

Key Polynomials

As we have seen, the Round Two integral basis algorithm involves repeatedly solving $n^2 \times n$ systems of linear equations and so (with the use of up-to-date matrix algorithms) would be expected to take no fewer than $O(n^{1+\log_2 7})$ operations (with $1 + \log_2 7 \approx 3.81$). The Round Four algorithm does considerably better, terminating after

$$O(m^{1+\epsilon}n^3 + m^{2+\epsilon}n^2)$$

operations, where $m = v_p(\text{disc } f)$. The behavior of Round Four is dominated by the cost of computing polynomial resultants, which are required in determining the p -adic values of the various elements that arise.

Another approach is to construct sequences of *valuations*, avoiding the explicit construction of individual elements but working instead with their minimal or characteristic polynomials. Determination of p -adic (and other) values is via Newton polygons, of both elementary and “higher order” types; computation of polynomial resultants is thus avoided.

1. Discrete Valuations of $\mathbb{Q}[x]$

Suppose W is a (non-trivial) discrete valuation of $\mathbb{Q}[x]$. We can approximate W by a sequence V_0, V_1, V_2, \dots , of *inductive valuations* of $\mathbb{Q}[x]$.

We define the valuation V_0 as

$$V_0(a_n x^n + \dots + a_0) = \min \{W(a_n), \dots, W(a_0)\}.$$

Next we let $\phi_1(x) = x$ and $\mu_1 = W(x)$ and define

$$V_1(a_n x^n + \dots + a_0) = \min \{V_0(a_i) + i\mu_1 \mid i = 0, \dots, n\}.$$

For $k > 1$ we assume $W \neq V_{k-1}$. We choose a monic polynomial ϕ_k of minimal degree such that $W(\phi_k) > V_{k-1}(\phi_k)$ and let $\mu_k = W(\phi_k)$.

We define V_k to be the (ϕ_k, μ_k) -*augmentation* of V_{k-1} , denoted

$$(1) \quad V_k = [V_{k-1}, \phi_k \rightarrow \mu_k]$$

and given by

$$(2) \quad V_k(f) = \min \{V_{k-1}(f_i) + i\mu_k \mid i = 0, \dots, n\}$$

for $f(x)$ with ϕ_k -adic expansion

$$(3) \quad f(x) = f_n(x)\phi_k^n + f_{n-1}(x)\phi_k^{n-1} + \dots + f_0(x).$$

Each valuation in the chain V_0, V_1, \dots, V_k is called an *inductive valuation*. If the construction of the successive inductive valuations does not terminate with $W = V_k$ then we define the valuation V_∞ by

$$V_\infty(f) = \lim_{k \rightarrow \infty} V_k(f)$$

for each $f(x)$ in $\mathbb{Q}[x]$.

If $W(f) > V_k(f)$ for all k then $V_{k+1}(f) > V_k(f)$ for all k and, since W is discrete, this implies the limit is ∞ . But this limit is bounded by $W(f)$, hence $W(f) = \infty$ and $f = 0$. Otherwise $W(f) = V_k(f)$ for all $k \geq t$ for some t . Thus $W = V_\infty$.

Theorem (M_1). *Every non-trivial discrete valuation of $\mathbb{Q}[x]$ can be represented either as an inductive valuation or as the limit of an infinite sequence of inductive valuations.*

2. Homogeneous Form

Definitions. A valuation V of $\mathbb{Q}[x]$ induces certain relations on $\mathbb{Q}[x]$.

$$\text{equivalence in } V : a \approx_v b \iff V(b-a) > V(b).$$

$$\text{equivalence-divisibility in } V : a \parallel_v b \iff b \approx_v ca \text{ for some } c(x) \in \mathbb{Q}[x].$$

Theorem (M₁). In the inductive valuation V_k any nonzero polynomial $f(x)$ in $\mathbb{Q}[x]$ has a unique (ϕ_1, \dots, ϕ_k) -adic expansion

$$(4) \quad f(x) = \sum_j c_j p^{m_{0j}} \phi_1^{m_{1j}} \phi_2^{m_{2j}} \dots \phi_k^{m_{kj}}$$

with $c_j \in \mathbb{Q}$, $v_p(c_j) = 0$, and $0 \leq m_{ij} < \deg \phi_{i+1} / \deg \phi_i$ for $i = 1, \dots, k-1$.

