
<td>zerolen</td>
<td>The text or image column contains a null value after an operation at the primary database. For example, after a text pointer has been allocated, there may be data values in a text or image column and an application at the primary database sets them to null.</td>
</tr>
</tbody>
</table>

`rs_insert` and `rs_update` also carry the replication status and additional `text_status` information for text and image columns. `rs_delete` carries the replication status only.

The `rs_writetext` form of the applied subcommand carries the text or image data. `rs_writetext` can carry up to 4K of the text or image data, so the data can be segmented and carried in multiple `rs_writetext` iterations. The syntax for `rs_writetext` is:

```
distribute command_tags applied 'table' rs_writetext
    append [first] [last] [changed] [with log]
    [textlen=100] column_list
```

- `append` indicates that there are more segments of text or image data to follow.
- `first` marks the first segment of data for the text or image column.
- `last` marks the last segment of data for the text or image column.
• changed indicates that the text or image column changed value. If the changed keyword is omitted, the text or image value did not change. This flag is used by the minimal columns feature to discard the data after Replication Server determines it is not needed.

• with log indicates that the modification is logged at the primary database transaction log. It is required only in the first segment of text or image data.

• textlen indicates the total size of the text or image column. It is required only in the first segment of text or image data.

• column_list contains the column name, followed by the text or image data. The data begins with a token header, which is constructed in this order:
  • The tilda (~) character, which denotes a structured token.
  • The period (.) character, if it is text data, or the slash (/) character, if it is image data.
  • The 3-byte character representation of the length of the text or image data segment being carried in an rs_writetext command. The 3-byte length is calculated by converting the length into base-64 representation, then adding the base character ! to each digit to ensure it is a printable character. The formula for calculating the base-64 number is the resultant 3 digits in d3, d2, and d1:
    
    d3 = len/(64^2)
    len = len - (d3*64^2)
    d2 = len/64
    d1 = len%64

    For example, the length of the text segment is 126, and d3= 0, d2=1, and d1=62. The base character ! is added to the digits (the integer value for the base character is 33), and they become !, , and _ (0+33, 1+33, and 2+33). Thus, the structure token header is represented as:

    ~.!."

• Following is an example of the LTL commands generated from the writetext command to update a text column called blurb:
execute subcommand

The execute subcommand is used to send a replicated function or stored procedure call to another Replication Server. This subcommand is used with the preferred method for distributing stored procedures—applied and request functions—and with the older method—request stored procedures.

This is the syntax for the execute subcommand:

```
distribute command_tags execute
   {[repfunc] function | [replication_definition.]function | sys_sp stored_procedure} [param_list]
```
• The repfunc keyword (available only with LTL version 103 or later) indicates that the function name that follows is a user-defined function associated with a function replication definition. When you create a function replication definition for a replicated stored procedure, a user-defined function with the same name is created for you. In this case, the execute subcommand does not include the function replication definition name.

For applied functions, Replication Server distributes the execute repfunc subcommand from a primary Replication Server to any replicate Replication Servers with subscriptions for the associated function replication definition.

For request functions, Replication Server distributes the execute repfunc subcommand from a replicate Replication Server to the primary Replication Server for the function replication definition.

• When the repfunc keyword is omitted, the function name that follows is a user-defined function associated with a table replication definition, and replication_definition is the name of the replication definition. Refer to the Replication Server Administration Guide for a detailed description of user-defined functions.

Without the repfunc keyword, the execute subcommand is used only for request stored procedures associated with table replication definitions. (Applied stored procedures associated with table replication definitions use the applied subcommand.) Replication Server distributes the execute subcommand from a replicate Replication Server to the primary Replication Server for the table replication definition.

If the execute subcommand does not specify a replication definition, Replication Server searches its system tables for the function name and then finds the associated table replication definition. If the function name is not unique, and the replication definition is not specified, an error message reports that the function name is valid for more than one replication definition.

• function is the name of both the user-defined function and the replicated stored procedure. When Replication Server receives the execute command, it maps the function name to a user-defined function previously created by either the create function replication definition command or the create function command.

• With LTL version 200 or later, RepAgent uses sys_sp to send system stored procedures to the standby database.
• *param_list* is a list of the data values supplied when the procedure was executed. You must enclose parameter values in parentheses.

