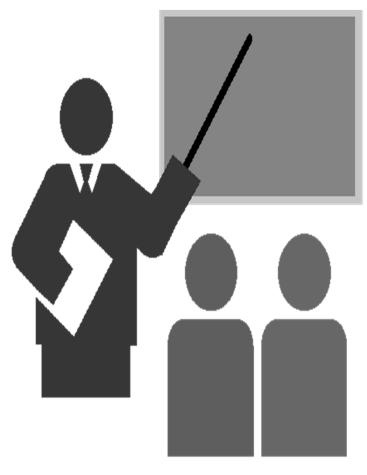
# Virtualization Technologies (ENCS 691K – Chapter 3)

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# The Key Technologies on Which Cloud Computing Relies



Web Services

Virtualization

# References (Virtualization)

- 1. . M. Pearce et al., Virtualization: Issues, Security, Threats, and Solutions, ACM Computing Survey, February 2013
- 2. . A. Khan et al., Network Vitrtualization: A Hypervisor for the Internet?, IEEE Communications Magazine, January 2012
- 3. . N.M Chowdhury and r. Boutaba, Network Virtualization: State of the Art and Research Challenges, IEEE Communications Magazine, July 2009
- 4. J. Carapinha et al., Network Virtualization A View from the Bottom, VISA '09 Proceedings of the 1st ACM workshop on Virtualized infrastructure systems and architectures, Pages 73-80



## Virtualization



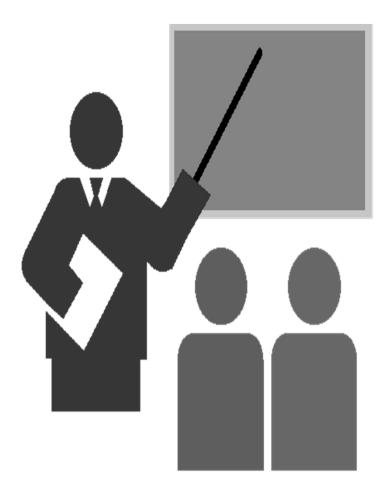
## **Outline**



1. Systems virtualization

2. Network virtualization

# On Systems Virtualization



Key concepts

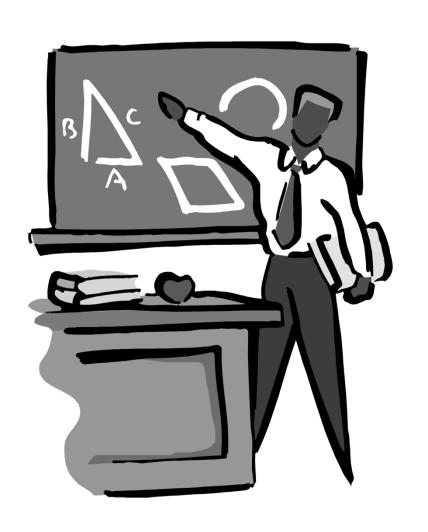
Type I (bare metal) vs. Type 2 (hosted)

Full virtualization vs. para-virtualization

# References (Systems Virtualization)

- 1. M. Pearce et al., Virtualization: Issues, Security, Threats, and Solutions, ACM Computing Survey, February 2013
- 2. P. Barham et al., XEN and the Art of Virtualization, SOSP '03 Proceedings of the nineteenth ACM symposium on Operating systems principles, Pages 164-177

# **Basic concepts**



- 1. On operating systems
- 2. Virtual machine, virtual machine monitor/hypervisor
- 4. Examples of benefits

# **Operating systems**

#### Some of the motivations

- Only one single tread of CPU can run at a time on any single core consumer machine
- Machine language is tedious

# **Operating systems**

Operating systems bring a level of abstraction on which multiple processes can run at a time – Deal among other things with:

- Multiplexing
- Hardware management issues

However only one operating system can run on a bare single core consumer machine

# virtual machines and hypervisors

Systems virtualization dates back to the 60s

- IBM experimentation with "time sharing systems"
  - Need for virtual machines to test how applications / users can time share a real machine

