

This can be written as

$$\frac{\sigma^2}{D} = 2^{+2R}$$

where $\frac{\sigma^2}{D}$ is the Signal-to-Quantization-Noise Ratio. Changing into decibel, we get

$$SQNR = 10 \log \frac{\sigma^2}{D} = 10 \log 2^{+2R} = +20R \log 2$$

$$SQNR = +6R \text{ dB}$$

This means that for each 1 bit increase in resolution, we gain 6 dB (4 times improvement) in SQNR.

X 11th lecture, Nov. 11, 2003

Simultaneous Quantization (Description) of Independent Gaussian Random variables:

Take normal sources X_1, X_2, \dots, X_m where $X_i \sim N(0, \sigma_i^2)$. We would like to find

$$R(D) = \min_{\{\hat{x}^m | x^m\}: E d(x^m, \hat{x}^m) \leq D} I(X^m; \hat{x}^m)$$

We have:

$$\begin{aligned} I(X^m; \hat{x}^m) &= h(X^m) - h(X^m | \hat{x}^m) \\ &= \sum_{i=1}^m h(X_i) - \sum_{i=1}^m h(X_i | X_{i-1}, \hat{x}^m) \\ (a) \quad &\geq \sum_{i=1}^m h(X_i) - \sum_{i=1}^m h(X_i | \hat{x}_i) \end{aligned}$$

$$= \sum_{i=1}^m \bar{I}(x_i; \hat{x}_i)$$

$$(b) \geq \sum_{i=1}^m R(D_i) = \sum_{i=1}^m \left(\frac{1}{2} \log \frac{\sigma_i^2}{D_i} \right)^+$$

We can achieve equality in (a) by letting

$$f(x^m | \hat{x}^m) = \prod_{i=1}^m f(x_i | \hat{x}_i)$$

and we get equality in (b) by letting

$$\hat{x}_i \sim N(0, \sigma_i^2 - D_i).$$

So, the problem is to minimize

$$\sum_{i=1}^m \max \left\{ \frac{1}{2} \log \frac{\sigma_i^2}{D_i}, 0 \right\}$$

w.r.t. the constraint $\sum_{i=1}^m D_i = D$

let

$$J(D) = \sum_{i=1}^m \frac{1}{2} \log \frac{\sigma_i^2}{D_i} + \mu \sum_{i=1}^m D_i$$

$$\frac{\partial J(D)}{\partial D_i} = -\frac{1}{2} \frac{1}{D_i} \times \frac{1}{\ln 2} + \mu = 0$$

$$\text{or } D_i = \Theta \quad \text{where } \Theta = (2\mu \ln 2)^{-1}$$

select Θ to satisfy

we use the Kuhn-Tucker condition: to get

$$\frac{\partial J}{\partial D_i} = \begin{cases} 0 & \text{if } D_i < \sigma_i^2 \\ < 0 & \text{if } D_i \geq \sigma_i^2 \end{cases}$$

$D_i = \theta$ means that we allocate equal distortion to each source, i.e., we assign the number of bits such that the distortion is the same. But, this is possible if θ is less than all σ_i^2 's (since $\frac{1}{2} \log \frac{\sigma_i^2}{D_i} = \frac{1}{2} \log \frac{\sigma_i^2}{\theta}$ will be $-\infty$ if $\sigma_i^2 < \theta$ for some i). When θ is greater than some of the σ_i^2 's we set the distortion for those i 's equal to σ_i^2 , i.e., we don't encode these and let the decoder use its best guess.

Therefore,

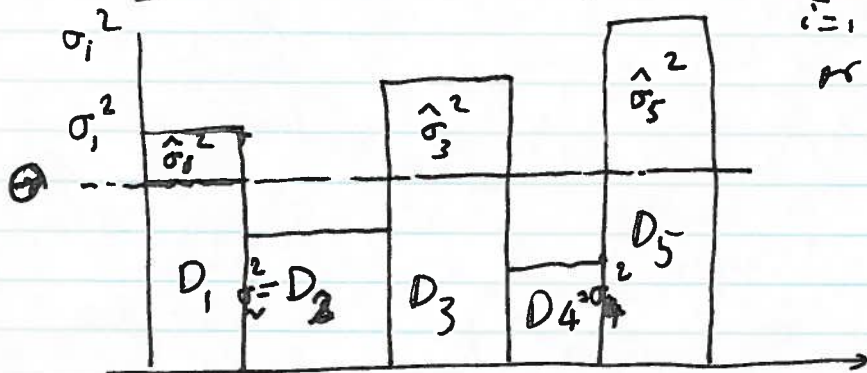
$$R(D) = \sum_{i=1}^m \frac{1}{2} \log \frac{\sigma_i^2}{D_i} = \sum_{i=1}^m \left(\frac{1}{2} \log \frac{\sigma_i^2}{\min(\sigma_i^2, \theta)} \right)^+$$

