

Cryptanalysis of FLEXAEAD

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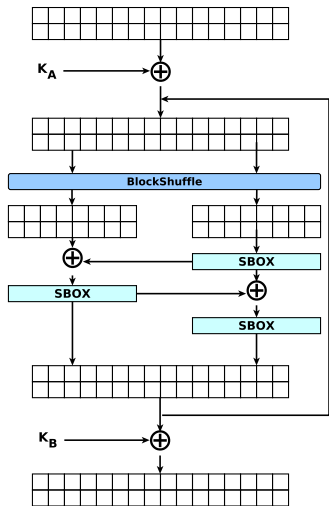


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Introduction

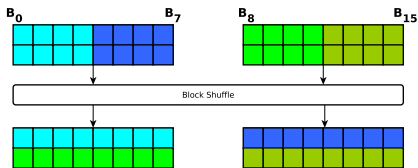
- ▶ FLEXAEAD is round 1 candidate of NIST LWC
- ▶ The underlying Blockcipher is *Internal Keyed Permutation*
- ▶ Block Size can be 64-bit, 128-bit or 256-bit
- ▶ Reported Key Recovery Attack for each variant
- ▶ The attacks are of two type
 1. Iterated Truncated Differential
 2. Yoyo Attacks

Internal *Keyed* Permutation of FLEXAEAD



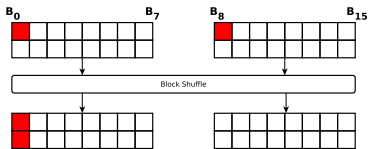
1. x -bit Flex state is called FLEX- x
2. FLEX-128 round function
3. State Bifurcation
4. AES Sbox is used
5. Repeated several times

BlockShuffle



Key Observations

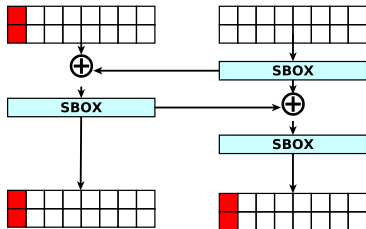
Effect of BlockShuffle



- ▶ Same Nibble in “Symmetric Bytes” transits to a single byte
- ▶ Number of active bytes can be decreased from two to one

Key Observations

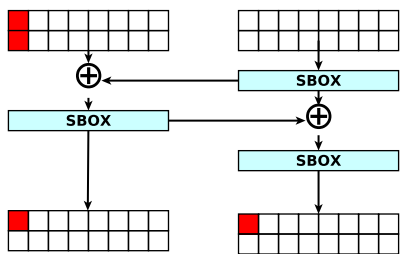
Effect of SBoxes



- ▶ Due to the effect of XOR, one active byte activates two bytes
- ▶ A pair of “Symmetric Byte” activates a pair of “Symmetric Byte”

Key Observations

Effect of SBoxes: Byte to Nibble Transition



- Only upper or lower nibbles of “Symmetric Bytes” are activated
- If initially a pair of “Symmetric Bytes” are active, this event occurs with equal probability

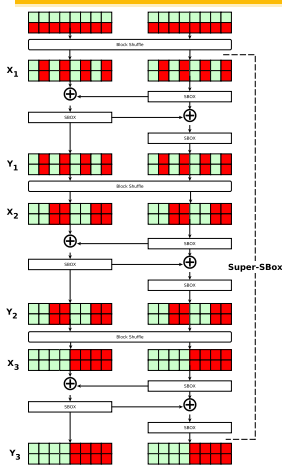
Exploiting AES Sbox

$$\left| \left\{ (x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0xf0 = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8} \right\} \right| = 4096$$
$$\left| \left\{ (x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0x0f = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8} \right\} \right| = 4096$$

With probability 2^{-7} two bytes transit to either upper or lower nibble

Key Observations

SuperSBox

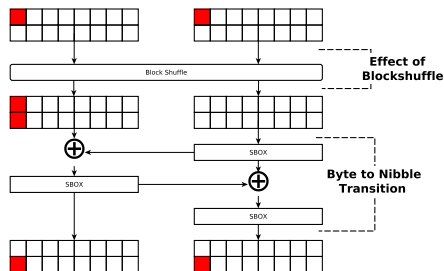


- ▶ Two Super-Sbox exists in $F_{\text{LEX-128}}$
- ▶ Initial BlockShuffle Layer is not considered in the Super-Sbox
- ▶ Super-Sbox spans over 2.5 round
- ▶ Each Super-Sbox is of 64-bit
- ▶ Super-Sbox in $F_{\text{LEX-64}}$ and $F_{\text{LEX-256}}$ spans over 1.5 and 3.5 round respectively