See the *Replication Server Reference Manual* for more information about the `create function replication` definition and `create function` commands. Also see the *Replication Server Administration Guide* for more information about replicated functions and stored procedures.

**Processing the rs_marker function**

When the `rs_marker` stored procedure is executed, RepAgent processes it so that Replication Server synchronizes subscription materialization cycles and warm standby applications.

`rs_marker` is executed with one `varchar(255)` parameter named `@rs_api`. RepAgent passes the parameter to the Replication Server in a `distribute` command with the following syntax:

```
distribute command_tags param_string
```

For example, if a client executes `rs_marker` with the following command:

```
rs_marker @rs_api='0x1234567'
```

The Adaptive Server RepAgent submits the following command to the Replication Server:

```
distribute command_tags 0x1234567
```

Notice that the `@rs_api` parameter is not in quotes.

**sqlddl append subcommand**

The `sqlddl append` subcommand (LTL version 200 or later) is used to apply DDL commands such as `create table` to the warm standby application as original text strings. Because long DDL may span several commands, `sqlddl append` lets you indicate the first and last text strings for the DDL command you want to apply.

`sqlddl append` has the following syntax:

```
sqlddl append [ first | last ] ddl_string
```

• *first* indicates the first part of a DDL sequence.
• *last* indicates the last part of a DDL sequence.
• *ddl_string* is part of a DDL command.
dump subcommand

The dump subcommand is used at a primary site to initiate a coordinated dump. RepAgent retrieves the dump record from the transaction log and sends a dump subcommand to the Replication Server to distribute rs_dumptran or rs_dumpdb functions to all sites with subscriptions for data in the database.

\[
distribute \text{command\_tags} \text{dump [database | transaction] database\_name, dump\_label id}
\]

- \text{database\_name} is the name of the database that was dumped.
- \text{dump\_label} is a varchar(30) value that contains some information to identify the dump. RepAgent uses the date and time of the dump for this variable.
- \text{id} is a varbinary(36) timestamp for the dump.

A dump command must be sent within a transaction.

purge subcommand

The purge subcommand instructs the Replication Server to purge all open transactions in the inbound queue for which the origin queue ID (\text{oqid}) of the begin record is less than that specified in the command.

\[
distribute \text{command\_tags} \text{purge open\_xact to oqid}
\]

- \text{oqid} is the origin queue ID number below which you want to purge all open transactions.

The purge subcommand requires LTL version 102 or higher.

Sample RepAgent session

This section contains a sample dialog between a RepAgent and a Replication Server. The example transfers two transactions to the Replication Server.

The transaction log contains two concurrent transactions, “T1” and “T2”, with log records:

\[
\begin{align*}
T1: & \text{ begin transaction} \\
T2: & \text{ begin transaction} \\
T1: & \text{ insert into authors ('karsen', '510 534-9219')} \\
T2: & \text{ update authors set phone = '510 986-7020'} \\
& \text{ where name = 'green' and phone = '415 986-7020'}
\end{align*}
\]
Sample RepAgent session

T2: commit transaction
T1: commit transaction

**Note**  LTL commands are usually submitted to Replication Server by a RepAgent using Open Client Client-Library routines. However, you can submit LTL commands interactively using `isql`.

distribute
@origin_time='Dec 10 1992  8:48:12:750AM',
@origin_qid=0x00000000000000000000000000000001,
@tran_id=0x000000111111
begin transaction 'T1' for 'user'/'password'
distribute
@origin_time='Dec 10 1992  8:48:12:750AM',
@origin_qid=0x00000000000000000000000000000002,
@tran_id=0x000000222222
begin transaction 'T2' for 'user'/'password'
distribute
@origin_time='Dec 10 1992  8:48:13:750AM',
@origin_qid=0x00000000000000000000000000000003,
@tran_id=0x000000111111
applied 'authors'.rs_insert yielding
  after @name='karsen', @phone='510 534-9219'
distribute
@origin_time='Dec 10 1992  8:48:13:750AM',
@origin_qid=0x00000000000000000000000000000004,
@tran_id=0x000000222222
applied 'authors'.rs_update yielding
  before @name='green', @phone='415 986-7020'
  after @name='green', @phone='510 986-7020'
distribute
@origin_time='Dec 10 1992  8:48:14:750AM',
@origin_qid=0x00000000000000000000000000000005,
@tran_id=0x000000222222
commit transaction
distribute
@origin_time='Dec 10 1992  8:48:14:750AM',
@origin_qid=0x00000000000000000000000000000006,
@tran_id=0x000000111111
commit transaction