# virtual machines and hypervisors

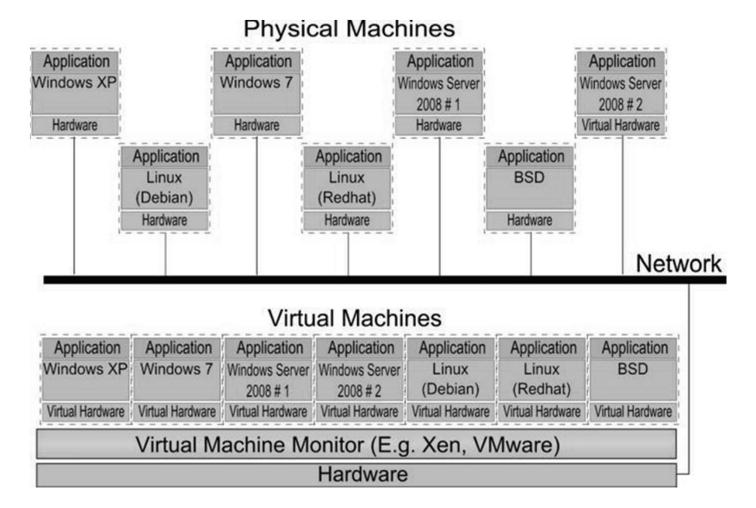
Virtual machine (VM) (sometimes called virtual hardware)

- Software that provides same inputs / outputs and behaviour expected from hardware (i.e. real machine) and that supports operations such as:
  - Create
  - Delete
  - Migrate
  - Increase resources

Virtual machine monitor (also called hypervisor)

 Software environment that enables operations on virtual machines (e.g. XEN, VMWare) and ensures isolation

# virtual machines, hypervisors



From reference [1] – Note: There is a small error in the figure

# **Examples of Benefits**

All benefits are due to the possibility to manipulate virtual machine (e.g. create, delete, increase resources, migrate), e.g.

- Co-existence of operating systems
- Operating systems research
- Software testing and run-time debugging
- Optimization of hardware utilization
- Job migration

# **Advanced concepts**



- 1. Bare metal vs. hosted hypervisor
- 2. Full virtualization vs. Paravirtualization

#### Types of hypervisor

- Type I bare metal
  - Installed on bare hardware
  - Examples
    - Citrix XEN server
    - VMWARE ESX/ESXI

#### Types of hypervisor

- Type 2 hosted
  - Runs on top of host operating system
  - Examples:
    - VMWare workstation
    - VirtualBox

#### Type I - Bare metal

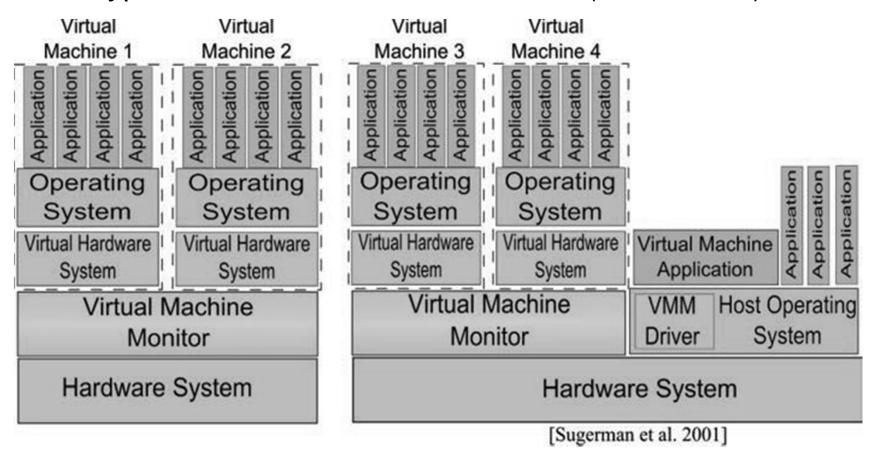
- Hypervisor installed on bare hardware
  - Advantages (compared to type II)
    - Performance (No additional software layer to go through)
    - Security (No possible attack through host operating system)
  - Drawbacks (compared to type II)
    - Host operating system needs to be "ported" on top of hypervisor
    - Complexity depends on the type of virtualization (Full virtualization vs. para-virtualization)