where

$$D_i = \begin{cases} \theta & \text{if } \theta < \sigma_i^2 \\ \sigma_i^2 & \text{if } \theta \geq \sigma_i^2 \end{cases} = \min(\sigma_i^2, \theta)$$

Note: for per symbol $R(D)$ expressions for R & D need to be divided by m .

where θ is selected so that $\sum_{i=1}^n D_i = D$.



$$D = \sum_{i=1}^n \min(\sigma_i^2, \theta)$$

Example of JPEG and MPEG.

Correlated Gaussian Source

Assume that $(X_1, \dots, X_n) = \underline{X}^n \sim N(0, K_n)$

Since K_n can be written as

$$K_n = U \Lambda_n U^{-1}$$

where $\Lambda_n = \text{diag}[\lambda_1, \dots, \lambda_n]$

$$\underline{Y}^n = U^{-1} \underline{X}^n$$

also, $\hat{\underline{Y}}^n = U^{-1} \hat{\underline{X}}^n$

Since U^{-1} is invertible $I(\underline{X}^n; \hat{\underline{X}}^n) = I(\underline{Y}^n; \hat{\underline{Y}}^n)$

It is also easy to show that $d(\underline{Y}^n, \hat{\underline{Y}}^n) = d(\underline{X}^n, \hat{\underline{X}}^n)$

So, the rate distortion function for X is the same as the rate distortion function for Y .

$$R(D) = \frac{1}{n} \sum_{i=1}^n \left(\frac{1}{2} \log \frac{\lambda_i}{\theta} \right)^+ = \frac{1}{n} \sum \max\left(\frac{1}{2} \log \frac{\lambda_i}{\theta}, 0 \right)$$

$$D = \frac{1}{n} \sum_{i=1}^n \min(\lambda_i, \theta)$$

Toeplitz distribution Theorem

Let K_∞ be an infinite Toeplitz matrix with entry K_m on the m th diagonal. The eigenvalues of K_∞ are contained in the interval $\delta \leq \lambda \leq \Delta$

where δ and Δ denote the essential infimum and supremum, respectively of the function

$$S(\omega) = \sum_{m=-\infty}^{\infty} K_m e^{-jm\omega}$$

Moreover, if both δ and Δ are finite and $G(\lambda)$ is any ^{continuous} function of $\lambda \in [\delta, \Delta]$, then

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{m=1}^n G(\lambda_m^{(n)}) = \frac{1}{2\pi} \int_{-\pi}^{\pi} G[S(\omega)] d\omega$$

where $\lambda_m^{(n)}$ are the eigenvalues of the n -th order matrix K_n centred about the main diagonal of K_∞ .

(See Grenander & Szegő: Toeplitz forms and their applications. See also T. Berger: Rate - Distortion Theory).

Applying this theorem to the Equations in the previous page, we find the rate-distortion of a discrete-time stationary source as:

$$D_\theta = \frac{1}{2\pi} \int_{-\pi}^{\pi} \min[\theta, S(\omega)] d\omega$$
$$R(D_\theta) = \frac{1}{4\pi} \int_{-\pi}^{\pi} \max[0, \log \frac{S(\omega)}{\theta}] d\omega$$

where θ traverses the interval $0 \leq \theta \leq \Delta$.

Example: $R(D)$ for a first-order Gauss-Markov source.

