

***COEN 212:***  
***DIGITAL SYSTEMS DESIGN I***  
***Lecture 7: Common Combinational Logic Circuits***  
***(Adders and Multipliers)***

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# Lecture 7:

## Objectives of this lecture

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- **In this lecture, we talk about:**
  - Adders.
  - Subtractors.
  - Multipliers and,
  - Magnitude Comparators.
- **In the next lecture, we will talk about:**
  - Decoders, Encoders, Multiplexers

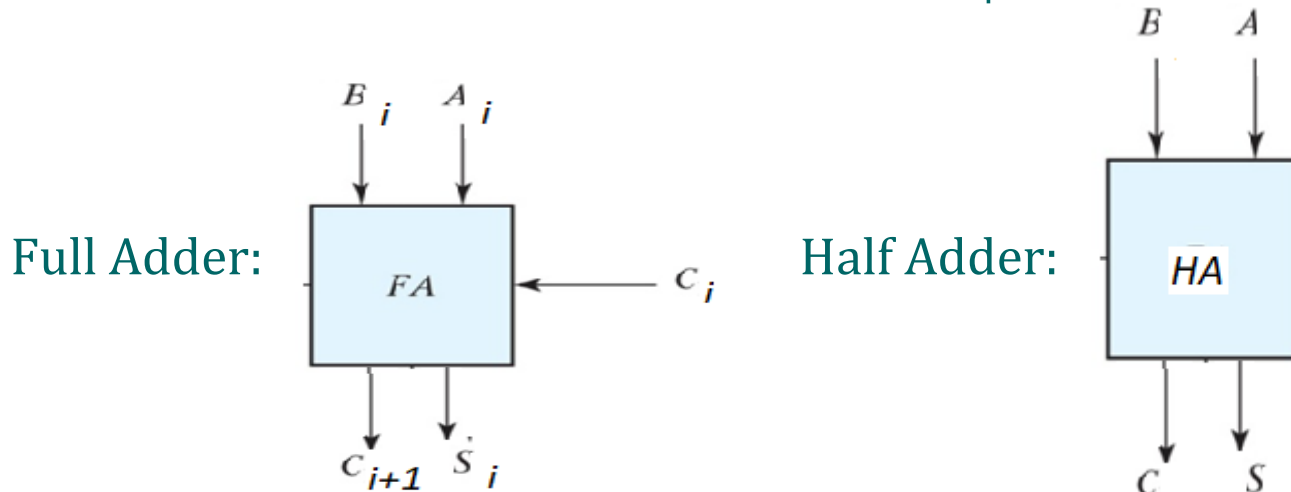
# Lecture 7: Reading for this lecture

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- **Digital Design by M. Morris R. Mano and Michael D. Ciletti, 6th Edition, Pearson, 2018:**
  - **Chapter 4 (4.5 to 4.8)**

# Lecture 7: Binary Adders

- Binary numbers are added bit by bit.
- To add two bits, we need a Full Adder.
- It takes in two bits and a carry and outputs a sum and a carry:
- A full adder can be implemented either directly or using two half adder. A half adder has two inputs and two outputs.



# Lecture 7: Addition: HA

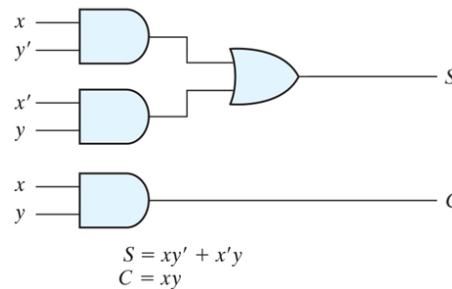
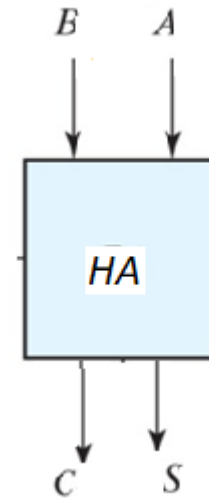
- Truth table for a half adder is:

| A | B | C | S |
|---|---|---|---|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 0 |