#### Type II - Hosted

- Hypervisor installed on top of host operating system
  - Drawbacks (compared to type I)
    - Performance (need to go through host operating system)
    - Security (i.e. Possibility to attack through host operating system)
  - Advantages (compared to type I)
    - Host operating system is re-used as it is (No need to port it)
    - No change required to applications running on top of host operating system

# Type I vs Type II Hypervisor (Summary)

Types of hypervisor/virtual machine monitor (From ref. 2)



#### More on operating systems fundamentals

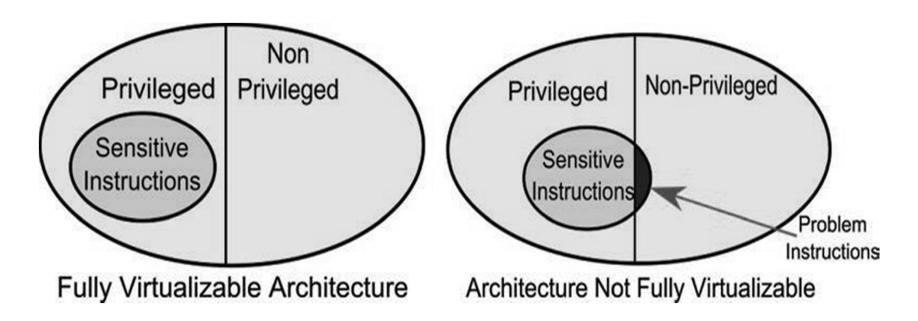
- Privileged vs. non privileged instruction
  - Privileged
    - If called in user mode, the CPU needs to trap it and switch control to supervisory software (e.g. hypervisor) for its execution

#### More on operating systems fundamentals

- Sensitive vs. non sensitive instruction
  - Sensitive
    - Has the capacity to interfere with supervisor software functioning (e.g. Hypervisor)
      - Write hypervisor memory vs. read hypervisor memory

### Could all CPU architectures be fully virtualized?

 Could be fully virtualized only if the set of sensitive instructions is a subset of the privileged instructions



From reference [1]

### Could all CPU architectures be fully virtualized?

- The case of Intel x86 CPU architectures
  - Cannot be fully virtualized
    - "Certain instructions must be handled by the VMM for correct virtualization, but these with insufficient privilege fail silently rather than causing a convenient trap" – Reference [2]

#### **Definitions**

#### **Full virtualization**

- Hypervisor enables virtual machines identical to real machine
  - Problematic for architectures such as Intel x86

#### **Definitions**

#### Para-virtualization

- Hypervisor enables virtual machine that are similar but not identical to real machine
  - A solution to the problem of CPU architectures that cannot be virtualized
    - Prevents user programs from executing sensitive instructions
  - Note:
    - Para-virtualization is not the only solution to the problem

#### **Full virtualization**

- Advantages
  - Possibility to host guest operating systems with no change since virtual machines are identical to real machines
- Disadvantages
  - Not always feasible (e.g. Intel x86)
    - There are work around (e.g. binary translation)
  - Some guest operating systems might need to see both virtual resources and real resources for real time applications

#### Para - virtualization

- Advantages
  - Feasible for all CPU architectures
  - Performance Compared to:
    - Full virtualization
    - Other approaches to architectures that could not be virtualized (e.g. binary translation)
- Disadvantages
  - Need to modify guest operating systems

#### Para - virtualization

- Alternatives to para-virtualization
  - Binary translation (e.g. VMWare ESX server)
    - Leads to full virtualization
    - No need to re-write "statically" guest operating systems
      - i.e. guest OS can be installed without change
    - Interpretation of guest code (OS + application)
      - "Rewrites" dynamically guest code and insert traps when necessary