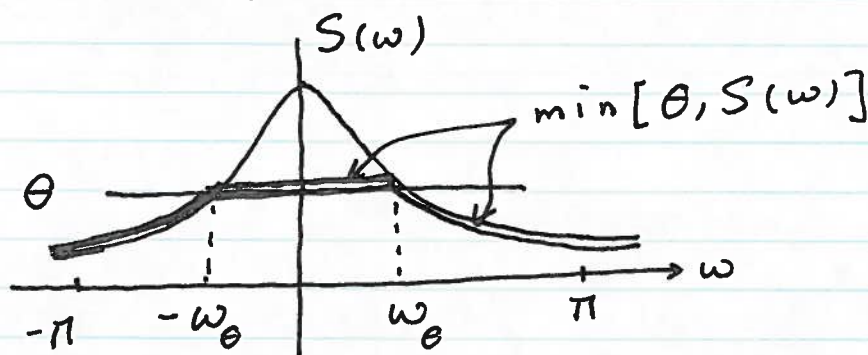
$$X_n = \rho X_{n-1} + Z_n$$

where Z_n is zero-mean Gaussian

$$K_m = \sigma^2 |\rho|^m$$

So:

$$S(\omega) = \sigma^2 \sum_{m=-\infty}^{\infty} |\rho|^m e^{-jm\omega} = \frac{\sigma^2(1-\rho^2)}{1-2\rho\cos\omega + \rho^2}$$



For small values of D , when;

$$\theta \ll S(\pi) = \frac{\sigma^2(1-\rho)}{1+\rho}$$

we have,

$$D = \frac{1}{2\pi} \int_{-\pi}^{\pi} \theta d\omega = \theta$$

and

$$R(D) = \frac{1}{4\pi} \int_{-\pi}^{\pi} \log \frac{\sigma^2(1-\rho^2)}{\theta(1+\rho^2-2\rho\cos\omega)} d\omega = \frac{1}{2} \log \frac{(1-\rho^2)}{D} \sigma^2$$

$$\int_0^{\pi} (1+\rho^2-2\rho\cos\omega) d\omega = 2\pi \ln 1 = 0 \rightarrow$$

So, for a Gauss-Markov source

$$R(D) = \frac{1}{2} \log \frac{(1-p^2)\sigma^2}{D} \quad D \leq \frac{1-p}{1+p} \sigma^2$$

for larger values of D , we need to vary θ from zero to maximum of $S(\omega)$ and derive the $R(D)$ parametrically.

SQNR:

$$\frac{\sigma^2}{D} = \frac{2^{2R}}{(1-p^2)}$$

$$\begin{aligned} \text{SQNR} &= 10 \log \frac{\sigma^2}{D} = 20R \log 2 - 10 \log(1-p^2) \\ &= 6R - 10 \log(1-p^2) \end{aligned}$$

Example: $p = 0.9$ and $R = 1$ bit/sample

$$\text{SQNR} = 6 + 7.2 = 13.2 \text{ dB}$$

results using vector Quantization and Finite State (Feed back) VQ.

k	VQ (SQNR)	FSVQ
1	4.4	10
2	7.9	10.8
3	9.2	11.4
4	10.2	12.1
5	10.6	
6	10.9	
7	11.2	

Characterization of the Rate-Distortion Function

The problem is to minimize

$I(x; \hat{x})$ subject to $\sum_x \sum_{\hat{x}} p(x) q(\hat{x}|x) d(x, \hat{x}) \leq D$
by proper choice of $\{q(\hat{x}|x)\}$. We use the method
of Lagrange multipliers.

$$\begin{aligned} J(q) &= \sum_x \sum_{\hat{x}} p(x) q(\hat{x}|x) \log \frac{q(\hat{x}|x)}{\sum_x p(x) q(\hat{x}|x)} \\ &\quad + \lambda \sum_x \sum_{\hat{x}} p(x) q(\hat{x}|x) d(x, \hat{x}) \\ &\quad + \sum_x v(x) \sum_{\hat{x}} q(\hat{x}|x). \end{aligned}$$

$$\begin{aligned} \frac{\partial J}{\partial q(\hat{x}|x)} &= p(x) \log \frac{q(\hat{x}|x)}{q(\hat{x})} + p(x) - \sum_{x'} p(x') q(\hat{x}|x') \frac{1}{q(\hat{x})} p(x) \\ &\quad + \lambda p(x) d(x, \hat{x}) + v(x) = 0. \end{aligned}$$