You can use the `get truncation` command to verify that the truncation point is set to the `origin_qid` from the last `distribute` command.
### Glossary

<table>
<thead>
<tr>
<th>Term</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>active database</td>
<td>In a warm standby application, a database that is replicated to a standby database. See also <a href="#">warm standby application</a>.</td>
</tr>
<tr>
<td>Adaptive Server</td>
<td>The Sybase version 11.5 and later relational database server. At least one Adaptive Server is required for each Replication Server, to manage the RSSD and any other Adaptive Server databases.</td>
</tr>
<tr>
<td>application programming interface (API)</td>
<td>A predefined interface through which users or programs communicate with each other. Open Client and Open Server are examples of APIs that communicate in a client/server architecture. RCL, the Replication Command Language, is the Replication Server API.</td>
</tr>
<tr>
<td>applied function</td>
<td>A replicated function, associated with a function replication definition, that Replication Server delivers from a primary database to a subscribing replicate database. The function passes parameter values to a stored procedure that is executed at the replicate database. See also <a href="#">replicated function delivery</a>, <a href="#">request function</a>, and <a href="#">function replication definition</a>.</td>
</tr>
<tr>
<td>article</td>
<td>A replication definition extension for tables or stored procedures that can be an element of a publication. Articles may or may not contain <code>where</code> clauses, which specify a subset of rows that the replicate database receives.</td>
</tr>
<tr>
<td>asynchronous procedure delivery</td>
<td>A method of replicating, from a source to a destination database, a stored procedure that is associated with a table replication definition.</td>
</tr>
<tr>
<td>asynchronous command</td>
<td>A command that a client submits where the client is not prevented from proceeding with other operations before the completion status is received. Many Replication Server commands function as asynchronous commands within the replication system.</td>
</tr>
</tbody>
</table>
atomic materialization
A materialization method that copies subscription data from a primary to a replicate database through the network in a single atomic operation, using a `select` operation with a holdlock. No changes to primary data are allowed until data transfer is complete. Replicate data may be applied either as a single transaction or in increments of ten rows per transaction, which ensures that the replicate database transaction log does not fill. Atomic materialization is the default method for the `create subscription` command. See also nonatomic materialization, bulk materialization and no materialization.

autocorrection
Autocorrection is a setting applied to replication definitions, using the `set autocorrection` command, to prevent failures caused by missing or duplicate rows in a copy of a replicated table. When autocorrection is enabled, Replication Server converts each update or insert operation into a delete followed by an insert. Autocorrection should only be enabled for replication definitions whose subscriptions use nonatomic materialization.

base class
A function-string class that does not inherit function strings from a parent class. See also function-string class.

bitmap subscription
A type of subscription that replicates rows based on bitmap comparisons. Create columns using the int datatype, and identify them as the rs_address datatype when you create a replication definition. When you create a subscription, compare each rs_address column to a bitmask using a bitmap comparison operator (&) in the `where` clause. Rows matching the subscription’s bitmap are replicated.

bulk materialization
A materialization method whereby subscription data in a replicate database is initialized outside of the replication system. For example, data may be transferred from a primary database using media such as magnetic tape, diskette, CD-ROM, or optical storage disk. Bulk materialization involves a series of commands, starting with `define subscription`. You can use bulk materialization for subscriptions to table replication definitions or function replication definitions. See also atomic materialization, nonatomic materialization, and no materialization.

centralized database system
A database system where data is managed by a single database management system at a centralized location.

class
See error class and function-string class.

class tree
A set of function-string classes, consisting of two or more levels of derived and parent classes, that derive from the same base class. See also function-string class.
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### Glossary