The polynomial $f(x)$ is *homogeneous in V_k* if all terms in the expansion (4) have the same value in V_k and each coefficient c_j belongs to $\{1, \dots, p-1\}$.

Each class of polynomials equivalent in V_k contains a unique representative in homogeneous form. The representative of the class of a polynomial $f(x)$ is its *k -homogeneous part*, formed by omitting from the expansion (4) all terms with value greater than $V_k(f)$ and in each remaining term replacing the coefficient c_j by $c_j \bmod p$.

We denote the k -homogeneous part of $f(x)$ by $f^{(V_k)}$. In general,

$$f \approx_{V_k} f^{(V_k)}, \quad V_k(f) = V_k(f^{(V_k)}), \quad f \approx_{V_k} g \text{ if and only if } f^{(V_k)} = g^{(V_k)}.$$

Exercises. Let V_k defined as in section 1. Prove the following.

1. V_k is a (discrete) non-archimedean valuation of $\mathbb{Q}[x]$, i.e.,
 - $V_k(f) = \infty$ if and only if $f = 0$,
 - $V_k(fg) = V_k(f) + V_k(g)$,
 - $V_k(f+g) \geq \min \{V_k(f), V_k(g)\}$.
2. $W(f) \geq V_k(f)$ for all $f(x)$ in $\mathbb{Q}[x]$.
3. If $\deg f < \deg \phi_k$ then $W(f) = V_k(f)$.
4. $W(\phi_i) = V_k(\phi_i)$ for $i = 1, \dots, k$.
5. $W(f) > V_{k-1}(f)$ if and only if $\phi_k \parallel_{V_{k-1}} f$.
6. $V_k(f) > V_{k-1}(f)$ if and only if $\phi_k \parallel_{V_{k-1}} f$.
7. $\phi_k(x) \not\approx_{V_{k-1}} \phi_{k-1}(x)$.
8. If $\phi_k \parallel_{V_{k-1}} f$ and $f \neq 0$ then $\deg f \geq \deg \phi_k$.
9. If $\phi_k \parallel_{V_{k-1}} fg$ then $\phi_k \parallel_{V_{k-1}} f$ or $\phi_k \parallel_{V_{k-1}} g$.

Note: The exercises appear as lemmas and theorems in MacLane's 1936 papers.

BIBLIOGRAPHY

- [M₁] S. MacLane, *A construction for absolute values in polynomial rings*, Transactions of the American Mathematical Society **40**(3) (1936) 363–395.
- [M₂] S. MacLane, *A construction for prime ideals as absolute values of an algebraic field*, Duke Mathematical Journal **2** (1936) 492–510.

3. Key Polynomials and Augmented Valuations

The construction of inductive valuations outlined in the previous section relies on the properties of *key polynomials*.

Definition. A *key polynomial* over a valuation V of $\mathbb{Q}[x]$ is a non-constant monic polynomial $\phi(x)$ in $\mathbb{Z}[x]$ that is *minimal* in V , i.e.,

$$\phi \parallel_v f \text{ and } f \neq 0 \implies \deg f \geq \deg \phi,$$

and *equivalence-irreducible* in V , i.e.,

$$\phi \parallel_v fg \implies \phi \parallel_v f \text{ or } \phi \parallel_v g.$$

Definition. For V a valuation of $\mathbb{Q}[x]$, $\phi(x)$ a non-constant polynomial in $\mathbb{Q}[x]$, and $\mu \in \mathbb{Q}$, the (ϕ, μ) -*augmentation* of V is the map

$$W = [V, \phi \rightarrow \mu]$$

given by

$$W(f) = \min \{ V(f_i) + i\mu \mid i = 0, \dots, n \}$$

for $f(x)$ with ϕ -adic expansion

$$f(x) = f_n(x)\phi^n + f_{n-1}(x)\phi^{n-1} + \dots + f_0(x).$$

Exercises. Let $\phi(x)$ be a key polynomial over V , let $\mu > V(\phi)$, and let W be the (ϕ, μ) -augmentation of V . Prove the following.