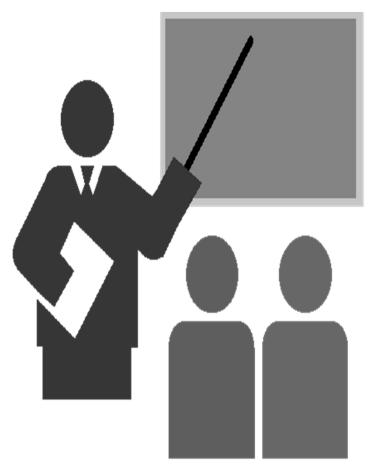
#### Para - virtualization

- Alternatives to para-virtualization
  - Binary translation
    - Disadvantages / penalties
      - Performance
      - However, optimization is possible, e.g.
        - » Adaptive translation (i.e. optimize the code being translated)

#### Para – virtualization

- A detailed case study on para-virtualization
  - XEN (Reference 2)

## On Network Virtualization



Motivations and basic components

Prior to network virtualization

A case stys

# References (Network Virtualization)

- 1. N.M Chowdhury and r. Boutaba, Network Virtualization: State of the Art and Research Challenges, IEEE Communications Magazine, July 2009
- 2. J. Carapinha et al., Network Virtualization A View from the Bottom, VISA '09 Proceedings of the 1st ACM workshop on Virtualized infrastructure systems and architectures, Pages 73-80
- 3. G. Schaffrat et al., Network Virtualization Architecture: Proposal and Initial Prototype, Proceeding VISA '09 Proceedings of the 1st ACM workshop on Virtualized infrastructure systems and architectures, Pages 63-72
- 4. J. Kurose and K. Ross, Computer Networking: A Top Down Approach, Pearson, 6<sup>th</sup> Edition, 2013
- 5. Venkateswanan, Virtual Private Networks, IEEE Potentials, Issue 20, no1

## On Network virtualization



1. Motivations

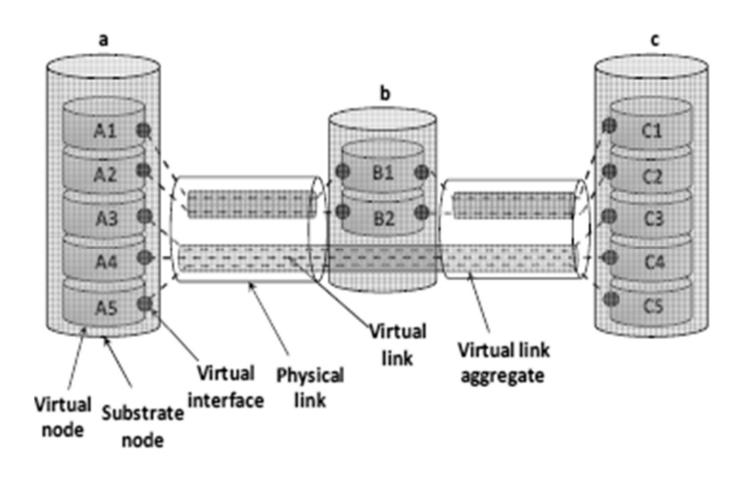
2. Basic components

## **Motivations**

Bring the benefits of systems virtualization to the networking world, e.g.