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<tr>
<th>Term</th>
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<tr>
<td><strong>Data Server Interface (DSI)</strong></td>
<td>Replication Server threads corresponding to a connection between a Replication Server and a database. DSI threads submit transactions from the DSI outbound queue to a replicate data server. They consist of a scheduler thread and one or more executor threads. The scheduler thread groups the transactions by commit order and dispatches them to the executor threads. The executor threads map functions to function strings and execute the transactions in the replicate database. DSI threads use an Open Client connection to a database. See also outbound queue and connection.</td>
</tr>
<tr>
<td><strong>data source</strong></td>
<td>A specific combination of a database management system (DBMS) product such as a relational or non-relational data server, a database residing in that DBMS, and the communications method used to access that DBMS from other parts of a replication system. See also database and data server.</td>
</tr>
<tr>
<td><strong>decision support application</strong></td>
<td>A database client application characterized by ad hoc queries, reports, and calculations and few data update transactions.</td>
</tr>
<tr>
<td><strong>default function string</strong></td>
<td>The function string that is provided by default for the system-provided classes <code>rs_sqlserver_function_class</code> and <code>rs_default_function_class</code> and classes that inherit function strings from these classes, either directly or indirectly. See also function string.</td>
</tr>
<tr>
<td><strong>dematerialization</strong></td>
<td>The optional process, when a subscription is dropped, whereby specific rows that are not used by other subscriptions are removed from the replicate database.</td>
</tr>
<tr>
<td><strong>derived class</strong></td>
<td>A function-string class that inherits function strings from a parent class. See also function-string class and parent class.</td>
</tr>
<tr>
<td><strong>direct route</strong></td>
<td>A route used to send messages directly from a source to a destination Replication Server, with no intermediate Replication Servers. See also indirect route and route.</td>
</tr>
<tr>
<td><strong>disk partition</strong></td>
<td>See partition.</td>
</tr>
<tr>
<td><strong>distributed database system</strong></td>
<td>A database system where data is stored in multiple databases on a network. The databases may be managed by data servers of the same type (for example, Adaptive Server) or by heterogeneous data servers.</td>
</tr>
<tr>
<td><strong>Distributor (DIST)</strong></td>
<td>A Replication Server thread that helps to determine the destination of each transaction in the inbound queue.</td>
</tr>
</tbody>
</table>
**dump marker**
A message written by Adaptive Server in a database transaction log when a dump is performed. In a warm standby application, when you are initializing the standby database with data from the active database, you can specify that Replication Server use the dump marker to determine where in the transaction stream to begin applying transactions in the standby database. See also **warm standby application**.

**error action**
A Replication Server response to a data server error. Possible Replication Server error actions are `ignore`, `warn`, `retry_log`, `log`, `retry_stop`, and `stop_replication`. Error actions are assigned to specific data server errors.

**error class**
A name for a collection of data server error actions that are used with a specified database.

**exceptions log**
A set of three Replication Server system tables that holds information about transactions that failed on a data server. The transactions in the log must be resolved by a user or by an intelligent application. You can use the `rs_helpexception` stored procedure to query the exceptions log.

**fault tolerance**
The ability of a system to continue to operate correctly even though one or more of its component parts is malfunctioning.

**function**
A Replication Server object that represents a data server operation such as insert, delete, select, or begin transaction. Replication Server distributes such operations to other Replication Servers as functions. Each function consists of a function name and a set of data parameters. In order to execute the function in a destination database, Replication Server uses function strings to convert a function to a command or set of commands for a type of database. See also `user-defined function`, and **replicated function delivery**.

**function replication definition**
A description of a replicated function used in replicated function delivery. The function replication definition, maintained by Replication Server, includes information about the parameters to be replicated and the location of the primary version of the affected data. See also **replicated function delivery**.

**function scope**
The range of a function’s effect. Functions have replication definition scope or function-string class scope. A function with replication definition scope is defined for a specific replication definition, and cannot be applied to other replication definitions. A function with function-string class scope is defined once for a function-string class and is available only within that class.
function string
A string that Replication Server uses to map a database command to a data server API. For the `rs_select` and `rs_select_with_lock` functions only, the string contains an input template, used to match function strings with the database command. For all functions, the string also contains an output template, used to format the database command for the destination data server.

function-string class
A named collection of function strings used with a specified database connection. Function-string classes include those provided with Replication Server and those you have created. Function-string classes can share function string definitions through function-string inheritance. The three system-provided function-string classes are `rs_sqlserver_function_class`, `rs_default_function_class`, and `rs_db2_function_class`. See also base class, class tree, derived class, function-string inheritance, and parent class.

function-string inheritance
The ability to share function string definitions between classes, whereby a derived class inherits function strings from a parent class. See also derived class, function-string class, and parent class.

function-string variable
An identifier used in a function string to represent a value that is to be substituted at run time. Variables in function strings are enclosed in question marks (?). They represent column values, function parameters, system-defined variables, or user-defined variables.