- So:

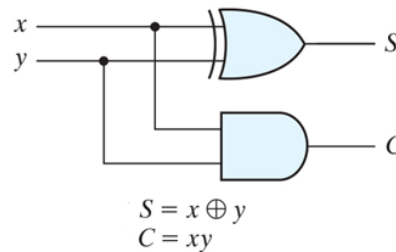
$$S = A'B + AB' = A \oplus B$$
$$C = AB$$

- The implementation is:

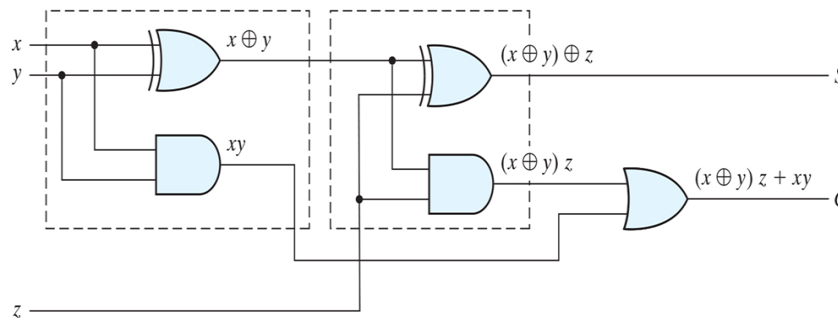


# Lecture 7: Addition: HA

- Half adder can also be implemented using XOR:



- Implementing a FA using two HA's:



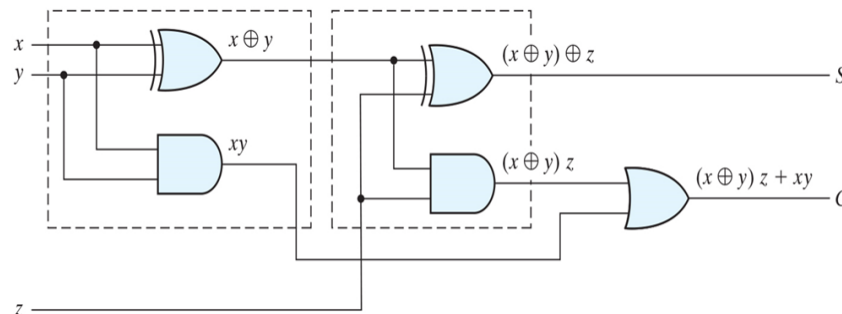
# Lecture 7:

## Addition: FA implementation using HA's

- Half adder can also be implemented using XOR:

$$S_i = A_i \oplus B_i \oplus C_i$$

- Also, :  $C_{i+1} = A_i B_i + A_i C_i + B_i C_i$   
 $= A_i(B_i + C_i) + B_i(A_i + C_i)$   
 $= A_i(B_i + B_i' C_i) + B_i(A_i + A_i' C_i)$   
 $= C_i(A_i B_i' + A_i' B_i) + A_i B_i$   
 $= C_i(A_i \oplus B_i) + A_i B_i$



# Lecture 7:

## Addition: FA direct implementation

- Truth Table of FA:

| x | y | z | C | S |
|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 |

K-map of  $S_y$

| x \ yz |   | y          |            |            |            |
|--------|---|------------|------------|------------|------------|
|        |   | 00         | 01         | 11         | 10         |
| x      | 0 | $m_0$      | $m_1$<br>1 | $m_3$      | $m_2$<br>1 |
|        | 1 | $m_4$<br>1 | $m_5$      | $m_7$<br>1 | $m_6$      |

$$S = x'y'z + x'yz' + xy'z' + xyz$$

$$S = x'y'z + x'yz' + xy'z' + xyz$$

K-map of C

| x \ yz |   | y     |            |            |            |
|--------|---|-------|------------|------------|------------|
|        |   | 00    | 01         | 11         | 10         |
| x      | 0 | $m_0$ | $m_1$      | $m_3$<br>1 | $m_2$      |
|        | 1 | $m_4$ | $m_5$<br>1 | $m_7$<br>1 | $m_6$<br>1 |