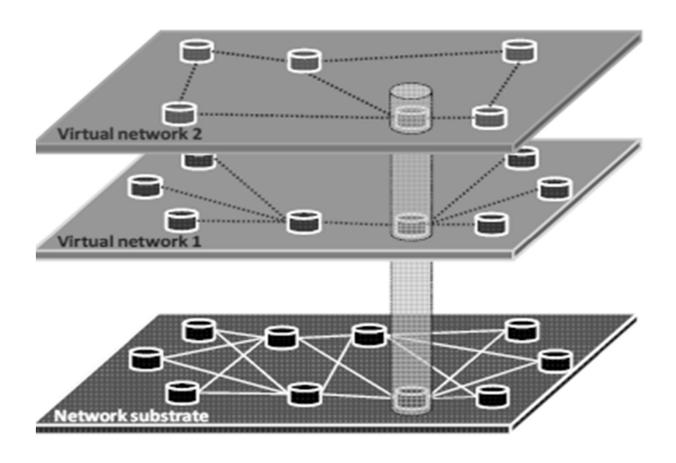
- Co-existence of virtual networks on top of a same real network
  - Note: Virtual Private Networks (VPNs) do not rely on virtualization and have several limitations
    - Different technologies and protocol stacks cannot be used for instance
- Networking research
- Optimization of networking resources utilization
  - Nodes
  - Links

# **Basic components**



From reference 9

### **Basic components**



From reference 9

#### On Network virtualization



1. Prior to network virtualization

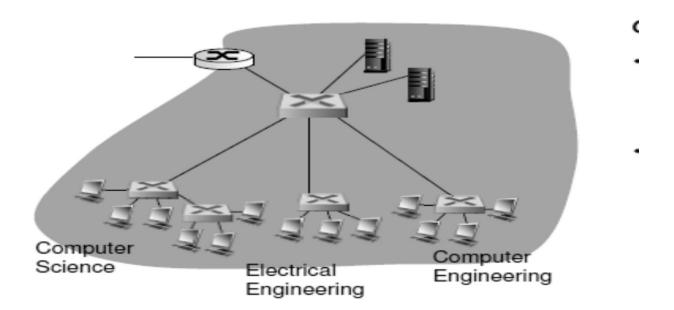
2. A Case study

#### Virtual Local Area Networks (VLANs)

- Possibility to define several VLANs over a same physical LAN infrastructure
  - Each VLAN has its broadcast domain and has an id.
- However
  - Each physical node is part of one and only VLAN
    - No efficient resource usage

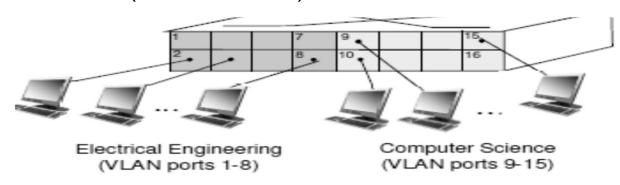
#### Virtual Local Area Networks (VLANs)

A LAN (Reference 4)

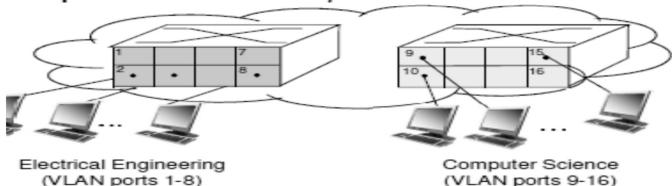


#### Virtual Local Area Networks (VLANs)

A VLAN (Reference 4)