function subscription
A subscription to a function replication definition (used in applied function delivery).

generation number
See database generation number.

heterogeneous data servers
Multi-vendor data servers used together in a distributed database system.

hibernation mode
A Replication Server state in which all DDL commands except `admin` and `sysadmin` commands, are rejected; all routes and connections are suspended; most service threads, such as DSI and RSI, are suspended; and RSI and RepAgent (or LTM) users are logged off and not allowed to log on. Used during route upgrades, and may be turned on for a Replication Server to debug problems.

hot standby application
A database application in which the standby database can be placed into service without interrupting client applications and without losing any transactions. See also warm standby application.
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<th>Term</th>
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<tr>
<td>ID Server</td>
<td>One Replication Server in a replication system is the ID Server. In addition to performing the usual Replication Server tasks, the ID Server assigns unique ID numbers to every Replication Server and database in the replication system, and maintains version information for the replication system.</td>
</tr>
<tr>
<td>inbound queue</td>
<td>A stable queue used to spool messages from a Replication Agent to a Replication Server.</td>
</tr>
<tr>
<td>indirect route</td>
<td>A route used to send messages from a source to a destination Replication Server, through one or more intermediate Replication Servers. See also direct route and route.</td>
</tr>
<tr>
<td>interfaces file</td>
<td>A file containing entries that define network access information for server programs in a Sybase client/server architecture. Server programs may include Adaptive Servers, SQL Servers, gateways, Replication Servers, and Replication Agents such as LTM for SQL Server. The interfaces file entries enable clients and servers to connect to each other in a network.</td>
</tr>
<tr>
<td>latency</td>
<td>The measure of the time it takes to distribute to a replicate database a data modification operation first applied in a primary database. The time includes replication agent processing, Replication Server processing, and network overhead.</td>
</tr>
<tr>
<td>local-area network (LAN)</td>
<td>A system of computers and devices, such as printers and terminals, connected by cabling for the purpose of sharing data and devices.</td>
</tr>
<tr>
<td>locator value</td>
<td>The value stored in the <code>rs_locater</code> table of the Replication Server’s RSSD that identifies the latest log transaction record received and acknowledged by the Replication Server from each previous site during replication.</td>
</tr>
<tr>
<td>logical connection</td>
<td>A database connection that Replication Server maps to the connections for the active and standby databases in a warm standby application. See also connection and warm standby application.</td>
</tr>
<tr>
<td>login name</td>
<td>The name that a user or a system component such as Replication Server uses to log in to a data server, Replication Server, or replication agent.</td>
</tr>
<tr>
<td>Log Transfer Language (LTL)</td>
<td>A subset of the Replication Command Language (RCL). A replication agent such as RepAgent or LTM for SQL Server uses LTL commands to submit to Replication Server the information it retrieves from primary database transaction logs.</td>
</tr>
<tr>
<td>Log Transfer Manager (LTM)</td>
<td>The replication agent program for Sybase SQL Server. See also replication agent and RepAgent thread.</td>
</tr>
</tbody>
</table>
maintenance user

A data server login name that Replication Server uses to maintain replicate data. In most applications, maintenance user transactions are not replicated.

materialization

The process of copying data specified by a subscription from a primary database to a replicate database, thereby initializing the replicate table. Replicate data can be transferred over a network, or, for subscriptions involving large amounts of data, loaded initially from media. See also atomic materialization, bulk materialization, no materialization, and nonatomic materialization.

materialization queue

A stable queue used to spool messages related to a subscription being materialized or dematerialized.

missing row

A row missing from a replicated copy of a table but present in the primary table.

mixed-version system

A replication system containing Replication Servers of different software versions that have different capabilities based on their different software versions and site versions. Mixed-version support is available only if the system version is 11.0.2 or greater.