$$C = xy + xz + yz$$

$$C = xy + xz + yz$$

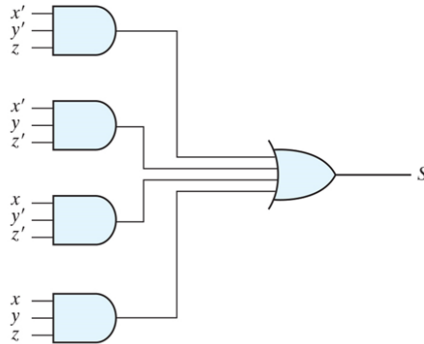


# Lecture 7:

## Addition: FA direct implementation

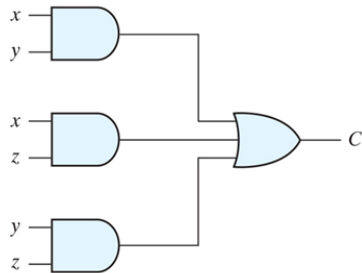
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- The implementation for  $S$



$$S = x'y'z + x'yz' + xy'z' + xyz$$

- The implementation for  $C$

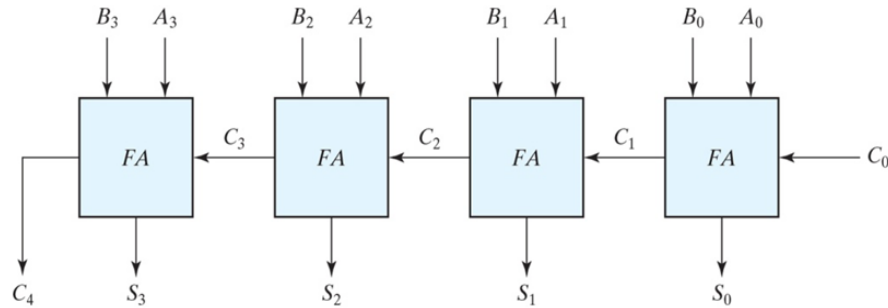


$$C = xy + xz + yz$$

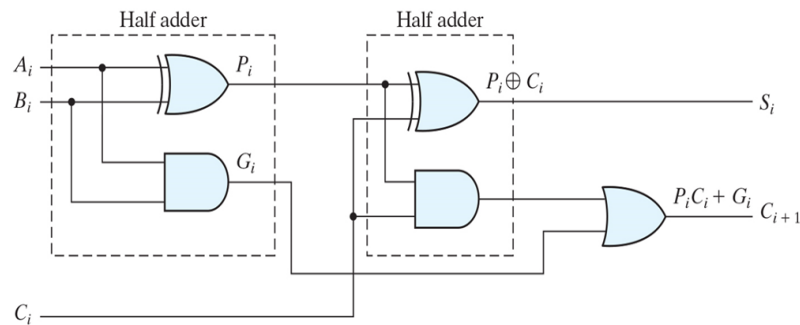
# Lecture 7:

## Binary adders/carry propagation

- Four-bit adder: Delay  $2m = 2 \times 4 = 8$



- Carry Lookahead



# Lecture 7:

## Binary adders/carry Lookahead

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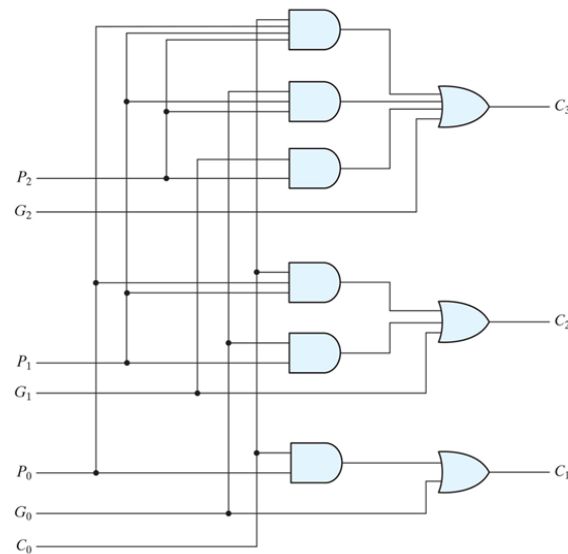
- Note that  $G_i$  generates a carry when both the inputs  $A_i$  and  $B_i$  are 1 regardless of the value of  $C_i$ .
- $G_i$  is called a carry generator.
- $P_i$  decides whether a carry will propagate from stage  $i$  to stage  $i + 1$ . It is called a carry propagator.
- Note that  $P_i$  and  $G_i$  are generated from  $A_i$  and  $B_i$  in one gate delay.
- Now, let's consider the example of 4-bit adder.
- $C_0 =$  the input carry
- $C_1 = G_0 + P_0C_0$
- $C_2 = G_1 + P_1C_1 = G_1 + P_1(G_0 + P_0C_0) = G_1 + P_1G_0 + P_1P_0C_0$
- $C_3 = G_2 + P_2C_2 = G_2 + P_2G_1 + P_2P_1G_0 + P_2P_1P_0C_0$

# Lecture 7:

## Binary adders/carry Lookahead

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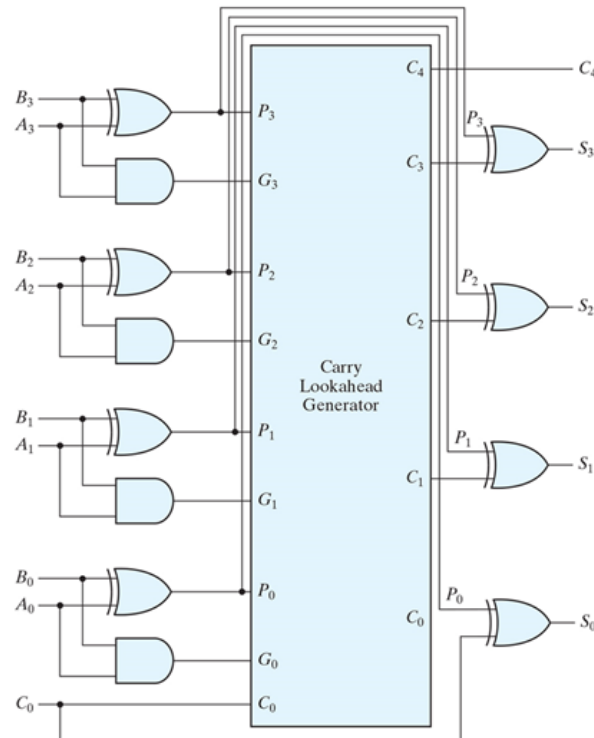
- carry lookahead generator:



# Lecture 7:

## Binary adders/carry Lookahead

- carry lookahead generator:





# Lecture 7:

## Subtraction: Example

- Example: 8-bit adder.
- Numbers are from -127 to +127 (from 11111111 to 01111111).
- The result may fall out of range.
- Example: Adding +60 and +70 = 130 > 127

$$\begin{array}{r} +60 \qquad 00111100 \\ +70 \qquad 01000110 \\ \hline +130 \qquad 10000010 \quad \rightarrow \quad -2 \end{array}$$

- Subtraction +70 from -60

$$\begin{array}{r} -60 \qquad 11000100 \\ -70 \qquad 10111010 \\ \hline -130 \qquad 10111110 \quad \rightarrow \quad +1 \end{array}$$

sign bit

overflow

# Lecture 7: BCD Adder

- Adding two decimal digits and a carry, we get a number between 0 and 19.