1. If $f(x) \neq 0$ then

(i) $V(f \bmod \phi) \geq V(f)$, and

(ii) $V(f \bmod \phi) > V(f)$ if and only if $\phi \parallel_v f$.

2. Let $a(x)$ and $b(x)$ be polynomials with $\deg a < \deg \phi$ and $\deg b < \deg \phi$, and let $a(x)b(x) = q(x)\phi + r(x)$ with $r(x) = a(x)b(x) \bmod \phi$. Then

$$V(q\phi) \geq V(ab) = V(r).$$

3. If $a(x)$ and $b(x)$ are polynomials with $\deg a < \deg \phi$ and $\deg b < \deg \phi$ then

$$W(a\phi^s \cdot b\phi^t) = W(a\phi^s) + W(b\phi^t).$$

4. If the polynomials $f(x)$, $g(x)$, and $f(x)g(x)$ have ϕ -adic expansions

$$f(x) = \sum_j f_j(x)\phi^j, \quad g(x) = \sum_k g_k(x)\phi^k, \quad f(x)g(x) = \sum_m h_m(x)\phi^m$$

respectively, and if s and t are the largest integers such that

$$W(f_s\phi^s) = W(f), \quad W(g_t\phi^t) = W(g)$$

respectively, then

$$W(h_{s+t}\phi^{s+t}) = W(f) + W(g).$$

5. W is a valuation of $\mathbb{Q}[x]$.

Exercise. Let V and W be valuations of $\mathbb{Q}[x]$ such that $W(f) \geq V(f)$ for all $f(x)$ and let $\phi(x)$ be a monic polynomial of minimum degree such that $W(\phi) > V(\phi)$. Show that $\phi(x)$ is a key polynomial over V , as follows.

1. Show that $W(f) > V(f)$ if and only if $\phi \parallel_v f$.

2. Show that $\phi(x)$ is equivalence-irreducible in V .

3. Show that $\phi(x)$ is minimal in V .

4. Non-finite Valuations

Definition. A non-finite valuation of $\mathbb{Q}[x]$ is a map $W : \mathbb{Q}[x] \rightarrow \mathbb{Q} \cup \{\infty\}$ such that

- $W(0) = \infty$,
- $W(fg) = W(f) + W(g)$,
- $W(f + g) \geq \min \{W(f), W(g)\}$

for all $f(x), g(x)$ in $\mathbb{Q}[x]$.

Suppose $G(x)$ is the defining polynomial for an algebraic extension \mathcal{K} of \mathbb{Q} , given by $\mathcal{K} = \mathbb{Q}(\xi)$ for some root ξ of $G(x)$. We are interested in extending the p -adic valuation v_p to \mathcal{K} . Any such extension gives rise to a non-finite valuation W of $\mathbb{Q}[x]$, defined by

$$W(f) = v_p(f(\xi))$$

for $f(x) \in \mathbb{Q}[x]$. The non-finite valuation W can be approximated by a sequence of inductive valuations, in just the same way a discrete valuation of $\mathbb{Q}[x]$ can.

Note that W depends on the choice of ξ , and if $\mu(x)$ is the minimal polynomial of ξ over \mathbb{Q}_p then $W(f) = \infty$ if and only if $\mu(x)$ divides $f(x)$ in $\mathbb{Q}_p[x]$.