For example, a replication system containing Replication Servers version 11.5 or later and version 11.0.2 is a mixed-version system. A replication system containing Replication Servers of releases earlier than release 11.0.2 is not a mixed-version system, because any newer Replication Servers are restricted by the system version from using certain new features. See also site version and system version.

name space

The scope within which an object name must be unique.

nonatomic materialization

A materialization method that copies subscription data from a primary to a replicate database through the network in a single operation, without a holdlock. Changes to the primary table are allowed during data transfer, which may cause temporary inconsistencies between replicate and primary databases. Data is applied in increments of ten rows per transaction, which ensures that the replicate database transaction log does not fill. Nonatomic materialization is an optional method for the create subscription command. See also autocorrection, atomic materialization, no materialization, and bulk materialization.

network-based security

Secure transmission of data across a network. Replication Server supports third-party security mechanisms that provide user authentication, unified login, and secure message transmission between Replication Servers.
no materialization

A materialization method that lets you create a subscription when the subscription data already exists at the replicate site. Use the `create subscription` command with the `without materialization` clause. You can use this method to create subscriptions to table replication definitions and function replication definitions. See also `atomic materialization` and `bulk materialization`.

online transaction processing (OLTP) application

A database client application characterized by frequent transactions involving data modification (inserts, deletes, and updates).

origin queue ID (qid)

Formed by the RepAgent, the qid uniquely identifies each log record passed to the Replication Server. It includes the date and timestamp and the database generation number. See also `database generation number`.

orphaned row

A row in a replicated copy of a table that does not match an active subscription.

outbound queue

A stable queue used to spool messages. The DSI outbound queue spools messages to a replicate database. The RSI outbound queue spools messages to a replicate Replication Server.

parallel DSI

Configuring a database connection so that transactions are applied to a replicate data server using multiple DSI threads operating in parallel, rather than a single DSI thread. See also `connection` and `Data Server Interface (DSI)`.

parameter

An identifier representing a value that is provided when a procedure executes. Parameter names are prefixed with an `@` character in function strings. When a procedure is called from a function string, Replication Server passes the parameter values, unaltered, to the data server. See also `searchable parameter`.

parent class

A function-string class from which a derived class inherits function strings. See also `function-string class` and `derived class`.

partition

A raw disk partition or operating system file that Replication Server uses for stable queue storage. Only use operating system files in a test environment.

physical connection

See `connection`.

primary data

The definitive version of a set of data in a replication system. The primary data is maintained on a data server that is known to all of the Replication Servers with subscriptions for the data.
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<tr>
<th><strong>term</strong></th>
<th><strong>definition</strong></th>
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<tr>
<td>primary database</td>
<td>Any database that contains data that is replicated to another database via the replication system.</td>
</tr>
<tr>
<td>primary fragment</td>
<td>A horizontal segment of a table that holds the primary version of a set of rows.</td>
</tr>
<tr>
<td>primary key</td>
<td>A set of table columns that uniquely identifies each row.</td>
</tr>
<tr>
<td>primary site</td>
<td>A Replication Server where a function-string class or error class is defined. See error class and function-string class.</td>
</tr>
<tr>
<td>principal user</td>
<td>The user who starts an application. When using network-based security, Replication Server logs in to remote servers as the principal user.</td>
</tr>
<tr>
<td>projection</td>
<td>A vertical slice of a table, representing a subset of the table’s columns.</td>
</tr>
<tr>
<td>publication</td>
<td>A group of articles from the same primary database. A publication lets you collect replication definitions for related tables and/or stored procedures and then subscribe to them as a group. You collect replication definitions as articles in a publication at the source Replication Server and subscribe to them with a publication subscription at the destination Replication Server. See also article and publication subscription.</td>
</tr>
<tr>
<td>publication subscription</td>
<td>A subscription to a publication. See also article and publication.</td>
</tr>
<tr>
<td>query</td>
<td>In a database management system, a query is a request to retrieve data that meets a given set of criteria. The SQL database language includes the select command for queries.</td>
</tr>
<tr>
<td>quiescent</td>
<td>A quiescent replication system is one in which all updates have been propagated to their destinations. Some Replication Server commands or procedures require that you first quiesce the replication system.</td>
</tr>
<tr>
<td>remote procedure call (RPC)</td>
<td>A request to execute a procedure that resides in a remote server. The server that executes the procedure could be a Adaptive Server, a Replication Server, or a server created using Open Server. The request can originate from any of these servers or from a client application. The RPC request format is a part of the Sybase Client/Server Interfaces.</td>
</tr>
<tr>
<td>RepAgent thread</td>
<td>The replication agent for Adaptive Server databases. RepAgent is an Adaptive Server thread; it transfers transaction log information from the primary database to a Replication Server for distribution to other databases.</td>
</tr>
<tr>
<td>replicate database</td>
<td>Any database that contains data that is replicated from another database via the replication system.</td>
</tr>
</tbody>
</table>
replicated function delivery: A method of replicating, from a source to a destination database, a stored procedure that is associated with a function replication definition. See also applied function, request function, and function replication definition.