Let $v(x) = p(x) \log \mu(x)$ to get

$$p(x) \left[\log \frac{q(\hat{x}|x)}{q(\hat{x})} + \lambda d(x, \hat{x}) + \log \mu(x) \right] = 0$$

or

$$q(\hat{x}|x) = \frac{q(\hat{x}) e^{-\lambda d(x, \hat{x})}}{\mu(x)} \quad (A)$$

Summing over \hat{x} , we get

$$\sum_{\hat{x}} q(\hat{x}|x) = 1 = \frac{\sum_{\hat{x}} q(\hat{x}) e^{-\lambda d(x, \hat{x})}}{\mu(x)}$$

or

$$\mu(x) = \sum_{\hat{x}} q(\hat{x}) e^{-\lambda d(x, \hat{x})}$$

or,

$$q(\hat{x}|x) = \frac{q(\hat{x}) e^{-\lambda d(x, \hat{x})}}{\sum_{\hat{x}} q(\hat{x}) e^{-\lambda d(x, \hat{x})}}$$

Multiplying (A) by $p(x)$ and summing over x , we get

$$q(\hat{x}) = \sum_x p(x) q(\hat{x}|x) = q(\hat{x}) \sum_x \frac{p(x) e^{-\lambda d(x, \hat{x})}}{\mu(x)}$$

if $q(\hat{x}) \neq 0$, we get

$$\sum_x \frac{p(x) e^{-\lambda d(x, \hat{x})}}{\mu(x)} = 1$$

$$\Rightarrow \sum_x r(x) e^{-\lambda d(x, \hat{x})} = 1 \text{ for all } \hat{x}$$

in general, when some $q(\hat{x}) = 0$, we get

$$\sum_x \frac{p(x) e^{-\lambda d(x, \hat{x})}}{\mu(x)} \begin{cases} = 1 & q(\hat{x}) > 0 \\ \leq 1 & q(\hat{x}) = 0 \end{cases}$$

Parametric form of $R(D)$:

$$\begin{aligned}
 R(D) &= \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log \frac{q(\hat{x}|x)}{q(\hat{x})} \\
 &= \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log \frac{e^{-\lambda d(x, \hat{x})}}{\mu(\hat{x})} \\
 &= \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log \frac{P(x)}{P(\hat{x})} e^{-\lambda d(x, \hat{x})} \\
 &= + \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log P(x) \\
 &\quad + \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log P(\hat{x}) \\
 &\quad - \lambda \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) d(x, \hat{x})
 \end{aligned}$$

$$\begin{aligned}
 &= \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) \log r(x) \\
 &\quad + H(X) - \lambda D
 \end{aligned}
 \quad = \sum_x P(x) \log r(x) + H(X) - \lambda D$$

and

$$\begin{aligned}
 D &= \sum_x \sum_{\hat{x}} P(x) q(\hat{x}|x) d(x, \hat{x}) \\
 &= \sum_x \sum_{\hat{x}} P(x) \frac{q(\hat{x}) r(x)}{P(\hat{x})} e^{-\lambda d(x, \hat{x})} d(x, \hat{x})
 \end{aligned}$$

$$D = \sum_x \sum_{\hat{x}} q(\hat{x}) r(x) e^{-\lambda d(x, \hat{x})} P(x, \hat{x})$$

$$R = \sum_x P(x) \log r(x) + H(X) - \lambda D$$

$$\text{let } \frac{p(x)}{\mu(x)} = r(x)$$

then, when all $q(\hat{x}) \neq 0$, we get

$$\sum_x r(x) e^{-\lambda d(x, \hat{x})} = 1$$

define a matrix A with elements $e^{-\lambda d(x, \hat{x})}$

$$\textcircled{1} \quad A^T \cdot \underline{r} = \underline{\Pi} \quad \text{where } \underline{\Pi} = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

and

$$\textcircled{2} \quad \underline{A} \cdot \underline{Q} = \underline{M} \quad M = \begin{bmatrix} \mu(1) \\ \vdots \\ \mu(n) \end{bmatrix}$$

$$Q = \begin{bmatrix} q(1) \\ \vdots \\ q(n) \end{bmatrix}$$

From $\textcircled{1}$, we can find

$\{r(x)\}$ and from $r(x) = \frac{p(x)}{\mu(x)}$ find $\mu(x)$

and then from $\textcircled{2}$ find $q(1) \dots, q(n)$.