| K | Binary Sum     |                |                |                | C | BCD Sum        |                |                |                | Decimal |
|---|----------------|----------------|----------------|----------------|---|----------------|----------------|----------------|----------------|---------|
|   | Z <sub>8</sub> | Z <sub>4</sub> | Z <sub>2</sub> | Z <sub>1</sub> |   | S <sub>8</sub> | S <sub>4</sub> | S <sub>2</sub> | S <sub>1</sub> |         |
| 0 | 0              | 0              | 0              | 0              | 0 | 0              | 0              | 0              | 0              | 0       |
| 0 | 0              | 0              | 0              | 1              | 0 | 0              | 0              | 0              | 1              | 1       |
| 0 | 0              | 0              | 1              | 0              | 0 | 0              | 0              | 1              | 0              | 2       |
| 0 | 0              | 0              | 1              | 1              | 0 | 0              | 0              | 1              | 1              | 3       |
| 0 | 0              | 1              | 0              | 0              | 0 | 0              | 1              | 0              | 0              | 4       |
| 0 | 0              | 1              | 0              | 1              | 0 | 0              | 1              | 0              | 1              | 5       |
| 0 | 0              | 1              | 1              | 0              | 0 | 0              | 1              | 1              | 0              | 6       |
| 0 | 0              | 1              | 1              | 1              | 0 | 0              | 1              | 1              | 1              | 7       |
| 0 | 1              | 0              | 0              | 0              | 0 | 1              | 0              | 0              | 0              | 8       |
| 0 | 1              | 0              | 0              | 1              | 0 | 1              | 0              | 0              | 1              | 9       |
| 0 | 1              | 0              | 1              | 0              | 1 | 0              | 0              | 0              | 0              | 10      |
| 0 | 1              | 0              | 1              | 1              | 1 | 0              | 0              | 0              | 1              | 11      |
| 0 | 1              | 1              | 0              | 0              | 1 | 0              | 0              | 1              | 0              | 12      |
| 0 | 1              | 1              | 0              | 1              | 1 | 0              | 0              | 1              | 1              | 13      |
| 0 | 1              | 1              | 1              | 0              | 1 | 0              | 1              | 0              | 0              | 14      |
| 0 | 1              | 1              | 1              | 1              | 1 | 0              | 1              | 0              | 1              | 15      |
| 1 | 0              | 0              | 0              | 0              | 1 | 0              | 1              | 1              | 0              | 16      |
| 1 | 0              | 0              | 0              | 1              | 1 | 0              | 1              | 1              | 1              | 17      |
| 1 | 0              | 0              | 1              | 0              | 1 | 1              | 0              | 0              | 0              | 18      |
| 1 | 0              | 0              | 1              | 1              | 1 | 1              | 0              | 0              | 1              | 19      |

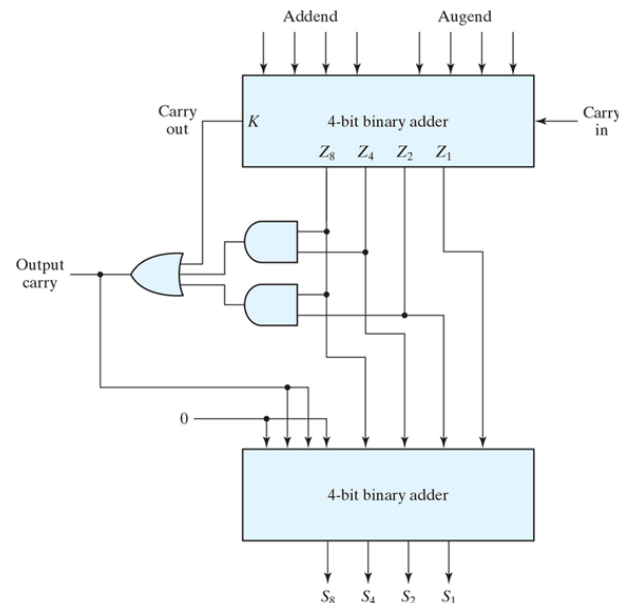


# Lecture 7: BCD Adder

- $C = 1$  whenever  $K = 1$  or  $Z_8 = 1$ .
- $C = 0$  when  $Z_8 = 1$  and both  $Z_4$  and  $Z_2$  are zero.
- So,