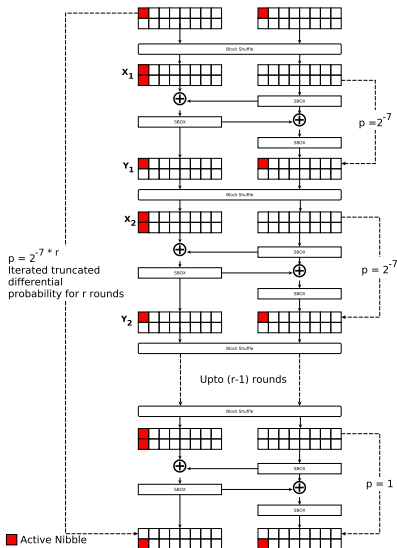
Iterated Truncated Differential

One Round Truncated Differential

- *Effect of BlockShuffle and Byte to Nibble Transition is Combined*
- The active nibbles in initial state and final state are in same position at the cost of 2^{-7}

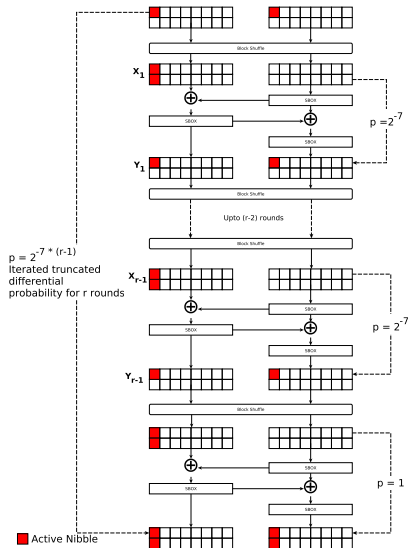


Iterated Truncated Differential



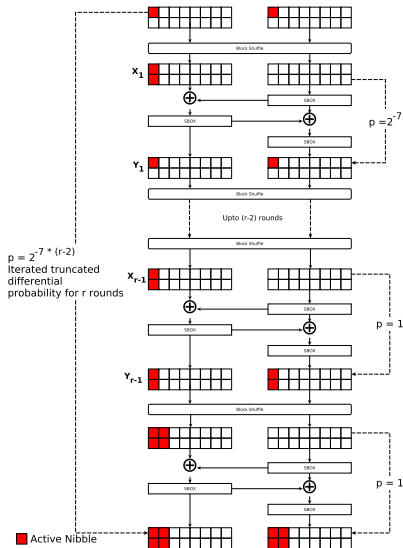
- ▶ The truncated differential can be iterated for r rounds
- ▶ Paying probability for r rounds
- ▶ Cost of the trail is $2^{-7 * r}$
- ▶ Some rounds at the end can be made free

Iterated Truncated Differential: Free Rounds=1



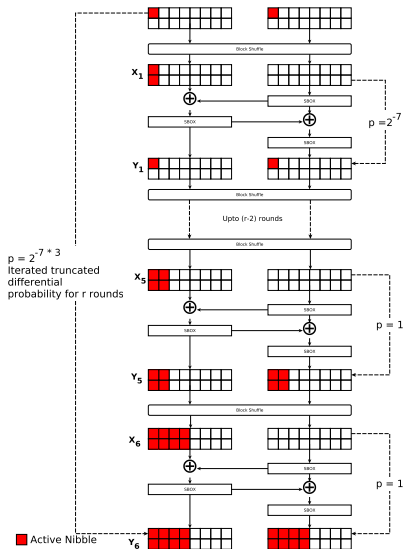
- ▶ 2 bytes are fully active
- ▶ Paying probability for $r - 1$ rounds
- ▶ Cost of the trail is $2^{-7*(r-1)}$