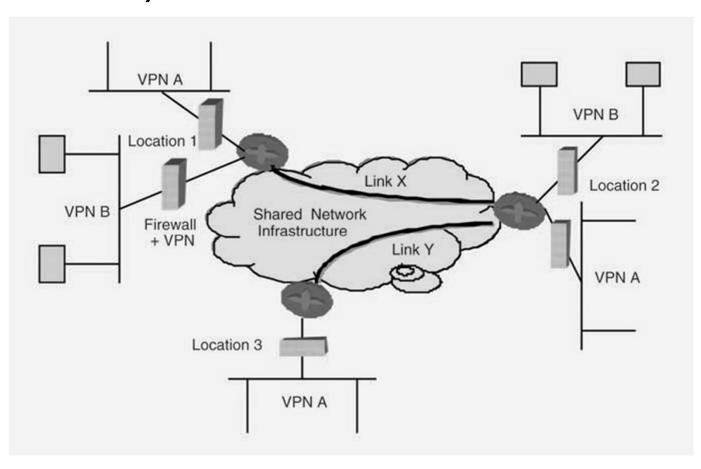
. operates as multiple virtual switches



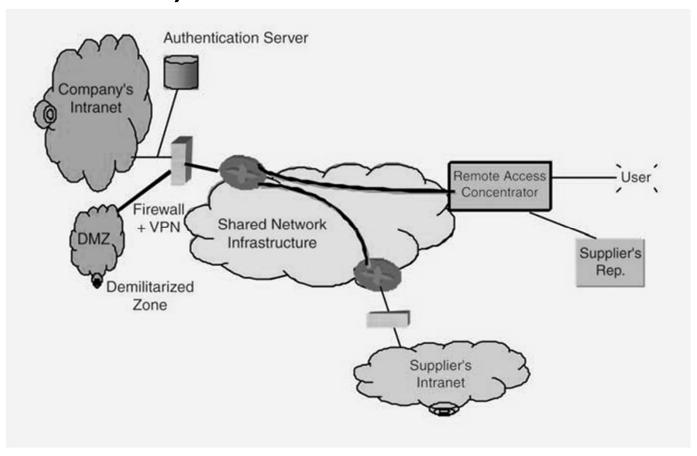
#### **Virtual Private Networks**

- Possibility to build virtual networks using a shared infrastructure (usually Internet, but might be a dedicated networks)
  - Site interconnection
  - Extranets
- But:
  - No real insolation between the different networks traffic over the shared infrastructure

### Virtual Private Networks – Reference 5 (LAN Interconnection)



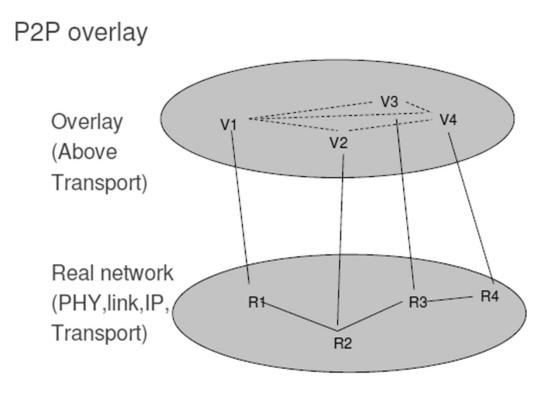
### Virtual Private Networks – Reference 5 (LAN Interconnection)



#### **Overlays**

- Logical networks built on top of real networks (e.g. skype)
- A same physical node might be part of several overlays
- But:
  - Overlays might interact in a harmful way
  - Used mainly at application layer and does not enable experimentation of lower layer protocols

#### **Overlays**



#### **Overlays**

#### P2P overlay

- Characteristics
  - own topology that may be different from the topology of the real network
  - Own protocols that may be different from the protocols used in the real network
  - May come with an application embedded in it (e.g. Skype) or as an infrastructure that can be used by other applications (e.g. CHORD)
  - APIs, toolkits are provided when the application is not embedded in the overlay

# A Case Study on Network Virtualization (Reference 3)

#### Business model of current Internet:

- Internet Service Providers (ISPs) (e.g. Bell, Rogers)
- Service Providers (eg. Google, Akamai)

## A Case Study on Network Virtualization Reference 3

New business model (4 roles):