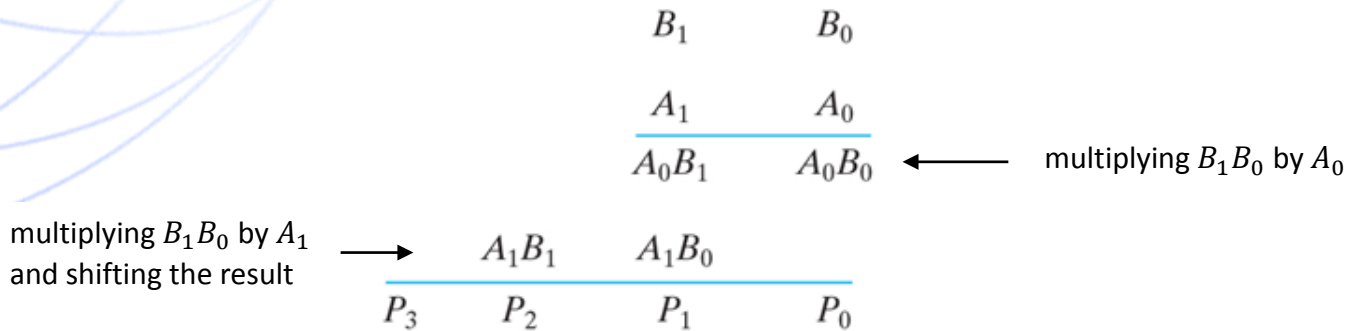
$$C = K + Z_8(Z_2 + Z_4) = K + Z_8Z_4 + Z_8Z_2$$

- When  $C = 1$ , we need to add 6 (0110) to  $Z_8Z_4Z_2Z_1$  to get  $S_8S_4S_2S_1$ .

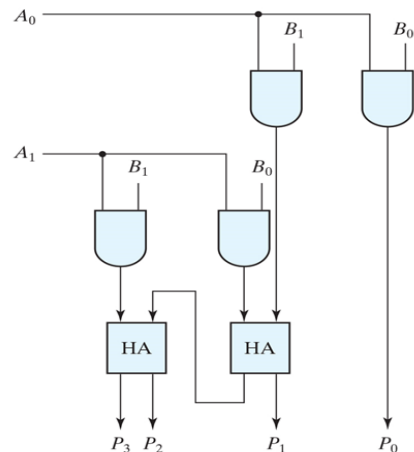


# Lecture 7: Binary multiplication

- $C = 1$  whenever  $K = 1$  or  $Z_8 = 1$ .



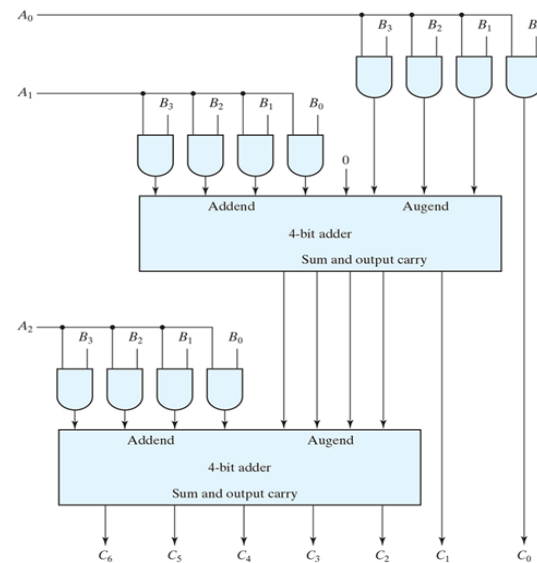
- 2-by-2 multiplier:



# Lecture 7: Binary multiplication

- Example: 4-bit by 3-bit multiplier

|       |          |          |          |          |          |       |  |
|-------|----------|----------|----------|----------|----------|-------|--|
|       |          | $B_3$    | $B_2$    | $B_1$    | $B_0$    |       |  |
|       |          | $A_2$    | $A_1$    | $A_0$    |          |       |  |
|       |          |          |          |          |          |       |  |
|       |          | $A_0B_3$ | $A_0B_2$ | $A_0B_1$ | $A_0B_0$ |       |  |
|       | $A_1B_3$ | $A_1B_2$ | $A_1B_1$ | $A_1B_0$ |          |       |  |
|       | $A_2B_3$ | $A_2B_2$ | $A_2B_1$ | $A_2B_0$ |          |       |  |
|       |          |          |          |          |          |       |  |
| $C_6$ | $C_5$    | $C_4$    | $C_3$    | $C_2$    | $C_1$    | $C_0$ |  |