Example: Binary memoryless source with $p(0) = p$

$p(1) = 1 - p$ and $d(x, \hat{x}) = \delta_{x, \hat{x}}$.

$$\underline{d} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$A = \left\{ e^{-\lambda d(x, \hat{x})} \right\} = \begin{bmatrix} 1 & e^{-\lambda} \\ e^{-\lambda} & 1 \end{bmatrix} = \begin{bmatrix} 1 & a \\ a & 1 \end{bmatrix}$$

$$A^T \cdot \underline{r} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix} \begin{bmatrix} r(0) \\ r(1) \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$\begin{bmatrix} r_1 \\ r_2 \end{bmatrix} = \frac{1}{1-\alpha^2} \begin{bmatrix} 1 & -\alpha \\ -\alpha & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} \frac{1}{1+\alpha} \\ \frac{1}{1+\alpha} \end{bmatrix}$$

$$r(0) = \frac{1}{1+\alpha}, \quad r(1) = \frac{1}{1+\alpha}$$

$$r(0) = \frac{p(0)}{\mu(0)} \Rightarrow \mu(0) = p(1+\alpha)$$

$$r(1) = \frac{p(1)}{\mu(1)} \Rightarrow \mu(1) = (1-p)(1+\alpha)$$

$$\begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix} \begin{bmatrix} q(0) \\ q(1) \end{bmatrix} = \begin{bmatrix} p(1+\alpha) \\ (1-p)(1+\alpha) \end{bmatrix}$$

$$q(0) = \frac{p - \alpha(1-p)}{1-\alpha}$$

$$q(1) = \frac{-p\alpha + (1-p)}{1-\alpha}$$

Now,

$$D = r(0) q(0) \cancel{\alpha} \times 0 + r(0) q(1) \alpha + r(1) q(0) \alpha + r(1) q(1) \cancel{\alpha} \times 0 = \frac{1}{1+\alpha} \times \frac{-p\alpha + (1-p)}{1-\alpha} + \frac{1}{1+\alpha} \times \frac{p - \alpha(1-p)}{1-\alpha}$$

$$D = \frac{\alpha}{1+\alpha} \Rightarrow \alpha = \frac{D}{1-D}$$

$$R(D) = \sum_x p(x) \log r(x) + H(X) - \lambda D$$

$$= p(0) \log r(0) + p(1) \log r(1) + H_b(p) - \lambda D$$

$$= -p \log(1+d) - (1-p) \log(1+d) + H_b(p) - \lambda D$$

$$= -\log(1+d) + H_b(p) - \lambda D$$

$$= -\log \frac{1}{1-d} + H_b(p) - \lambda D$$

$$= \log(1-d) + H_b(p) + \log d D$$

$$= \log(1-d) + H_b(p) + D \log \frac{D}{1-d}$$

$$= \log(1-d) + H_b(p) + D \log D - D \log(1-d)$$

$$= H_b(p) - H_b(D)$$

Exercise: Consider a source with four letters with probability vector $\underline{p} = \frac{1}{2} \begin{bmatrix} p \\ 1-p \\ 1-p \\ p \end{bmatrix}$

and the distortion measure

$$\underline{d} = \begin{bmatrix} 0 & 1/2 & 1/2 & 1 \\ 1/2 & 0 & 1 & 1/2 \\ 1/2 & 1 & 0 & 1/2 \\ 1 & 1/2 & 1/2 & 0 \end{bmatrix}$$

Show that:

$$R(D) = \log 2 - H_b(p) - 2 H_b(D)$$

$$0 \leq D \leq D_1 = \frac{1}{2} (1 - \sqrt{1-2p})$$

and

$$R(D) = \eta(1-p) - \frac{1}{2} [\eta(2D-p) + \eta(2(1-D)-p)]$$

for

$$D_1 \leq D \leq D_{\max} = \frac{1}{2}$$

and

$$\eta(x) \triangleq -x \log x.$$

The above formulation provides us with a method for finding a lower bound on $R(D)$

Assume that $\{r(x)\}$, $x \in \mathcal{X}$ satisfy

$$\sum_x r(x) e^{-\lambda d(x, \hat{x})} \leq 1 \quad \text{all } \hat{x}$$

Then

$$R(D) \geq H(X) + \sum_x p(x) \log[r(x)] - \lambda D = R_L(D)$$