Iterated Truncated Differential: Free Rounds=2



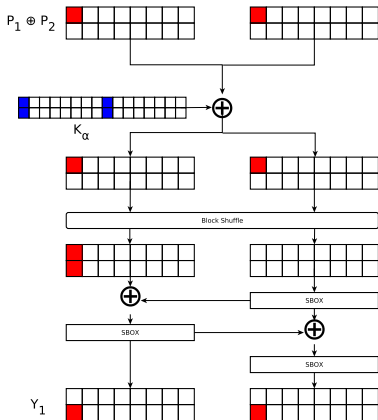
- ▶ 4 bytes are fully active
- ▶ Paying probability for $r - 2$ rounds
- ▶ Cost of the trail is $2^{-7*(r-2)}$

Iterated Truncated Differential: Distinguisher



- ▶ Number of free rounds is 3
- ▶ Probability of 6-round FLEX-128 distinguisher is 2^{-7*3}
- ▶ In similar way, number of free rounds in 5-round FLEX-64 and 7-round FLEX-256 is 2 and 4 respectively

Iterated Truncated Differential: Key Recovery



- ▶ Find a right pair (P_1, P_2) , such that difference is in byte 0 and 8
- ▶ Guess Key byte 0 and 8 (2^{16} possible guesses)
- ▶ Run one round encryption and check whether same of byte 0 and 8 are active or not in Y_1 (2^9 key candidates remain)
- ▶ Use two more right pairs to reduce key candidates to 1
- ▶ Repeat the procedure for 8 more byte pairs

Iterated Truncated Differential Attacks: Summary

Block Size	#rounds	Data Complexity		Time Complexity	Memory Complexity
		Encs	Decs	MA's	
64	7	$2^{30.5}$		$2^{34.5}$	$2^{18.5}$
128	16	$2^{93.5}$		$2^{108.5}$	$2^{20.5}$
256	21	$2^{109.5}$		$2^{125.5}$	$2^{22.5}$

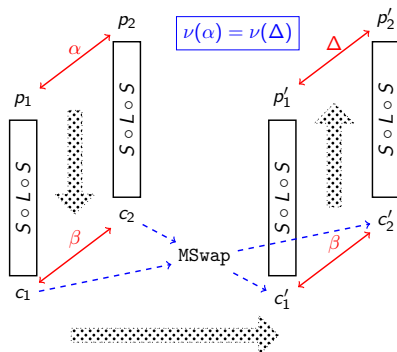
Yoyo Attacks



$$G'_2 = L \circ S \circ L \circ S$$

Two full generic Rounds

$$G_2 = S \circ L \circ S \quad \leftarrow \text{Dropping final linear layer (to simplify)}$$



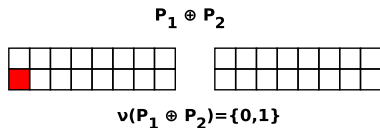
- ν is the Zero Difference Pattern

Applied to AES

- First key-independent Yoyo distinguishers of AES
- 5-round Key Recovery

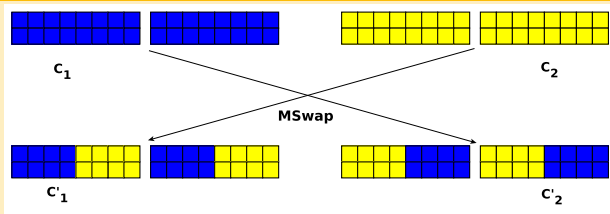
The Yoyo Trick

Zero Difference Pattern



- ▶ Two Super-Sbox in FLEX-128 state
- ▶ A fully inactive Super-Sbox is denoted by 1; otherwise, 0