replicated stored procedure: An Adaptive Server stored procedure that is marked as replicated using the sp_setrepproc or the sp_setreplicate system procedure. Replicated stored procedures can be associated with function replication definitions or table replication definitions. See also replicated function delivery and asynchronous procedure delivery.

replicated table: A table that is maintained by Replication Server, in part or in whole, in databases at multiple locations. There is one primary version of the table, which is marked as replicated using the sp_setreptable or the sp_setreplicate system procedure; all other versions are replicated copies.

replication agent: A program or module that transfers transaction log information representing modifications made to primary data from a database server to a Replication Server for distribution to other databases. RepAgent is the replication agent for Adaptive Server databases. LTM (Log Transfer Manager) is the replication agent for SQL Server databases.

Replication Command Language (RCL): The commands used to manage information in Replication Server.

replication definition: Usually, a description of a table for which subscriptions can be created. The replication definition, maintained by Replication Server, includes information about the columns to be replicated and the location of the primary version of the table.

You can also create function replication definitions; sometimes the term “table replication definition” is used to distinguish between table and function replication definitions. See also function replication definition.

Replication Server: The Sybase server program that maintains replicated data, typically on a LAN, and processes data transactions received from other Replication Servers on the same LAN or on a WAN.

Replication Server Interface (RSI): A thread that logs in to a destination Replication Server and transfers commands from the RSI outbound stable queue to the destination Replication Server. There is one RSI thread for each destination Replication Server that is a recipient of commands from a primary or intermediate Replication Server. See also outbound queue and route.
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<td><strong>Replication System Administrator</strong></td>
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<td><strong>Replication Server System Database (RSSD)</strong></td>
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<td><strong>Replication Server system Adaptive Server</strong></td>
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<td><strong>SQL Server</strong></td>
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schema
The structure of the database. DDL commands and system procedures change system tables stored in the database. Supported DDL commands and system procedures can be replicated to standby databases when you use Replication Server version 11.5 or later and Adaptive Server version 11.5 or later.

searchable column
A column in a replicated table that can be specified in the where clause of a subscription or article to restrict the rows replicated at a site.

searchable parameter
A parameter in a replicated stored procedure that can be specified in the where clause of a subscription to help determine whether or not the stored procedure should be replicated. See also parameter.

secondary truncation point
See truncation point.

site
An installation consisting of, at minimum, a Replication Server, data server, and database, and possibly a Replication Agent, usually at a discrete geographic location. The components at each site are connected over a WAN to those at other sites in a replication system. See also primary site.

site version
The version number for an individual Replication Server. Once the site version has been set to a particular level, the Replication Server enables features specific to that level, and downgrades are not allowed. See also software version, route version, and system version.

software version
The version number of the software release for an individual Replication Server. See also site version and system version.

Stable Queue Manager (SQM)
A thread that manages the stable queues. There is one Stable Queue Manager (SQM) thread for each stable queue accessed by the Replication Server, whether inbound or outbound.

Stable Queue Transaction (SQT) interface
A thread that reassembles transaction commands in commit order. A Stable Queue Transaction (SQT) interface thread reads from inbound stable queues, puts transactions in commit order, then sends them to the Distributor (DIST) thread or a DSI thread, depending on which thread required the SQT ordering of the transaction.