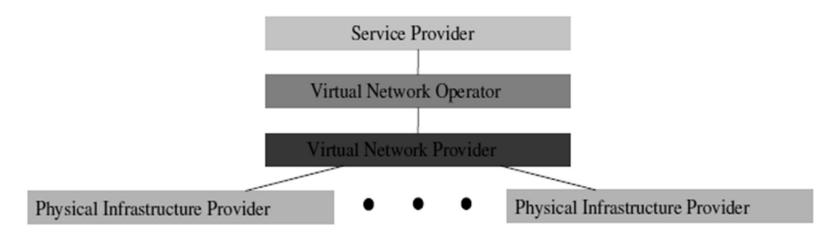
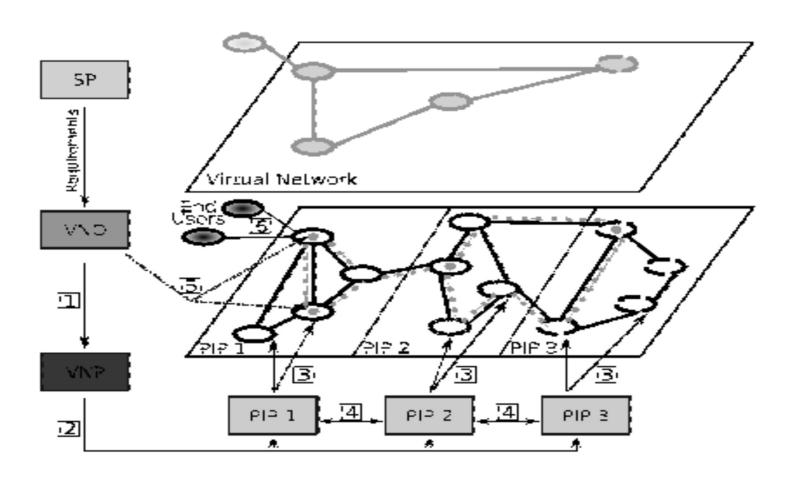


Figure 1: VNet Management and Business Roles

## A Case Study on Network Virtualization Reference 3

New business model (6 interfaces):



## A Case Study on Network Virtualization Reference 3

#### Simplified scenario

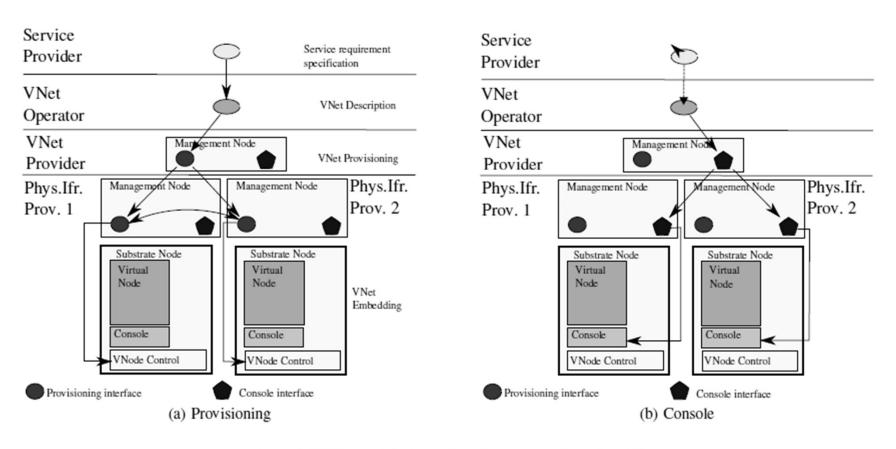


Figure 3: VNet provisioning (a) and console architecture (b).

# A Case Study on Network Virtualization (Reference 3)

#### Prototype

- Node level virtualization
  - XEN
- VNET description
  - XML

# A Case Study on Network Virtualization (Reference 3)

Topology used for Vnet instantiation measurements (end to end from Vnet request by service provider till full provisioning of VNET

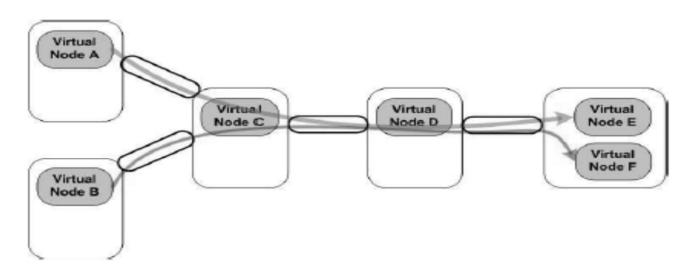


Figure 6: Experimental topology.

### The End