MSwap



- ▶ Bytes are swapped between two texts according Super-Sbox output

Yoyo Attacks: Deterministic Distinguisher



$P_1 \oplus P_2$



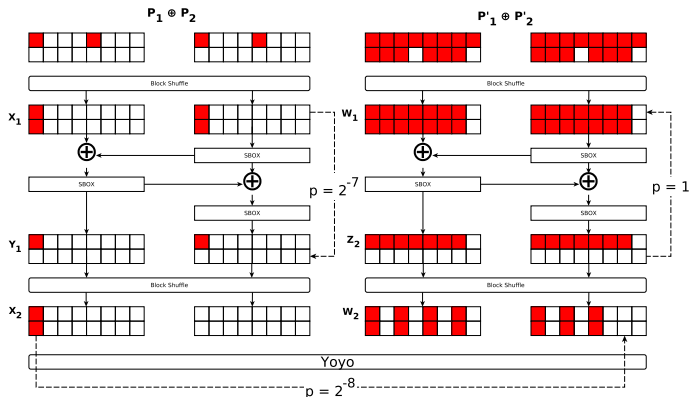
$P'_1 \oplus P'_2$

- ▶ Super-Sbox and BlockShuffle are considered as S and L layer respectively
- ▶ FLEX-128 Super-Sbox spans over 2.5 rounds
- ▶ 6-round FLEX-128 Deterministic Distinguisher
- ▶ Apply Yoyo game
 1. $P_1, P_2 \xrightarrow{ENC} C_1, C_2$
 2. $C_1, C_2 \xrightarrow{MSwap} C'_1, C'_2$
 3. $C'_1, C'_2 \xrightarrow{DEC} P'_1, P'_2$

Yoyo Attacks: Key Recovery

- ▶ 6-round Deterministic Distinguisher is the building block of 7-round FLEX-128 Key Recovery attack
- ▶ Byte to Nibble Transition is used to extend for 1 round
- ▶ Similar kinds of attacks exist for FLEX-64 and FLEX-256

Yoyo Attacks: Key Recovery



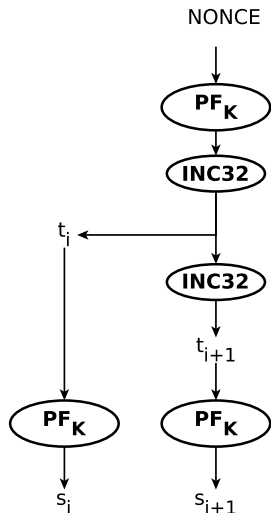
- Choose P_1, P_2 and encrypt them to obtain C_1, C_2
- Apply *MSwap* on C_1, C_2 and decrypt them to get P'_1, P'_2
- Any one of the 8 active Bytes in W_2 can be zero w.p. 2^{-5}
- Trail probability is 2^{-12}
- Key Recovery part is same as Iterated Truncated Differential

Yoyo Attacks: Summary

Block Size	#rounds	Data Complexity		Time Complexity	Memory Complexity
		Encs	Decs	MAAs	
64	5	2^{10}	$2^{16.5}$	$2^{15.5}$	2^{10}
128	7	$2^{10.5}$	$2^{16.5}$	$2^{16.5}$	$2^{11.5}$
256	9	2^{11}	$2^{16.5}$	$2^{17.5}$	2^{13}

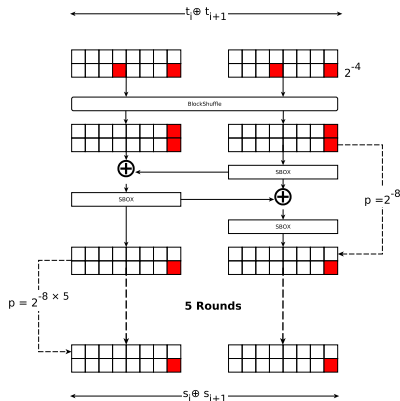
Forgery Attacks

Sequence Generation Step



- ▶ Sequence of bits are used for AE
- ▶ PF_k is used for sequence generation
- ▶ INC32 acts as XOR with probability 2^{-1}
- ▶ Last call to PF_k of two consecutive numbers differ by INC32

Differential Trail of Sequence Generation



- Differential Characteristics for Sequence Generation of FLEXAEAD-128
- Difference in Plaintext or Associated Data cancels out the difference in $S_i \oplus S_{i+1}$ with probability 2^{-8}

Forgery Attacks on FLEXAEAD

Scheme	Complexity
FLEXAEAD-64	2^{50}
FLEXAEAD-128	2^{60}
FLEXAEAD-256	2^{80}

Conclusion

1. Reported Iterated Truncated Differential which exploits AES Sbox and BlockShuffle operation
2. Generalized Yoyo Distinguishing Attack is applicable
3. All attacks are exploited to recover subkeys
4. Practical ones are experimentally verified
5. FLEXAEAD is out of 2nd round

Thank You