stable queues
Store-and-forward queues where Replication Server stores messages destined for a route or database connection. Messages written into a stable queue remain there until they can be delivered to the destination Replication Server or database. Replication Server builds stable queues using its disk partitions. See also inbound queue, outbound queue, and materialization queue.
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<tr>
<th><strong>standalone mode</strong></th>
<th>A special Replication Server mode used for initiating recovery operations.</th>
</tr>
</thead>
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<tr>
<td><strong>standby database</strong></td>
<td>In a warm standby application, a database that receives data modifications from the active database and serves as a backup of that database. See also warm standby application.</td>
</tr>
<tr>
<td><strong>stored procedure</strong></td>
<td>A collection of SQL statements and optional control-of-flow statements stored under a name in a Adaptive Server database. Stored procedures supplied with Adaptive Server are called system procedures. Some stored procedures for querying the RSSD are included with the Replication Server software.</td>
</tr>
<tr>
<td><strong>subscription</strong></td>
<td>A request for Replication Server to maintain a replicated copy of a table, or a set of rows from a table, in a replicate database at a specified location. You can also subscribe to a function replication definition, for replicating stored procedures.</td>
</tr>
<tr>
<td><strong>subscription dematerialization</strong></td>
<td>See dematerialization.</td>
</tr>
<tr>
<td><strong>subscription materialization</strong></td>
<td>See materialization.</td>
</tr>
<tr>
<td><strong>subscription migration</strong></td>
<td>See row migration.</td>
</tr>
<tr>
<td><strong>Sybase Central</strong></td>
<td>A graphical tool that provides a common interface for managing Sybase and Powersoft products. Replication Server uses Replication Server Manager as a Sybase Central plug-in. See also Replication Server Manager (RSM).</td>
</tr>
<tr>
<td><strong>synchronous command</strong></td>
<td>A command that a client considers complete only after the completion status is received.</td>
</tr>
<tr>
<td><strong>system function</strong></td>
<td>A function that is predefined and part of the Replication Server product. Different system functions coordinate replication activities, such as rs_begin, or perform data manipulation operations, such as rs_insert, rs_delete, and rs_update.</td>
</tr>
<tr>
<td><strong>system-provided classes</strong></td>
<td>Replication Server provides the error class rs_sqlserver_error_class and the function-string class es rs_sqlserver_function_class, rs_default_function_class, and rs_db2_function_class. Function strings are generated automatically for the system-provided function-string classes and for any derived classes that inherit from these classes, directly or indirectly. See also error class and function-string class.</td>
</tr>
<tr>
<td>Term</td>
<td>Definition</td>
</tr>
<tr>
<td>-----------------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>system version</td>
<td>The version number for a replication system that represents the version for which new features are enabled, for Replication Servers of release 11.0.2 or earlier, and below which no Replication Server can be downgraded or installed. For a Replication Server version 11.5, your use of certain new features requires a site version of 1150 and a system version of at least 1102. See also mixed-version system, site version, and software version.</td>
</tr>
<tr>
<td>table replication definition</td>
<td>See replication definition.</td>
</tr>
<tr>
<td>table subscription</td>
<td>A subscription to a table replication definition.</td>
</tr>
<tr>
<td>thread</td>
<td>A process running within Replication Server. Built upon Sybase Open Server, Replication Server has a multi-threaded architecture. Each thread performs a certain function such as managing a user session, receiving messages from a Replication Agent or another Replication Server, or applying messages to a database. See also Data Server Interface (DSI), Distributor (DIST), and Replication Server Interface (RSI).</td>
</tr>
<tr>
<td>transaction</td>
<td>A mechanism for grouping statements so that they are treated as a unit: either all statements in the group are executed or no statements in the group are executed.</td>
</tr>
<tr>
<td>Transact-SQL</td>
<td>The relational database language used with Adaptive Server. It is based on standard SQL (Structured Query Language), with Sybase extensions.</td>
</tr>
</tbody>
</table>
| truncation point            | An Adaptive Server database that holds primary data has an active truncation point, marking the transaction log location where Adaptive Server has completed processing. This is the primary truncation point.

The RepAgent for an Adaptive Server database maintains a secondary truncation point, marking the transaction log location separating the portion of the log successfully submitted to the Replication Server from the portion not yet submitted. The secondary truncation point ensures that each operation enters the replication system before its portion of the log is truncated. |
| user-defined function       | A function that allows you to create custom applications that use Replication Server to distribute replicated functions or asynchronous stored procedures between sites in a replication system. In replicated function delivery, a user-defined function is automatically created by Replication Server when you create a function replication definition. |
| variable                    | See function-string variable.                                                                                                               |
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</thead>
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<tr>
<td>version</td>
<td>See mixed-version system, site version, software version, and system version.</td>
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<tr>
<td>warm standby application</td>
<td>An application that employs Replication Server to maintain a standby database for a database known as the active database. If the active database fails, Replication Server and client applications can switch to the standby database.</td>
</tr>
<tr>
<td>wide-area network (WAN)</td>
<td>A system of local-area networks (LANs) connected together with data communication lines.</td>
</tr>
</tbody>
</table>
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