Exercise. Assume $1 \leq k \leq n - 1$. Show that

1. $V_n(\phi_k) = V_k(\phi_k)$, and
2. if $\deg f < \deg \phi_{k+1}$ then $V_n(f) = V_k(f)$.

The G -projection of V_k

Suppose $G(x)$ has ϕ_k -adic expansion

$$(5) \quad G(x) = g_m(x) \phi_k^m + g_{m-1}(x) \phi_k^{m-1} + \cdots + g_0(x)$$

and that the expression $V_k(g_i \phi_k^i)$ is minimal for the single value $i = e$. By the exercise and the triangle law, if $n > k$ then

$$V_n(G) = V_n(g_e \phi_k^e) = V_k(g_e \phi_k^e) = V_k(G)$$

and W cannot be the limit of the sequence V_0, V_1, V_2, \dots

Definition. The difference

$$\max \{ i \mid V_k(G) = V_k(g_i \phi_k^i) \} - \min \{ i \mid V_k(G) = V_k(g_i \phi_k^i) \}$$

from the expansion (5) is called the G -projection of V_k .

To approximate W we are constrained to choose only key polynomials ϕ_k and key values μ_k so that each valuation V_k will have positive G -projection.

Definition. V_k is called a k^{th} approximant to G if the G -projection of V_k is positive.

Key Values

The key polynomial ϕ_k having been determined the expansion (5) can be computed and its level k Newton polygon, the lower convex hull of the set

$$\{ (i, V_{k-1}(g_i)) \mid i = 0, \dots, m \},$$

can be drawn.

The G -projection constraint obliges us to choose μ_k so that the lower convex hull of the set

$$\{ (i, V_{k-1}(g_i) + i\mu_k) \mid i = 0, \dots, m \}$$

has a horizontal edge, and this is the case if and only if $-\mu_k$ is the slope of an edge of the level k Newton polygon.

It is also necessary to have $\mu_k > V_{k-1}(\phi_k)$.

Finding ϕ_k

Definition. A polynomial $e(x)$ with ϕ_k -adic expansion

$$e(x) = e_m(x) \phi_k^m + e_{m-1}(x) \phi_k^{m-1} + \cdots + e_0(x)$$

is an equivalence-unit in V_k if $V_k(e_0(x)) < V_k(e_j(x) \phi_k^j)$ for $j = 1, \dots, m$.

Lemma (M₂). The polynomial ϕ_k is a key polynomial over V_k .

Theorem (M₂). In the inductive valuation V_k every polynomial $f(x)$ has a decomposition

$$f(x) \approx_{V_k} e(x) \psi_1(x) \psi_2(x) \cdots \psi_t(x)$$

as a product of homogeneous polynomials, with $e(x)$ an equivalence-unit and each $\psi_i(x)$ a key polynomial, and this decomposition is unique except for the order of the factors.

Lemma (M₂). If V_k is a k^{th} approximant to G then $\phi_k \parallel_{V_{k-1}} G$.

Lemma (M₂). If $G(x)$ is not itself a key polynomial over V_{k-1} then

$$G(x) \approx_{V_{k-1}} e(x) \psi_1(x) \cdots \psi_t(x)$$

with $e(x)$ a homogeneous equivalence-unit and $\psi_1(x), \dots, \psi_t(x)$ homogeneous key polynomials over V_{k-1} .

Lemma (M₂). If ϕ_k is chosen to be one of ψ_1, \dots, ψ_t , but with $\phi_k \neq \phi_{k-1}$, and if the key value μ_k is chosen as described above, then

$$V_k = [V_{k-1}, \phi_k \rightarrow \mu_k]$$

is a k^{th} approximant to G .