# Lecture 7:

## Magnitude Comparator

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- It compares two numbers  $A = A_3A_2A_1A_0$  and  $B = B_3B_2B_1B_0$  and decides whether  $A = B$ ,  $A > B$  or  $A < B$ .
- $A = B$  if and only if  $A_3 = B_3, A_2 = B_2, A_1 = B_1, A_0 = B_0$ .
- $A_3 = B_3$  if  $X_3 = A_3B_3 + A_3'B_3'$ .
- Similarly  $X_i = A_iB_i + A_i'B_i'$ ,  $i = 0, 1, 2, 3$  shows  $A_i = B_i$ .
- So:  $(A = B) = X_3X_2X_1X_0$ .
- Case of:  $A > B: A > B$
- If  $A_3$  is equal to 1 and  $B_3 = 0$ . So, if  $A_3B_3' = 1$ , then  $A > B$ .
- If  $A_3 = B_3$ , i.e., if  $X_3 = 1$  and  $A_2 = 1$  and  $B_2 = 0$ , i.e., if  $X_3A_2B_2' = 1$ .
- if  $X_3X_2A_1B_1' = 1$  or  $X_3X_2X_1A_0B_0' = 1$ .

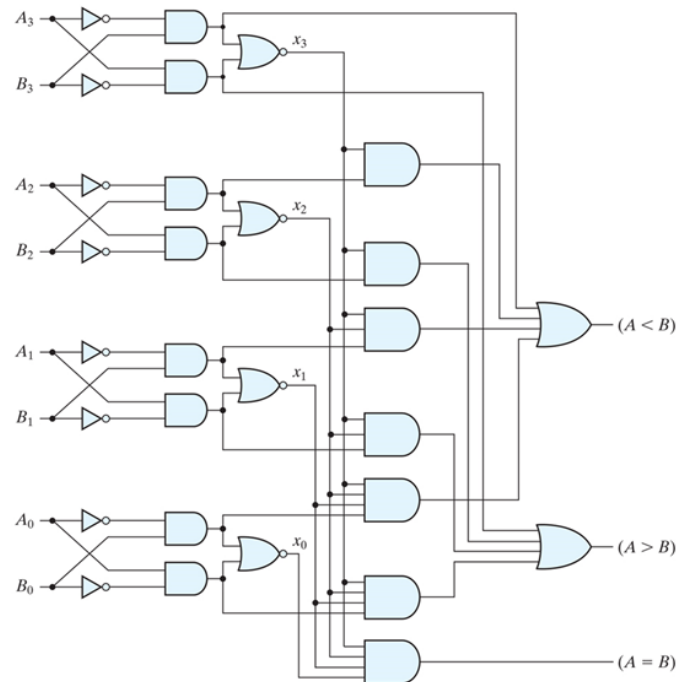
$$(A > B) = A_3B_3' + X_3A_2B_2' + X_3X_2A_1B_1' + X_3X_2X_1A_0B_0'$$

and

$$(A < B) = A_3'B_3 + X_3A_2'B_2 + X_3X_2A_1'B_1 + X_3X_2X_1A_0'B_0$$

# Lecture 7: Magnitude Comparator

- 4-bit Magnitude Comparator:



# Lecture 7: Knowledge Check

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- **Question 1:** To add two 8-bit numbers, the number of Half-Adders needed is:
  - a) 8, b) 15, c) 16, d) 17
- **Question 2:** If the gate delay is 10 ns (nano seconds), what would be the delay in adding two 8 bit numbers using carry look ahead?
  - a) 40 ns, b) 160 ns, c) 80 ns, d) 20 ns
- **Question 3:**