5. Residue-classes

Definitions. A valuation V of $\mathbb{Q}[x]$ induces certain relations on $\mathbb{Q}[x]$.

$$\text{congruence in } V : a \equiv_v b \iff V(b - a) > 0.$$

$$\text{congruence-divisibility in } V : a \parallel_v b \iff b \equiv_v ca \text{ for some } c(x) \in \mathbb{Q}[x].$$

Definitions. For a valuation V of $\mathbb{Q}[x]$, the *valuation ring* O_V of V , the prime ideal P_V of O_V , the *residue-class* $\llbracket a \rrbracket_V$ of a polynomial $a(x)$ in O_V , and the *residue-class ring* Δ_V are given by

$$O_V = \{ a(x) \in \mathbb{Q}[x] \mid V(a) \geq 0 \},$$

$$P_V = \{ a(x) \in \mathbb{Q}[x] \mid V(a) > 0 \},$$

$$\llbracket a \rrbracket_V = \{ b(x) \in O_V \mid V(b - a) > 0 \},$$

$$\Delta_V = O_V/P_V = \{ \llbracket a \rrbracket_V \mid a(x) \in O_V \}.$$

Definition. We let Γ_V denote the *value-group* of V , i.e.,

$$\Gamma_V = V(\mathbb{Q}[x]).$$

Definition. For $W = [V, \phi \rightarrow \mu]$ and $f(x)$ a polynomial with $W(f) \in \Gamma_V$, a *W-flattener* of f is a polynomial $f_w^b(x)$ such that

$$V(f_w^b) = W(f_w^b) = -W(f).$$

Proposition (M₁). Let $W = [V, \phi \rightarrow \mu]$, let $f(x)$ be a polynomial with $\phi \not\parallel_w f$, and let $f_w^b(x)$ be an arbitrary *W-flattener* of f .

(i) If $g(x) \in \mathbb{Q}[x]$ with $W(g) = 0$ then

$$f \parallel_w g \iff f_w^b f \parallel_w g.$$

(ii) The polynomial $f(x)$ is *equivalence-irreducible* in W if and only if

$$f_w^b f \parallel_w gh \implies f_w^b f \parallel_w g \text{ or } f_w^b f \parallel_w h$$

for all polynomials $g(x)$ and $h(x)$ with $W(g) = W(h) = 0$.

Definitions. For $W = [V, \phi_w \rightarrow \mu_w]$ we define F_W , τ_W , $\phi_{WW}^{\tau_W b}$, y_W as follows.

○ F_W is the subring of Δ_w given by

$$F_W = \{ \llbracket f \rrbracket_w \mid V(f) \geq 0 \} = \{ \llbracket f \rrbracket_w \mid f \in O_V \}.$$

○ τ_w denotes the smallest positive integer such that $\tau_w \mu_w \in \Gamma_V$.

○ $\phi_{WW}^{\tau_w b}(x)$ denotes an arbitrary *W-flattener* of $\phi_w^{\tau_w b}$.

○ y_w denotes the residue-class $\llbracket \phi_{WW}^{\tau_w b} \phi_w^{\tau_w} \rrbracket_w$.

Lemma (M₁). y_w is transcendental over F_W and $\Delta_w = F_W[y_w]$.

Lemma (M₁). If V is the (ϕ, μ) -augmentation of the valuation U and $\psi(x)$ is a key polynomial over V not equivalent in V to $\phi(x)$ then $V(\psi) \in \Gamma_U$.

Theorem (M₁). Let V be the (ϕ_V, μ_V) -augmentation of the valuation U , let W be the (ϕ_w, μ_w) -augmentation of V , with $\phi_w \not\approx_V \phi_V$, and let $\phi_{wV}^b(x)$ be an arbitrary *V-flattener* of ϕ_w . Then the following hold.

(i) The polynomial $\psi_w(y_w) = \llbracket \phi_{wV}^b \phi_w \rrbracket_V$ is irreducible in $F_V[y_w]$.

(ii) If θ_w is a root of ψ_w then $F_w = F_V(\theta_w)$.

(iii) If $m = \deg \psi_w$ then $\deg \phi_w = m \tau_V \deg \phi_V$.